

Perfect matchings and Hamiltonian cycles in the preferential attachment model

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Abstract

In this paper, we study the existence of perfect matchings and Hamiltonian cycles in the preferential attachment model. In this model, vertices are added to the graph one by one, and each time a new vertex is created it establishes a connection with m random vertices selected with probabilities proportional to their current degrees. (Constant m is the only parameter of the model.) We prove that if $m \geq 1,260$, then asymptotically almost surely there exists a perfect matching. Moreover, we show that there exists a Hamiltonian cycle asymptotically almost surely, provided that $m \geq 29,500$. One difficulty in the analysis comes from the fact that vertices establish connections only with vertices that are “older” (i.e. are created earlier in the process). However, the main obstacle arises from the fact that edges in the preferential attachment model are not generated independently. In view of that, we also consider a simpler setting—sometimes called uniform attachment—in which vertices are added one by one and each vertex connects to m older vertices selected uniformly at random and independently of all other choices. We first investigate the existence of perfect matchings and Hamiltonian cycles in the uniform attachment model, and then extend the argument to the preferential attachment version.

1 Introduction

Two of the most natural questions concerning models of sparse random graph are whether or not a random instance is Hamiltonian or has a perfect matching. The hamiltonicity threshold for the basic models of *random graphs* $G(n, m)$ and $G(n, p)$ was established quite precisely by Komlós and Szemerédi [24]. (See Bollobás [7], Ajtai, Komlós and Szemerédi [1] and Bollobás and Frieze [8] for refinements.) Perfect matchings were first investigated (using Tutte’s theorem) by Erdős and Rényi [18]. (A simpler argument using expansion properties was provided by Bollobás and Frieze [8].)

The existence of perfect matchings and Hamiltonian cycles has also been investigated for other relevant models of random graphs. For instance, Díaz, Mitsche and Pérez-Giménez [14] obtained the hamiltonicity threshold for the *random geometric graph* $G_{n,r}$, which is defined as follows: the vertices of $G_{n,r}$ are n points chosen independently and uniformly at random from the unit square,

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and any two vertices are joined by an edge if and only if they are within distance r from each other in the square. The result in [14] was further strengthened by Balogh, Bollobás, Krivelevich, Müller and Walters [2] and Müller, Pérez-Giménez and Wormald [29], who gave a more precise characterization of the emergence of the first Hamiltonian cycle in the random geometric graph. (The analogous result for perfect matchings follows from a more general statement in [29].)

Another well-studied sparse random graph is the *random regular graph*. Let $G_{n,d}$ denote a graph chosen uniformly at random from the set of d -regular graphs with vertex set $[n]$. Robinson and Wormald [34] showed that *asymptotically almost surely* (a.a.s.) $G_{n,d}$ is Hamiltonian for constant $d \geq 3$. (Here and in similar statements, an event occurs a.a.s. if it occurs with probability tending to 1 as n tends to infinity.) Allowing d to grow with n presented some challenges, but they have now been resolved (see Cooper, Frieze and Reed [13] and Krivelevich, Sudakov, Vu and Wormald [27]).

Yet another random graph model extensively studied from that perspective is the so-called *m -out graph*. This time, each vertex $v \in [n]$ independently chooses m random out-neighbours to create the random digraph $D_{m\text{-out}}^n$. We then obtain $G_{m\text{-out}}^n$ by ignoring orientations. The hamiltonicity of $G_{m\text{-out}}^n$ was first discussed by Fenner and Frieze [19]. They showed that $G_{23\text{-out}}^n$ is a.a.s. Hamiltonian. This was improved to $G_{10\text{-out}}^n$ by Frieze [20] and to $G_{5\text{-out}}^n$ by Frieze and Luczak [22]. Cooper and Frieze [12] showed that $G_{4\text{-out}}^n$ is a.a.s. Hamiltonian, and the last drop was squeezed by Bohman and Frieze who established that $G_{3\text{-out}}^n$ is a.a.s. Hamiltonian [6].

In this paper, we study the *preferential attachment model*, which, arguably, is the best-known model for complex networks. The first consideration of this model goes back to 1925 when Yule used it to explain the power-law distribution of the number of species per genus of flowering plants [36]. The application of the model to describe the growth of the World Wide Web was proposed by Barabási and Albert in 1999 [3]. We will use the following precise definition of the model, as considered by Bollobás and Riordan in [9] as well as Bollobás, Riordan, Spencer, and Tusnády [10].

Let G_1^1 be the graph with one vertex, 1, and one loop. The random graph process $(G_1^t)_{t \geq 1}$ is defined inductively as follows. Given G_1^{t-1} , we form G_1^t by adding vertex t together with a single edge between t and i , where i is selected randomly with the following probability distribution:

$$\Pr(i = s) = \begin{cases} \deg(s, t-1)/(2t-1) & 1 \leq s \leq t-1, \\ 1/(2t-1) & s = t, \end{cases}$$

where $\deg(s, t-1)$ denotes the degree of vertex s in G_1^{t-1} (i.e. the degree of s “at time $t-1$ ”, after vertex $t-1$ was added). In other words, we send an edge e from vertex t to a random vertex i , where the probability that a vertex is chosen as i is proportional to its degree at the time, counting e as already contributing one to the degree of t . In particular, each vertex may only be attached to itself or to an “older” vertex (i.e. a vertex created earlier in the process).

