Logic and Computability Lecture 5



Introduction to Z3

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What is Z3?

Solver for Satisfiability Modulo Theories

What is Z3?

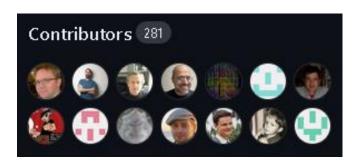
- Solver for Satisfiability Modulo Theories
 - we know how to check satisfiability
 - ... until now: Only propositional logic!
- Z3 allows us to efficiently answer decision problems including
 - Integers, Reals, Arithmetic
 - BitVectors, uninterpreted Functions, Arrays,
 - etc.
- More on Theories starting from next week
- Today: Basics Principles of Z3 and First Problems

Background

- Developed by Microsoft Research
 - https://github.com/Z3Prover/z3

Christoph Lev Leonardo Wintersteiger Nachmanson de Moura

Nikolaj Bj**ø**rner



■ SMT-LIB2 - A standardized language for Problems in SMT

Principles

- Is $\neg a \land (a \lor b)$ satisfiable?
- What do we need to describe a problem for the solver?
 - Variables (of a specific Sort),(declare-const a Bool)

(declare-const b Bool)

Constraints, and

```
(assert (not a) )
(assert (or a b) )
```

Checking for Satisfiability

```
(check-sat)
```

A Simple Example in SMT-LIB2

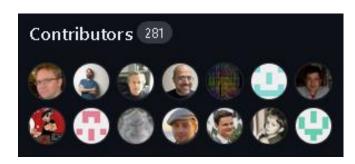
```
(declare-const a Bool)
(declare-const b Bool)
(assert (not a) )
(assert (or a b) )
(check-sat)
(get-model)
```

Background

- Developed by Microsoft Research
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- SMT-LIB2 A standardized language for Problems in SMT
- API for C++, Python, Julia, etc.

Installing

- We will use the Python API:
 - pip install z3-solver
- Optionally, you may install z3 natively:
 - sudo apt-get install z3 (Via aptitude for Ubuntu, etc.)
 - https://www.nuget.org/packages/Microsoft.Z3/ (Windows)
 - https://jfmc.github.io/z3-play (online)

Python API

- User-friendly interface for SMT-LIB2
- Used in the Programming Assignment
- Variables (of a specific Sort),

Constraints, and

```
(assert (not a) )
(assert (or a b) )
```

Checking for Satisfiability (check-sat)

```
solver = Solver()
solver.add(Not(a))
solver.add(Or(a,b))
```

solver.check()

Python API

```
from z3 import *
a, b = Bools("a b")
solver = Solver()
solver.add(Not(b))
solver.add(Or(a,b))
print(solver.sexpr())
result = solver.check()
model = solver.model()
print(result)
print(model)
```

Python API

Constraints

- Provides Methods for Connectives:
 - And(), Or(), Not(), Implies(), ==, ^, etc.
- Method to check whether two statements can be distinct:
 - Distinct(a,b)
- Operator overloading:
 - +, -, >>, <<, etc.
- Reference: https://z3prover.github.io/api/html/namespacez3py.html

A First Example

- We want to show that the following statements are equal:
 - $p \rightarrow q$
 - $\blacksquare \neg p \lor q$

A First Example

```
p \rightarrow q == \neg p \lor q?
 from z3 import *
 solver = Solver()
 a, b = Bools("a b")
1, r = Bools("l r")
 solver.add(l == Implies(a, b))
 solver.add(r == Or(Not(a), b))
 solver.add(Distinct(r,1) )
 result = solver.check()
print(result)
```

Back to SMT-LIB2

```
p \rightarrow q == \neg p \lor q?
from z3 import *
solver = Solver()
a, b = Bools("a b")
l, r = Bools("l r")
 solver.add(l == Implies(a, b))
 solver.add(r == Or(Not(a), b))
 solver.add(Distinct(r,1))
print(solver.sexpr())
result = solver.check()
print(result)
```

BitVectors

- Z3 allows us to use so-called theories
- We have a first look at bitvectors
- Syntax:
 - bv = BitVector("bv", <size>)
- BitVectors respect under-/overflow behaviour!
 - In contrast to Z3's integers