For $m \in \mathbb{N} \setminus \{1\}$, the process $(G_m^t)_{t \geq 1}$ is defined similarly with the only difference that m edges are added to G_m^{t-1} to form G_m^t (one at a time), counting previous edges as already contributing to the degree distribution. Equivalently, one can define the process $(G_m^t)_{t \geq 1}$ by considering the process $(G_1^t)_{t \geq 1}$ on a sequence $1', 2', \dots$ of vertices; the graph G_m^t is formed from G_1^{tm} by identifying vertices $1', 2', \dots, m'$ to form 1, identifying vertices $(m+1)', (m+2)', \dots, (2m)'$ to form 2, and so on. Although $(G_m^t)_{t \geq 1}$ is an infinite-time random process, we will restrict our attention to a finite number of time steps and consider $(G_m^t)_{1 \leq t \leq n}$ or G_m^n for large n . Note that in this model G_m^n is in general a multigraph, possibly with multiple edges between two vertices (if $m \geq 2$) and self-loops. For our purpose, as we are interested in Hamiltonian cycles and perfect matchings, loops can be ignored and multiple edges between two nodes can be treated as a single edge.

Perhaps the most remarkable feature of the preferential attachment mechanism is that it provides a simple illustration of the *rich-get-richer* principle, by which vertices that acquire higher

degree early in the process use this accumulated advantage to attract more edges during the process. This results in a heavy-tailed degree sequence with some vertices of very high degree, which do not typically occur in the previous models of random graphs described above. Indeed, it was shown in [10] that, for any constant $m \in \mathbb{N}$, a.a.s. the degree distribution of G_m^n follows a power law: the number of vertices with degree k falls off as $(1 + o(1))ck^{-3}n$ for some explicit constant $c = c(m)$ and large $k \leq n^{1/15}$.

For the purpose of this paper, the case $m = 1$ is easy to analyze, since G_1^n is a forest. Each node sends an edge either to itself or to an earlier node, so the graph consists of components which are trees, each with a loop attached. The expected number of components is then $\sum_{t=1}^n 1/(2t-1) \sim (1/2) \log n$ and, since events are independent, we derive that a.a.s. there are $(1/2 + o(1)) \log n$ components in G_1^n by Chernoff's bound. In particular, G_1^n is a.a.s. disconnected and thus contains no Hamiltonian cycle. A similar argument shows that a.a.s. many components of G_1^n have an odd number of vertices, so G_1^n has no perfect matching. In view of this, we will restrict our attention throughout the paper to the case $m \geq 2$, for which it is known [9] that G_m^n is a.a.s. connected.

We finally consider the *uniform attachment* graph $G_{m\text{-old}}^n$ on vertex set $[n]$, which can be thought of as an intermediate model between the m -out and preferential attachment models. In this setting, vertex 1 has m directed loops. Each vertex $v \in [n] \setminus \{1\}$ independently chooses m random out-neighbours (with repetition) from $[v-1]$, to create the random digraph $D_{m\text{-old}}^n$. We then obtain $G_{m\text{-old}}^n$ by ignoring orientations. (Note the sole purpose of having loops in the definition is to ensure that the number of edges of $G_{m\text{-old}}^n$ is precisely mn , but loops and multiple edges play no role in the existence of perfect matchings or Hamiltonian cycles and thus may be ignored.) We may also regard this model as a random process $(G_{m\text{-old}}^t)_{1 \leq t \leq n}$, in which at each step a new vertex is created and attached only to older vertices (as in preferential attachment), but the choice of these vertices is uniform (as in the m -out model). It turns out that the fact that edges are always generated towards older vertices causes one of the main difficulties in our analysis, which motivates a unified treatment for uniform and preferential attachment. In fact, we will often examine the uniform attachment case first for simplicity, and then adapt the argument to the non-uniform (rich-get-richer) distribution in preferential attachment. However, a more challenging and technical problem we encountered in transferring the result to not-uniform model was the fact that edges generated in the preferential attachment model are not generated independently. The two-round exposure technique (explained in Subsection 2.3), used in the argument, has to be dealt with care. The lemmas needed to overcome all difficulties might be interesting on their own rights, and could be useful in future applications of this powerful technique.

In this paper, we prove that there exists a constant $c \in \mathbb{N}$ such that a.a.s. both G_m^n and $G_{m\text{-old}}^n$ have a Hamiltonian cycle, provided that $m \geq c$. Of course, the existence of a Hamiltonian cycle implies the existence of a perfect matching¹. However, we treat both properties independently in order to obtain smaller constants for the weaker property. Specifically, for perfect matchings we get $c = 159$ for $G_{m\text{-old}}^n$ (Theorem 13) and $c = 1,260$ for G_m^n (Corollary 14). For Hamiltonian cycles we get $c = 3,214$ for $G_{m\text{-old}}^n$ (Theorem 16) and $c = 29,500$ for G_m^n (Corollary 17). We tried to tune parameters in the argument to get the corresponding c 's as small as possible but clearly, with more effort, one may improve them further. However, it seems that with the existing argument it is impossible to find the threshold c for any of the two properties and any of the two models. In particular, is it true that G_3^n or $G_{3\text{-old}}^n$ are a.a.s. Hamiltonian? This remains an open problem for now. On the other hand, all we managed to show is that a.a.s. G_2^n has no perfect matching (and so also no Hamiltonian cycle) and that $G_{2\text{-old}}^n$ has no Hamiltonian cycle (see discussion in Section 4).

¹In the sequel, a perfect matching will denote a matching of size $\lfloor n/2 \rfloor$ (thus allowing one unmatched vertex if n is odd).

2 Preliminaries

2.1 Basic definitions and notation

In order to state some of the intermediate results in the paper, it is convenient to extend the definitions of $G_{m\text{-old}}^n$ and G_m^n given in Section 1 to the case $m = 0$ by interpreting both $G_{0\text{-old}}^n$ and G_0^n as the graph on vertex set $[n]$ and no edges. Fix any constant $m \in \mathbb{Z}_{\geq 0}$. For any vertices $v, w \in [n]$ satisfying $w < v$, we say that w is *older* than v (or v is *younger* than w), since w is created earlier than v in the process leading to $G_{m\text{-old}}^n$ or G_m^n . We say that an edge $e = vw$ *stems* from v if $w \leq v$ (or equivalently e is one of the m edges that attach v to younger vertices or to v itself at the step of the process when v is added to the graph).

Given a graph $G = (V, E)$ and a set $C \subseteq V$, we will use $N(C)$ to denote the *neighbourhood* of C ; that is,

$$N(C) = \{w \in V \setminus C : vw \in E \text{ for some } v \in C\},$$

and we write $N(v) = N(\{v\})$ for simplicity. All asymptotics throughout are as $n \rightarrow \infty$ (we emphasize that the notations $o(\cdot)$ and $O(\cdot)$ refer to functions of n , not necessarily positive, whose growth is bounded). We also use the notations $f \ll g$ for $f = o(g)$ and $f \gg g$ for $g = o(f)$. For simplicity, we will write $f(n) \sim g(n)$ if $f(n)/g(n) \rightarrow 1$ as $n \rightarrow \infty$ (that is, when $f(n) = (1 + o(1))g(n)$). Since we aim for results that hold a.a.s. (see the definition in the previous section), we will always assume that n is large enough. We often write, for example, G_m^n when we mean a graph drawn from the distribution G_m^n . Given $k \in \mathbb{Z}_{\geq 0}$, we define $k!! = k(k-2)(k-4) \cdots 3 \cdot 1$ for odd k and $k!! = k(k-2)(k-4) \cdots 4 \cdot 2$ for even k . Finally, for any $x \in \mathbb{R}_+$, we use $\log x$ to denote the natural logarithm of x , and $[k] = \{1, 2, \dots, k\}$ for any $k \in \mathbb{N}$.

2.2 Useful bounds

Most of the time, we will use the following version of *Chernoff's bound*. Suppose that $X \in \text{Bin}(n, p)$ is a binomial random variable with expectation $\mu = np$. If $0 < \delta < 1$, then

$$\Pr[X < (1 - \delta)\mu] \leq \exp\left(-\frac{\delta^2\mu}{2}\right),$$

and if $\delta > 0$,

$$\Pr[X > (1 + \delta)\mu] \leq \exp\left(-\frac{\delta^2\mu}{2 + \delta}\right).$$

However, at some point we will need the following, stronger, version: for any $t \geq 0$, we have

$$\Pr[X \geq \mu + t] \leq \exp\left(-\mu\varphi\left(\frac{t}{\mu}\right)\right), \tag{1}$$

and for any $0 \leq t \leq \mu$, we have

$$\Pr[X \leq \mu - t] \leq \exp\left(-\mu\varphi\left(\frac{-t}{\mu}\right)\right), \tag{2}$$

where

$$\varphi(x) = (1 + x)\log(1 + x) - x. \tag{3}$$

Moreover, let us mention that all of these bounds hold for the general case in which X is a sum of Bernoulli(p_i) indicator random variables with (possibly) different p_i . These inequalities are well known and can be found, for example, in [23].

We will also use a standard martingale tool: the *Hoeffding-Azuma inequality*. As before, see for example [23] for more details. Let X_0, X_1, \dots be a martingale. Suppose that there exist $c_1, c_2, \dots, c_n > 0$ such that $|X_k - X_{k-1}| \leq c_k$ for each $1 \leq k \leq n$. Then, for every $x > 0$,

$$\Pr[X_n > \mathbf{E}X_n + x] \leq \exp\left(-\frac{x^2}{2\sum_{k=1}^n c_k^2}\right). \quad (4)$$

The Hoeffding-Azuma inequality can be generalized to include random variables close to martingales. One of our proofs, proof of Lemma 5, will use the supermartingale method of Pittel et al. [33], as described in [35, Corollary 4.1]. Let X_0, X_2, \dots, X_n be a sequence of random variable. Suppose that there exist $c_1, c_2, \dots, c_n > 0$ and $b_1, b_2, \dots, b_n > 0$ such that

$$|X_k - X_{k-1}| \leq c_k \quad \text{and} \quad \mathbf{E}[X_k - X_{k-1} | X_{k-1}] \leq b_k$$

for each $1 \leq k \leq n$. Then, for every $x > 0$,

$$\Pr\left[\text{For some } t \text{ with } 0 \leq t \leq n: X_t - X_0 > \sum_{k=1}^t b_k + x\right] \leq \exp\left(-\frac{x^2}{2\sum_{k=1}^n c_k^2}\right). \quad (5)$$

Finally, we include the following auxiliary lemma that will simplify some calculations in Lemma 4. (We use the convention that the empty product is equal to 1 and the empty sum is equal to 0.)