Operations on BitVectors

- The BitVector Sort respects overloaded operators:
 - <,>, <=, +, -, <<, >>, /, etc.
 - Caution: These are signed interpretations
 - Use ULT, UGT, ULE for unsigned interpretations

Equivalence Checking for BitVectors

- We want to prove the equivalence of the following
 - (((y & x) * -2) + (y + x))
 - x ⊕ y

Weird XOR

```
from z3 import *
x = BitVec('x', 32)
y = BitVec('y', 32)
output = BitVec('output ', 32)
s = Solver()
s.add(x^y==output)
s.add(Distinct(((y \& x) * -2) + (y + x),output))
print(s.check())
```

Operations on BitVectors

- The BitVector Sort respects overloaded operators:
 - <,>, <=, +, -, etc.</p>
 - Caution: These are signed interpretations
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- Overflow and Underflow
 - BVAddNoOverflow, BVAddNoUnderflow
 - BVMulNoOverflow, BVMulNoUnderflow

Overflow Behaviour

- We want to check whether the statement TODO
 - \bullet (x + 1 < x 1)

Variables in a Satisfying Model

- Variables and Expressions are stored in z3-specific classes
- We can use solver.model().decls() to iterate through all declared variables
 - Use .as_long() to convert a BitVector to a Python Integer

```
model = solver.model()
for var in solver.model.decls():
    print(f"{var}: {model[var]}(:{type(model[var])})")
```

Overflow Behaviour

- We want to check whether the statement TODO
 - \bullet (x + 1 < x 1)

- We need to add
 - BVNoOverflow(x, 1, True)
 - BVNoUnderflow(x, 1, True)
- Functions that evaluate to False when Over-/Underflow would occur in the model

Assignment Sheet

- 4 Exercises + 1 Bonus Exercise
- You are allowed to work in groups of 2
 - If you do so, please add your information into the README
- Deadline: 05. 06. 2024

Outline - Part II

- IntSort + Z3 Built-in Sorts
- Quantifiers
- Custom Sorts
- Uninterpreted Functions

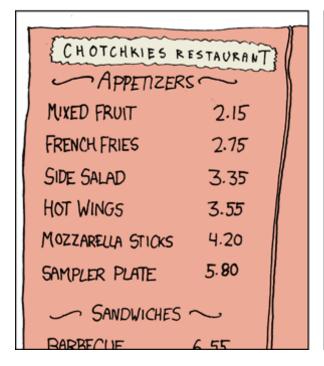
Working with Integers

- IntSort
 - <,>, <=, ==, +, -, etc.</p>

Working with Integers

- IntSort
 - <,>, <=, ==, +, -, etc.</p>

MY HOBBY: EMBEDDING NP-COMPLETE PROBLEMS IN RESTAURANT ORDERS





Example

```
#!/usr/bin/python3
from z3 import *
a,b,c,d,e,f = Ints('a b c d e f')
s = Solver()
s.add(215*a + 275*b + 335*c + 355*d + 420*e + 580*f == 1505,
a>=0, b>=0, c>=0, d>=0, e>=0, f>=0)
result = s.check()
if result == sat:
    print(s.model())
```

Variables in a Satisfying Model

- Variables and Expressions are stored in z3-specific classes
- We can use solver.model().decls() to iterate through all declared variables
 - Use .as_long() to convert a BitVector, Int, Real, etc. to a Python Integer

Example contd.

```
results=[]
while True:
    if s.check() == sat:
        m = s.model()
        print(m)
        results.append(m)
        block = [a != m[a].as long(), b != m[b].as long(), c != m[c].as long(), d !=
        m[d].as long(), e != m[e].as long(), f != m[f].as long()]
        11 11 11
        #Different approach: Iterate over all entries in the model
        block = []
        for d in m.decls():
            print(d, type(d), d(), type(d()), m[d], type(m[d]))
            c = d()
            block.append(c != m[d].as long())
        11 11 11
        s.add(Or(block))
    else:
        print ("All results enumerated, total=", len(results))
        break
```

Z3 Built-in Sorts

- BoolSort, BitVecSort, IntSort, RealSort
- Sequences, Strings
- Arrays

Assignment Sheet

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- Deadline: 05. 06. 2024 12 days

Assignment Sheet - Interviews

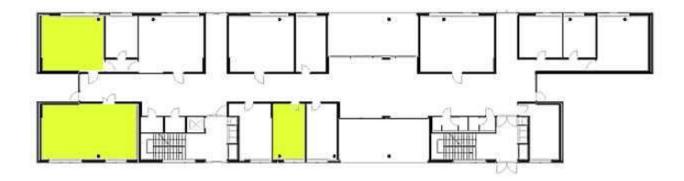
- 10 Min Interview Slots
- When? 06. 06. 12:00 17:00 Slots via TeachCenter

Assignment Sheet - Interviews

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- What? Discussion with (one of your) tutor or Stefan Bring your Laptop!