Lemma 1. *Given any integers $a = a(n)$ and $b = b(n)$ such that $0 \leq a \leq b$, define*

$$c_{a,b} = \prod_{i=a+1}^b \left(\frac{2i-1}{2i}\right).$$

Then,

$$c_{a,b} = e^{O(\varepsilon)} \sqrt{a/b},$$

where $\varepsilon = 1/(a+1)$, and the constants involved in the $O()$ notation do not depend on a or b . Moreover

$$\sum_{i=a+1}^b c_{a,i}^2 = e^{O(\varepsilon)} a \log(b/a).$$

Proof. The case $a = b$ is trivial, so we may assume $a < b$. Then,

$$\begin{aligned} c_{a,b} &= \prod_{i=a+1}^b \left(1 - \frac{1}{2i}\right) = \prod_{i=a+1}^b \exp\left(-\frac{1}{2i} + O(i^{-2})\right) = \exp\left(-\sum_{i=a+1}^b \frac{1}{2i} + O(\varepsilon)\right) \\ &= \exp\left(-\frac{1}{2} \log(b/a) + O(\varepsilon)\right) = e^{O(\varepsilon)} \sqrt{a/b}, \end{aligned}$$

where we used the fact that the harmonic sum satisfies $\sum_{i=1}^n 1/i = \log n + \gamma + O(1/n)$. This proves the first claim of the lemma. The following calculation yields the second claim:

$$\sum_{i=a+1}^b c_{a,i}^2 = e^{O(\varepsilon)} a \sum_{i=a+1}^b 1/i = e^{O(\varepsilon)} a \log(b/a). \quad \square$$

2.3 Two-round exposure

Fix constants $m, m_1, m_2 \in \mathbb{Z}_{\geq 0}$ such that $m = m_1 + m_2$, and recall the definitions in Section 2.1. For each vertex $v \in [n]$ of either model $G_{m\text{-old}}^n$ or G_m^n , we colour blue the first m_1 edges stemming from v and colour red the m_2 remaining ones. We denote the resulting edge-coloured versions of $G_{m\text{-old}}^n$ and G_m^n by $G_{m_1, m_2\text{-old}}^n$ and G_{m_1, m_2}^n , respectively. (Note that these edge colourings are not proper in the usual graph-theoretical sense.) Let $\pi_1(G_{m_1, m_2\text{-old}}^n)$ be the graph on vertex set $[n]$ whose edges are precisely the blue edges of $G_{m_1, m_2\text{-old}}^n$, and symmetrically let $\pi_2(G_{m_1, m_2\text{-old}}^n)$ be the graph with the red edges of $G_{m_1, m_2\text{-old}}^n$. Moreover, let $\pi(G_{m_1, m_2\text{-old}}^n)$ be the union of graphs $\pi_1(G_{m_1, m_2\text{-old}}^n)$ and $\pi_2(G_{m_1, m_2\text{-old}}^n)$ — or simply $G_{m_1, m_2\text{-old}}^n$ — after removing the colours from the edges. By construction, $\pi(G_{m_1, m_2\text{-old}}^n)$ is distributed as $G_{m\text{-old}}^n$. Finally, graphs $\pi_1(G_{m_1, m_2}^n)$, $\pi_2(G_{m_1, m_2}^n)$ and $\pi(G_{m_1, m_2}^n)$ are defined analogously by replacing $G_{m_1, m_2\text{-old}}^n$ by G_{m_1, m_2}^n .

Observe that, for each $\sigma \in \{1, 2\}$, $\pi_\sigma(G_{m_1, m_2\text{-old}}^n)$ has the same distribution as $G_{m_\sigma\text{-old}}^n$ (so we can identify $\pi_\sigma(G_{m_1, m_2\text{-old}}^n)$ and $G_{m_\sigma\text{-old}}^n$), and moreover the graphs $\pi_1(G_{m_1, m_2\text{-old}}^n)$ and $\pi_2(G_{m_1, m_2\text{-old}}^n)$ are independent of each other. In other words, the union of two independent uniform attachment graphs $G_{m_1\text{-old}}^n$ and $G_{m_2\text{-old}}^n$ is distributed as $G_{m\text{-old}}^n$ with $m = m_1 + m_2$. This allows us to build $G_{m\text{-old}}^n$ by first exposing the edges of $G_{m_1\text{-old}}^n$ (blue edges) and then adding the edges of $G_{m_2\text{-old}}^n$ (red edges). This technique is known as the two-round exposure method, and will be repeatedly used in Section 3 to build a perfect matching and a Hamiltonian cycle in $G_{m\text{-old}}^n$.

Unfortunately, the analogous property does not hold for the preferential attachment model (unless² m_1 or m_2 equals 0). That is, if $m_1, m_2 > 0$, then $G_{m\text{-old}}^n$ is not distributed as the union of independent $G_{m_1\text{-old}}^n$ and $G_{m_2\text{-old}}^n$. This is due to the fact that, when a new edge stemming from a vertex v is created, the choice of the second endpoint depends on all the edges (both blue and red) previously added during the process. We want to use a version of the two-round exposure technique, and build the preferential attachment graph G_m^n as the union of $\pi_1(G_{m_1, m_2}^n)$ and $\pi_2(G_{m_1, m_2}^n)$. However (again if $m_1, m_2 > 0$), for each $\sigma \in \{1, 2\}$, graph $\pi_\sigma(G_{m_1, m_2}^n)$ is not distributed as $G_{m_\sigma}^n$, and moreover $\pi_1(G_{m_1, m_2}^n)$ and $\pi_2(G_{m_1, m_2}^n)$ are not independent of each other. In spite of all these obstacles, our argument will analyze $\pi_1(G_{m_1, m_2}^n)$ and $\pi_2(G_{m_1, m_2}^n)$ separately, in a similar fashion as we deal with $G_{m_1\text{-old}}^n$ and $G_{m_2\text{-old}}^n$ for the uniform attachment model. This time though, we will need to use properties of $\pi_1(G_{m_1, m_2}^n)$ that hold a.a.s. conditional on a given $\pi_2(G_{m_1, m_2}^n)$ and vice-versa.