Assignment Sheet - Interviews

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- What? Discussion with (one of your) tutor or Stefan Bring your Laptop!
- Where? IAIK Ground Floor



Quantifiers

- Z3 offers ForAll() and Exists()
- Usage: ForAll (<vars>, <formula>)

Example

▶ 15. [M26] J. H. Quick noticed that $((x+2) \oplus 3) - 2 = ((x-2) \oplus 3) + 2$ for all x. Find all constants a and b such that $((x+a) \oplus b) - a = ((x-a) \oplus b) + a$ is an identity.

▶ 15. [M26] J. H. Quick noticed that $((x+2) \oplus 3) - 2 = ((x-2) \oplus 3) + 2$ for all x. Find all constants a and b such that $((x+a) \oplus b) - a = ((x-a) \oplus b) + a$ is an identity.

- We want to use Z3 to find all constants, s.t.
 - $\forall x ((x+a) \oplus b) a = ((x-a) \oplus b) + a$

• $\forall x ((x+a) \oplus b) - a = ((x-a) \oplus b) + a$

```
from z3 import *
s = Solver()
a, b = BitVecs('a b', 4)
x = BitVec('x', 4)
lhs, rhs = BitVecs('lhs rhs', 4)
s.add(lhs == ((x+a)^b)-a))
s.add(rhs == ((x-a)^b)+a))
s.push()
s.add(ForAll(x, lhs == rhs)
results=[]
while True:
  if s.check() == sat:
    m = s.model(); results.append(m)
    block = [a != m[a].as long(), b != m[b].as long()]
    s.add(Or(block))
  else:
    print ("results total=", len(results))
    break
```

- $\forall x ((x+a) \oplus b) a = ((x-a) \oplus b) + a$
 - Let's also use Z3 to find constants such that the equality does not hold

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 - Let's also use Z3 to find constants such that the equality does not hold
- Use solver.push() and solver.pop() to store and restore solver states

- $\forall x ((x+a) \oplus b) a = ((x-a) \oplus b) + a$
 - Let's also use Z3 to find constants such that the equality does not hold

```
from z3 import *
s = Solver()
a, b = BitVecs('a b', 4)
    = BitVec('x', 4)
X
s.add(lhs == ((x+a)^b)-a))
s.add(rhs == ((x-a)^b)+a))
s.push()
s.add(ForAll(x, lhs == rhs)
s.pop()
s.add(Exists(x, lhs != rhs)
result = s.check()
print(result)
print(s.sexpr())
if result == sat:
    print(s.model())
```

Custom Sorts – Datatypes

- Beyond the built-in Sorts
- Datatypes allow us to define more complex data structures, e.g. lists, trees, etc.

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- Datatypes allow us to define more complex data structures, e.g. lists, trees, etc.
- Simple Case: Enum

```
ColoursDatatype = Datatype("Colour")
ColoursDatatype.declare("RED")
ColoursDatatype.declare("GREEN")
ColoursDatatype.declare("BLUE")
ColoursDatatype.declare("MAGENTA")
ColoursSort = ColoursDatatype.create()
x = Const("x", ColoursSort)
```

Uninterpreted Functions

- Generally, we have:
 - $f: A_0 \times ... \times A_n \rightarrow B$
 - f maps values from $A_0 \times ... \times A_n$ to B

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 - Z3 decides the output based on the constraint
 - f can be seen as a lookup-table

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 - $f: A_0 \times ... \times A_n \to B$
 - f maps values from $A_0 \times ... \times A_n$ to B
- Uninterpreted Functions have no know "structure"
 - Z3 decides the output based on the constraint
 - f can be seen as a lookup-table
- f = Function('f', IntSort(), IntSort())
 f = Function('f', IntSort(), ColoursSort)