2.4 Relationship between G_m^n and $G_{m\text{-old}}^n$

In this section, we will analyze some basic features of $G_{m\text{-old}}^n$ and $\pi_\sigma(G_{m_1, m_2}^n)$, where $\sigma \in \{1, 2\}$ and m, m_1, m_2 are fixed nonnegative integers such that $m = m_1 + m_2$. As already mentioned, if $m_2 = 0$ then $\pi_1(G_{m_1, m_2}^n)$ is simply the preferential attachment model G_m^n . Hence, various lemmas stated below for G_{m_1, m_2}^n also apply to G_m^n .

Given any vertex $v \in [n]$ and any set $W \subseteq [v-1]$ of vertices older than v , the probability in $G_{m\text{-old}}^n$ that there is no edge between v and W is precisely

$$\left(1 - \frac{|W|}{v-1}\right)^m. \quad (6)$$

Moreover, the corresponding event is independent from all edges added during the construction of $G_{m\text{-old}}^n$ stemming from vertices different from v . This is a crucial fact in our argument, which

²If $m_2 = 0$ then $\pi_2(G_{m_1, m_2}^n)$ has no edges, so $\pi_1(G_{m_1, m_2}^n) = G_m^n = G_{m_1}^n$, and vice-versa.

unfortunately does not extend to the preferential attachment model G_m^n , since the attachments of vertex v are not independent from “the future” (i.e. they are not independent from the edges that stem from vertices younger than v , which are added later in the process). Fortunately, we can obtain a weaker statement that will suffice for our purposes, and may be useful for other applications.

Lemma 2. *Fix any constants $m, m_1, m_2 \in \mathbb{Z}_{\geq 0}$ with $m = m_1 + m_2$, and let $\sigma \in \{1, 2\}$. Given any vertex $v \in [n]$ and any set $W \subseteq [v - 1]$, let E be any event of G_{m_1, m_2}^n that does not involve any edges stemming from v in $\pi_\sigma(G_{m_1, m_2}^n)$. Then, conditional upon E , the probability that $\pi_\sigma(G_{m_1, m_2}^n)$ has no edges between v and W is at most*

$$\left(1 - \frac{|W|}{v+n}\right)^{m_\sigma} \leq \left(1 - \frac{|W|}{2n}\right)^{m_\sigma}.$$

Proof. The statement is trivially true if $m_\sigma = 0$, so we assume that $m_\sigma \geq 1$.

Let e_1, \dots, e_{m_σ} be the m_σ edges that stem from v in $\pi_\sigma(G_{m_1, m_2}^n)$. Recall that these edges are blue if $\sigma = 1$, and red if $\sigma = 2$. For each $i \in [m_\sigma]$ and each $u \in [v]$, let $B_{i,u}$ be the event that e_i is attached to vertex u , and let $B_i = \bigcup_{u \in W} B_{i,u}$ be the event that e_i is attached to some vertex in W . Also, let E'_i be any event that involves only edges (blue or red) of G_{m_1, m_2}^n that are different from e_i .

Claim 1. Fix $i \in [m_\sigma]$. The probability of B_i conditional upon E'_i is at least $|W|/(v+n)$.

The lemma follows immediately by repeatedly applying Claim 1 to each B_i ($i \in [m_\sigma]$) with

$$E'_i = E \cap \overline{B_1} \cap \overline{B_2} \cap \dots \cap \overline{B_{i-1}}.$$

We proceed to prove Claim 1. Since $i \in [m_\sigma]$ is fixed, we will omit for simplicity the subindex i from the notation. The colours of the edges play no role in this claim, so we may ignore them and regard G_{m_1, m_2}^n simply as G_m^n with $m = m_1 + m_2$. Also, we can assume that event E' fully determines all $mn - 1$ edges of the process $(G_m^t)_{1 \leq t \leq n}$ other than e . (Indeed, any other situation can be trivially reduced to this one by using the law of total probabilities.) Moreover, we may assume that E' not only determines edges that are introduced during the process but also an order in which they were added.

For each $u \in [n] \setminus \{v\}$, let d_u be the degree of u in $G_m^n - e$ conditional upon E' (each self-loop of u contributes with 2 to this degree). Define d_v analogously but adding an extra 1 to account for the contribution of the first endpoint of e . It is easy to observe from the definition of the process $(G_m^t)_{1 \leq t \leq n}$ that, given any $u \in [v - 1]$,

$$\Pr(B_u \cap E') = \frac{d_u \prod_{w \in [n]} f_w}{(mn - 1)!!}, \quad (7)$$

where each factor f_w depends only on E' . For instance, if vertex $w \in [n]$ has no self-loops in $G_m^n - e$ given E' then $f_w = (d_w - 1)!/(m - 1)!$. Similar expressions for f_w can be easily computed in the case that w has some self-loops (the value of f_w depends on which edges stemming from w create a loop). The case $u = v$ is slightly different, but we have

$$\Pr(B_v \cap E') \leq \frac{d_v \prod_{w \in [n]} f_w}{(mn - 1)!!}. \quad (8)$$

In view of (7) and (8), for every $u \in [v - 1]$,

$$\Pr(B_u | E') = \frac{\Pr(B_u \cap E')}{\sum_{w \in [v]} \Pr(B_w \cap E')} \geq \frac{d_u}{\sum_{w \in [v]} d_w} \geq \frac{d_u}{2mv + m(n - v)} = \frac{d_u}{m(v + n)}.$$

Therefore, summing over $u \in W$,

$$\Pr(B \mid E') \geq \sum_{u \in W} \frac{d_u}{m(v+n)} \geq \frac{m|W|}{m(v+n)} = \frac{|W|}{v+n},$$

as required. This yields Claim 1 and finishes the proof of the lemma. \square

Note that the bound in the statement of Lemma 2 could be much worse than (6), especially if v is much smaller than n . However, in those cases in which event E does not depend upon the “future” of vertex v (i.e. does not depend on edges stemming from younger vertices in the process), we can derive a much better bound which resembles that of (6).