1D – Colouring Problem

```
from itertools import combinations
from z3 import *
solver = Solver()
Colours = Datatype("Colours")
Colour = Colours.create()
     = Function('f', IntSort(), Colour)
variables = list()
for i in range (0,5):
    variables.append(Int(i))
    solver.add(0 <= variables[-1])</pre>
    solver.add(variables[-1] <= 5)</pre>
solver.add(Distinct(variables))
for combi in combinations (variables, 2):
    solver.add(Implies(Abs(combi[0] - combi[1]) == 1, f(combi[0]) != f(combi[1])))
result = solver.check()
if result == sat:
    print(solver.model())
```

Z3 Check Results

- So far we have only used solver.model()
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- Can we use insights from unsat results?

Z3 Check Results

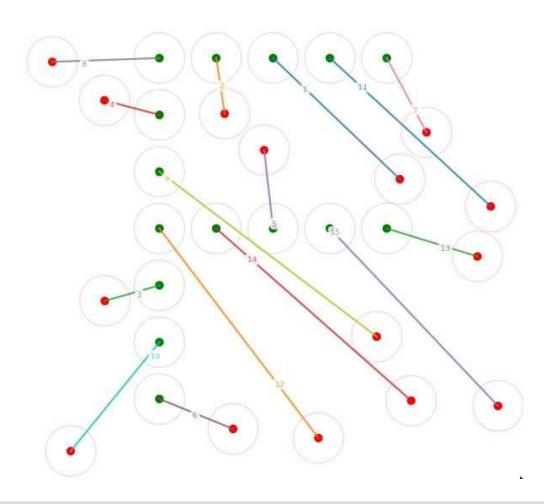
- So far we have only used solver.model()
 - Only available if assumptions are satisfiable
- Can we use insights from unsat results?
 - Yes! Z3 produces a so-called unsatisfiable core

Unsatisfiable Cores

- Can we use insights from unsat results?
 - Yes! Z3 produces a so-called unsatisfiable core
 - Tell Z3 to "track" certain assumptions:
 - solver.assert_and_track(<phi>, <Bool|String>)

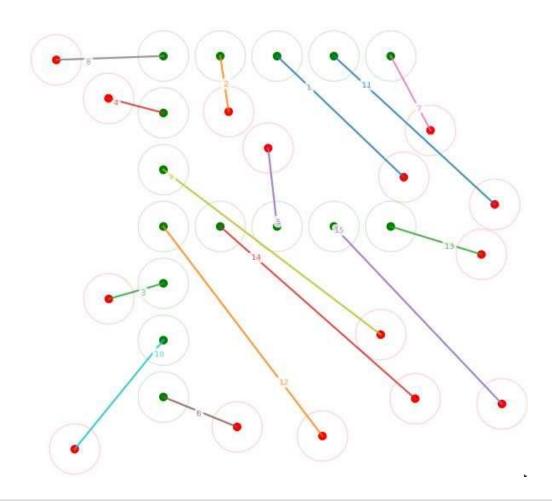
```
p, q = Bools('p q')
s = Solver()
s.add(Not(q))
s.assert_and_track(q, 'p1')
s.assert_and_track(p, 'p2')
print(s.check())
print(s.unsat core())
```

We are given a list of start positions s i and goal positions g i



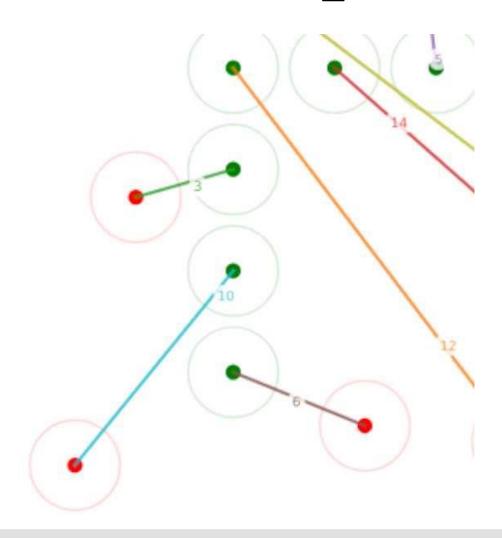
We are given a list of start positions s i and goal positions g i

- In each timestep j we want to move one molecule,
- without them crossing any other molecule while moving!