Lemma 3. *Fix any constants $m, m_1, m_2 \in \mathbb{Z}_{\geq 0}$ with $m = m_1 + m_2$ and let $\sigma \in \{1, 2\}$. Given any vertex $v \in [n]$ and any set $W \subseteq [v-1]$, let E be any event that involves only edges that stem from vertices in $[v-1]$ (i.e. older than v) in the process leading to G_{m_1, m_2}^n . Then, conditional upon E , the probability that $\pi_\sigma(G_{m_1, m_2}^n)$ has no edges between v and $W \subseteq [v-1]$ is at most*

$$\left(1 - \frac{|W|}{2v}\right)^{m_\sigma}.$$

Proof. Suppose that all (blue and red) edges of G_{m_1, m_2}^n stemming from vertices in $[v-1]$ have already been exposed. Additionally, suppose that some edges (or perhaps none) stemming from v have been exposed as well. We add a new edge e stemming from v . At that time, the total degree of W (i.e. the sum of the degrees of all its vertices) is at least $m|W|$, and the total degree of the graph (counting e as already contributing one to the degree of v) is at most $2mv$. Therefore, the conditional probability that e does not join v to a vertex in W is at most

$$1 - \frac{m|W|}{2mv} = 1 - \frac{|W|}{2v}.$$

The lemma follows immediately by considering the m_σ edges stemming from v in $\pi_\sigma(G_{m_1, m_2}^n)$. \square

In view of this, some of the proofs for $G_{m\text{-old}}^n$ in the paper that use (6) can be easily adapted to $\pi_\sigma(G_{m_1, m_2}^n)$ (and thus to G_m^n) by adjusting constants. Most of the time, Lemma 2 or Lemma 3 will be enough to do it. However, at some point, we will need an upper bound on the probability in $\pi_\sigma(G_{m_1, m_2}^n)$ that all the neighbours of v fall inside of W . In order to do that we will use Lemma 6 below but, before we can state it, we need the following lemma that bounds the total degree of a set W at the time vertex v is created in the process leading to G_m^n .

Lemma 4. *Let $A > 0$ be a sufficiently large constant. Fix any constant $m \in \mathbb{Z}_{\geq 0}$, and let $\omega = \omega(n)$ be any function tending to infinity as $n \rightarrow \infty$. The following property holds a.a.s. for G_m^n . For any $1 \leq k \leq t \leq n$ and any set $W \subseteq [t]$ of size k ,*

$$\sum_{w \in W} \deg(w, t) \leq \begin{cases} (1 + A/(mk + 1)) \left(2m\sqrt{kt} + \sqrt{8mkt} \log(et/k)\right) & \text{if } t \geq \omega, \\ 2m\omega & \text{otherwise.} \end{cases}$$

Proof. The statement is trivially true if $m = 0$ since G_0^n has no edges, so assume that $m \geq 1$. Moreover, the case $k = t$ (which corresponds to $W = [t]$) follows immediately from the case $W = [t-1]$ (at the only expense of replacing A by a slightly larger absolute constant), so we will ignore it.

Fix k and t such that $1 \leq k < t \leq n$, and assume $t \geq \omega \rightarrow \infty$ as $n \rightarrow \infty$. We will need to understand the behaviour of the following random variable:

$$Y_t(k) = \sum_{j=1}^k \deg(j, t).$$

(In other words, $Y_t(k)$ is the sum of the degrees of the k oldest vertices (at time t .) In view of the identification between the models G_m^n and G_1^{mn} , it will be useful to investigate the following collection of random variables: for $mk \leq s \leq mn$, let

$$X_s = \sum_{j'=1}^{mk} \deg_{G_1^t}(j', s).$$

Clearly, $Y_t(k) = X_{mt}$ and $X_{mk} = 2mk$. Moreover, for $mk < s \leq mn$,

$$X_s = \begin{cases} X_{s-1} + 1 & \text{with probability } \frac{X_{s-1}}{2s-1}, \\ X_{s-1} & \text{otherwise.} \end{cases}$$

The conditional expectation is given by

$$\mathbf{E}(X_s | X_{s-1}) = (X_{s-1} + 1) \frac{X_{s-1}}{2s-1} + X_{s-1} \left(1 - \frac{X_{s-1}}{2s-1}\right) = \frac{2s}{2s-1} X_{s-1}.$$

Define $\widehat{X}_{mk} = X_{mk} = 2mk$ and

$$\widehat{X}_s = \left(\prod_{i=mk+1}^s \frac{2i-1}{2i} \right) X_s \quad \text{for } mk < s \leq mn,$$

or equivalently $\widehat{X}_s = c_{mk,s} X_s$ with $c_{mk,s}$ defined as in Lemma 1. We have that $\mathbf{E}(\widehat{X}_s | \widehat{X}_{s-1}) = \widehat{X}_{s-1}$, and thus $\widehat{X}_{mk}, \dots, \widehat{X}_{mn}$ is a martingale with $\mathbf{E}\widehat{X}_s = 2mk$. Moreover, the difference between consecutive terms is

$$|\widehat{X}_s - \widehat{X}_{s-1}| = c_{mk,s} \left| X_s - \left(1 + \frac{1}{2s-1}\right) X_{s-1} \right| = c_{mk,s} \left| X_s - X_{s-1} - \frac{X_{s-1}}{2s-1} \right| \leq c_{mk,s},$$

since $X_s - X_{s-1} \in \{0, 1\}$ and $0 < \frac{X_{s-1}}{2s-1} \leq \frac{mk+s-1}{2s-1} < 1$. Also, in view of Lemma 1, there exist absolute constants $A_1, A_2 > 0$ such that

$$c_{mk,mt} \geq e^{-A_1/(mk+1)} \sqrt{k/t} \quad \text{and} \quad \sqrt{2 \sum_{i=mk+1}^{mt} c_{mk,i}^2} \leq e^{A_2/(mk+1)} \sqrt{2mk \log(t/k)}. \quad (9)$$

Assume that constant A in the statement is sufficiently large so that $1 + A/x \geq e^{(A_1+A_2)/x}$ for all $x \geq 1$ (in particular, $A = e^{A_1+A_2} - 1$ works), and define

$$\begin{aligned} D &:= (1 + A/(mk+1)) \left(2m\sqrt{kt} + \sqrt{8mkt} \log(et/k) \right) \\ &\geq e^{(A_1+A_2)/(mk+1)} \left(2m\sqrt{kt} + \sqrt{8mkt} \log(et/k) \right). \end{aligned} \quad (10)$$

We wish to bound the probability that $Y_t(k) > D$. Observe that this event implies that

$$\begin{aligned} \frac{\widehat{X}_{mt} - 2mk}{\sqrt{2 \sum_{i=mk+1}^{mt} c_{mk,i}^2}} &= \frac{c_{mk,mt} Y_t(k) - 2mk}{\sqrt{2 \sum_{i=mk+1}^{mt} c_{mk,i}^2}} > \frac{c_{mk,mt} D - 2mk}{\sqrt{2 \sum_{i=mk+1}^{mt} c_{mk,i}^2}} \geq \\ &\geq \frac{e^{A_2/(mk+1)} (2mk + \sqrt{8mk \log(et/k)}) - 2mk}{e^{A_2/(mk+1)} \sqrt{2mk \log(t/k)}} > \\ &> \frac{2\sqrt{k} \log(et/k)}{\sqrt{\log(t/k)}} > 2\sqrt{\log((et/k)^k)} \geq \sqrt{\log\left(\binom{t}{k}^4\right)} \geq \sqrt{\log\left(t^3 \binom{t}{k}\right)}, \end{aligned}$$

where we used (9), (10) and the facts that $1 \leq k < t$ and $t \leq \binom{t}{k} \leq (et/k)^k$. Hence, applying Hoeffding-Azuma inequality (4) to the martingale $\widehat{X}_{mk}, \dots, \widehat{X}_{mn}$ we get that

$$\Pr(Y_t(k) > D) \leq \Pr\left(\frac{\widehat{X}_{mt} - 2mk}{\sqrt{2 \sum_{i=mk+1}^{mt} c_{mk,i}^2}} > \sqrt{\log\left(t^3 \binom{t}{k}\right)}\right) \leq e^{-\log(t^3 \binom{t}{k})} = 1/t^3 \binom{t}{k}.$$

Observe that for any set of vertices W of size k , each random variable $\sum_{j' \in W} \deg_{G_1^t}(j', s)$ is stochastically dominated by X_s (for $mk \leq s \leq mn$). Therefore,

$$\Pr\left(\sum_{w \in W} \deg(w, t) > D\right) \leq \Pr(Y_t(k) > D) \leq 1/t^3 \binom{t}{k}.$$

Since there are at most $\binom{t}{k}$ sets $W \subseteq [t]$ of size k to consider, the desired property fails for a given t with probability at most $\sum_{k=1}^{t-1} t^{-3} \leq t^{-2}$. Hence, a.a.s. it does not fail for any $\omega \leq t \leq n$. As $Y_\omega(\omega) = 2m\omega$, the desired bound trivially holds (deterministically) for $t \leq \omega$, and the proof is finished. \square

We will also need a stronger result for sets of a certain type.

Lemma 5. *Fix any constant $c \in (0, 1)$ and $m \in \mathbb{Z}_{\geq 0}$. The following property holds a.a.s. for G_m^n . For any $cn \leq t \leq n$,*

$$Y_t := \sum_{w \in [cn]} \deg(w, t) \sim 2mn\sqrt{ct/n}.$$

Proof. In view of the identification between the models G_m^n (on the vertex set $1, 2, \dots, n$) and G_1^{mn} (on the vertex set $1', 2', \dots, mn'$), it will be useful to investigate the following random variable instead of Y_t : for $m \lfloor cn \rfloor \leq t \leq mn$, let

$$X_t = \sum_{j \in [cmn]} \deg_{G_1^t}(j', t).$$

Clearly, $Y_t = X_{tm}$. It follows that $X_{m \lfloor cn \rfloor} = Y_{\lfloor cn \rfloor} = 2m \lfloor cn \rfloor$. Moreover, for $m \lfloor cn \rfloor < t \leq mn$,

$$X_t = \begin{cases} X_{t-1} + 1 & \text{with probability } \frac{X_{t-1}}{2t-1}, \\ X_{t-1} & \text{otherwise.} \end{cases}$$