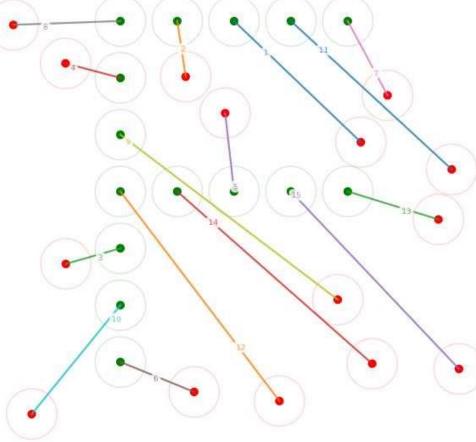
■ We are given a list of start positions s_i and goal positions g_i

- In each timestep j we want to move one molecule,
- without them crossing any other molecule while moving!
- E.g.: i=10 needs to move before 6



■ We are given a list of start positions s_i and goal positions g_i

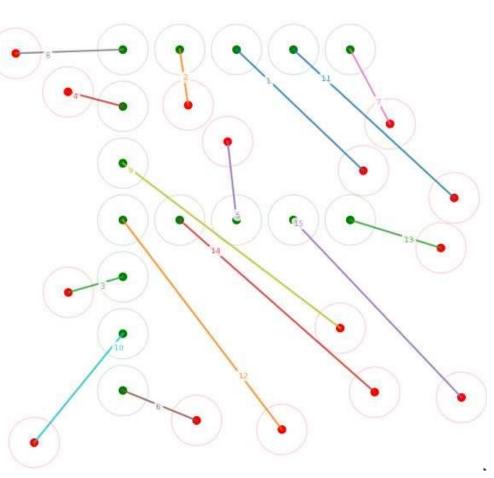
Which molecule i should move in which timestep j?



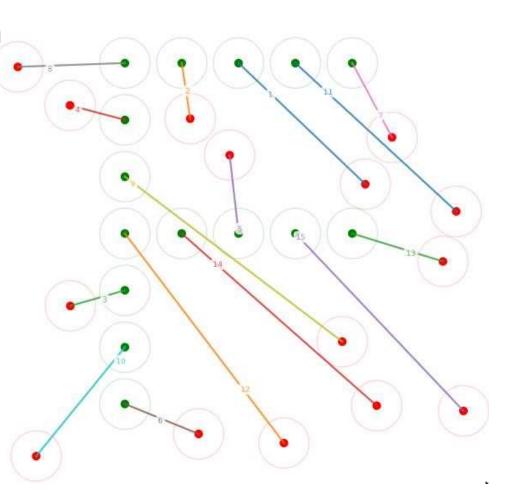
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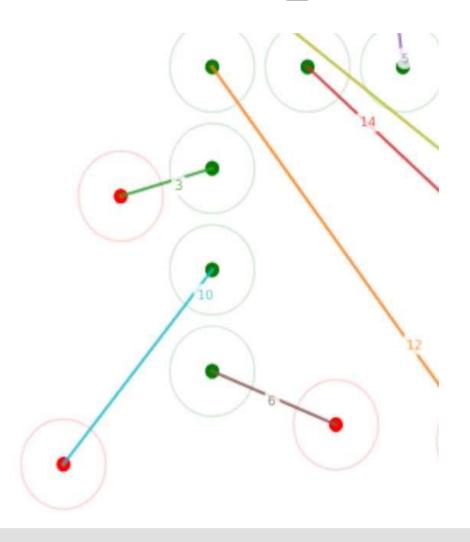
• What do we need to tell Z3 to give us a solution?



- We are given a list of start positions s_i and goal positions g_i
- Which molecule i should move in which timestep j?
- What do we need to tell Z3 to give us a solution?
- Booleans:
 - "Is molecule i at the start",
 - "is molecule i at the goal"
 - "Is molecule i scheduled to move at timestep j"



- We are given a list of start positions s_i and goal positions g_i
- Which molecule i should move in which timestep j?
- What do we need to tell Z3 to give us a solution?
- Constraints:
 - "If 10 moves at timestep j, then 6 has to move in a timestep > j"
 - "If a moves at timestep j, then b has to have moved in a timestep < j"</p>



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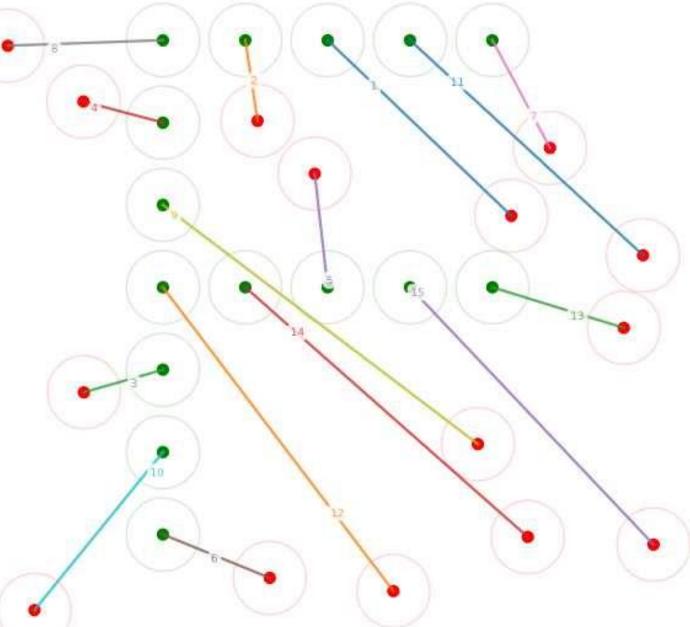
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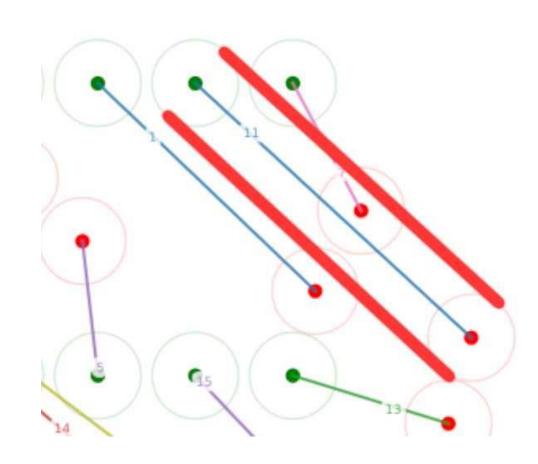
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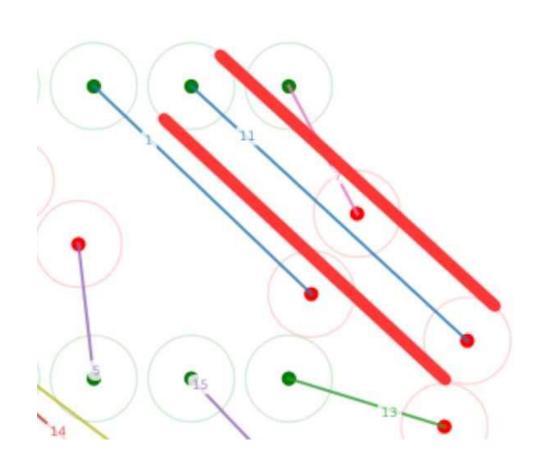
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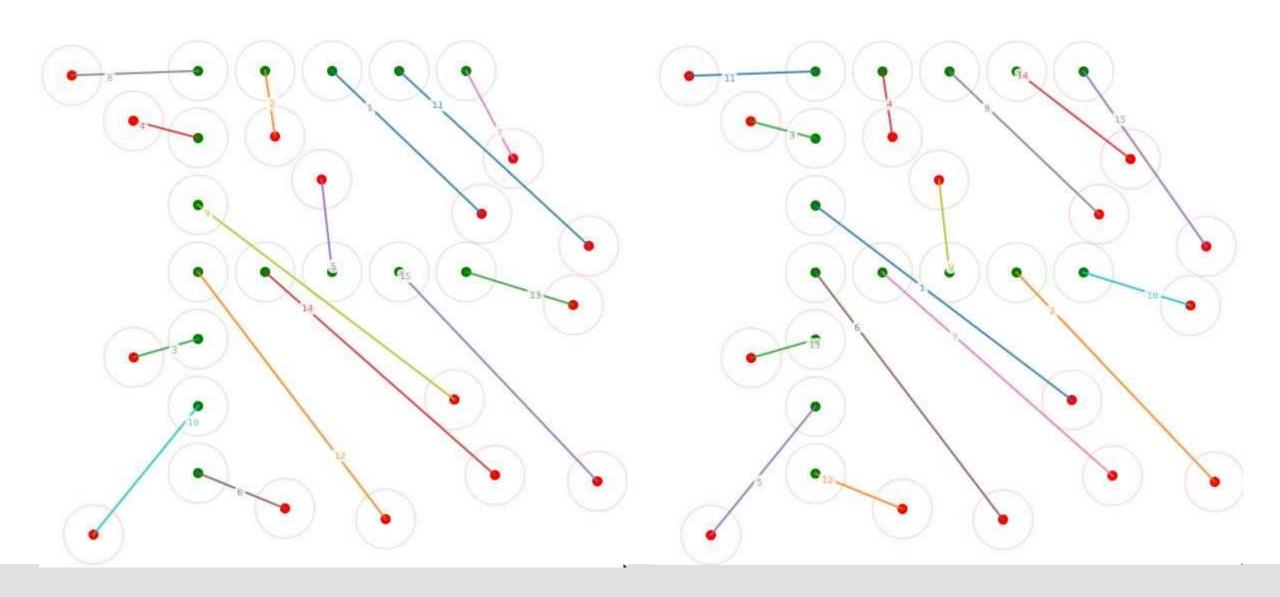


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- Which molecule i should move in which timestep j?
- What do we need to tell Z3 to give us a solution?
- Unsatisfiable:
 - 7 will never move out of the way for 11
 - [g-i1:7;i2:11, s-i1:7;i2:11]



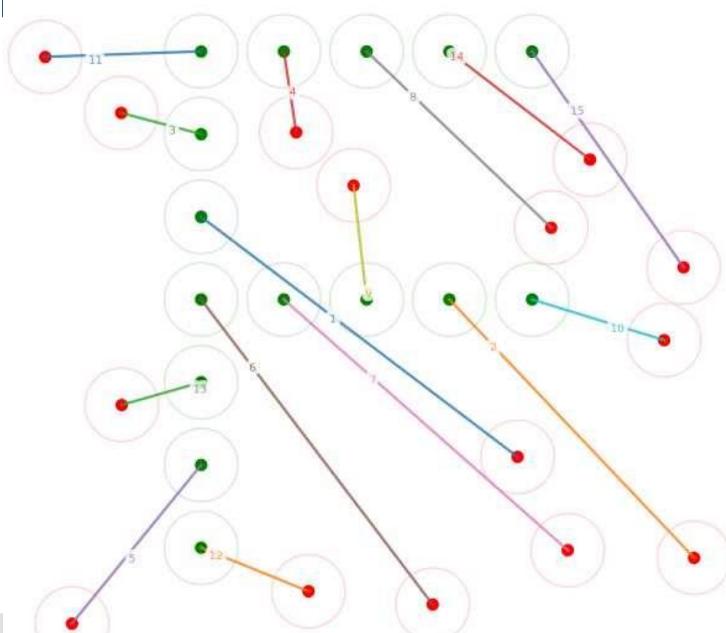
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- Which molecule i should move in which timestep j?
- What do we need to tell Z3 to give us a solution?
- Unsatisfiable:
 - 7 will never move out of the way for 11
 - [g-i1:7;i2:11, s-i1:7;i2:11]
 - Therefore: 7 and 11 need to switch goal positions!





Excu

vement



More Z3...

- solver = Solver()
 - A general, multi-purpose solver
- Z3 contains multiple proof methods that can be fine-tuned for a problem setting

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- solver = Solver()
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- Z3 contains multiple proof methods that can be fine-tuned for a problem setting
- Tactics and Tacticals:
 - A Tactic (<String>) describes a heuristic or algorithm to be applied
 - E.g.: ackermannize-bv, euf-completion (congruence-closure), simplify, solve-eqs, smt, sat, ...
 - A Tactical describes the application order of the different tactics:
 - Then ([<String>, ...]) apply tactics consecutively,
 - OrElse(t, s) apply t, if it fails apply s,
 - Repeat (t) repeat t until no changes happen
 - .solver() in order to get the according Solver

More Z3...

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