The conditional expectation is given by

$$\mathbf{E}(X_t | X_{t-1}) = (X_{t-1} + 1) \cdot \frac{X_{t-1}}{2t-1} + X_{t-1} \left(1 - \frac{X_{t-1}}{2t-1}\right) = X_{t-1} \left(1 + \frac{1}{2t-1}\right).$$

Taking expectation again, we derive that

$$\mathbf{E}X_t = \mathbf{E}X_{t-1} \left(1 + \frac{1}{2t-1}\right).$$

Hence, it follows that

$$\mathbf{E}(Y_t) = \mathbf{E}(X_{tm}) = 2m \lfloor cn \rfloor \prod_{s=m \lfloor cn \rfloor + 1}^{tm} \left(1 + \frac{1}{2s-1}\right) \sim 2cmn \left(\frac{tm}{cmn}\right)^{1/2} = 2mn\sqrt{ct/n}.$$

In order to transform X_t into something close to a martingale (to be able to apply the generalized Azuma-Hoeffding inequality (5)), we set for $m \lfloor cn \rfloor \leq t \leq mn$

$$Z_t = X_t - 2m \lfloor cn \rfloor - \sum_{k=m \lfloor cn \rfloor + 1}^t \sqrt{cmn/k}$$

(note that $Z_{m \lfloor cn \rfloor} = 0$) and use the following stopping time

$$T = \min \left\{ t > m \lfloor cn \rfloor : X_t \leq 2\sqrt{tcmn} + t^{2/3} \text{ or } t = mn \right\}.$$

Indeed, we have for $m \lfloor cn \rfloor < t \leq mn$

$$\mathbf{E}(Z_t - Z_{t-1} \mid Z_{t-1}) = \frac{X_{t-1}}{2t-1} - \sqrt{cmn/t} \leq (1/2 + o(1))t^{-1/3} < 0.51t^{-1/3},$$

provided $t \leq T$, and $|Z_t - Z_{t-1}| \leq 1$ as $t > cmn$. Let $t \wedge T$ denote $\min\{t, T\}$. We apply the generalized Azuma-Hoeffding inequality (5) to the sequence $(Z_{t \wedge T} : m \lfloor cn \rfloor \leq t \leq mn)$, with $c_t = 1$, $b_t = 0.51t^{-1/3}$ and $x = 0.1t^{2/3}$, to conclude that a.a.s. for all t such that $m \lfloor cn \rfloor \leq t \leq mn$

$$Z_{t \wedge T} - Z_{m \lfloor cn \rfloor} = Z_{t \wedge T} \leq \sum_{k \leq t} b_k + x \leq 0.77t^{2/3} + 0.1t^{2/3} \leq 0.9t^{2/3}.$$

To complete the proof we need to show that a.a.s., $T = mn$. The events asserted by the equation hold a.a.s. up until time T , as shown above. Thus, in particular, a.a.s.

$$\begin{aligned} X_T &= Z_T + 2m \lfloor cn \rfloor + \sum_{k=m \lfloor cn \rfloor + 1}^T \sqrt{cmn/k} \\ &\leq 0.9t^{2/3} + 2m \lfloor cn \rfloor + \sqrt{cmn} \int_{m \lfloor cn \rfloor}^T 1/\sqrt{k} dk + O(1) \\ &< 2\sqrt{tcmn} + t^{2/3}, \end{aligned}$$

which implies that $T = mn$ a.a.s. In particular, it follows that a.a.s., for any $cn \leq t \leq n$, $Y_t = X_{mt} < 2mn\sqrt{ct/n} + o(n)$. The lower bound can be obtained by applying the same argument symmetrically to $(-Z_{t \wedge T} : m \lfloor cn \rfloor \leq t \leq mn)$, and so the proof is finished. \square

Finally, we provide the last tool we need in order to extend the arguments in the paper from $G_{m\text{-old}}^n$ to $\pi_\sigma(G_{m_1, m_2}^n)$ (and thus to G_m^n if m_1 or m_2 equals 0).

Lemma 6. *Let $C > 0$ be a sufficiently large constant. Fix any constants $m, m_1, m_2 \in \mathbb{Z}_{\geq 0}$ with $m = m_1 + m_2 \geq 1$ and let $\sigma \in \{1, 2\}$. Let $\omega = \omega(n)$ be any function tending to infinity as $n \rightarrow \infty$, and suppose that the a.a.s. events in Lemma 4 hold in $G_m^n = \pi(G_{m_1, m_2}^n)$. Then, for any $\omega \leq j \leq n$, $R \subseteq [n] \setminus [j]$ and $\emptyset \neq Q \subseteq [n]$, the conditional probability that every vertex $v \in R$ has all neighbours in Q with respect to graph $\pi_\sigma(G_{m_1, m_2}^n)$ is at most*

$$\left(\left(1 + \frac{C}{|Q|} + o(1) \right) \left(2\sqrt{\frac{|Q|}{j}} + \sqrt{\frac{8|Q|}{mj}} \log \left(e \left(1 + \frac{j}{|Q|} \right) \right) \right) \right)^{m_\sigma r}.$$

Proof. Put $|R| = r$ and $|Q| = q \geq 1$. Since the lemma is trivial for $m_\sigma r = 0$, we will assume that $m_\sigma \geq 1$ and $r \geq 1$. For $t, m, x \geq 1$, define

$$f_{t,m}(x) = 2m\sqrt{xt} + \sqrt{8mxt} \log(e(1+t/x)).$$

We will use the following observation that we will be proved at the very end.

Claim 2. Given any constant $A > 0$, there exists $A' > 0$ sufficiently large such that for every $t, m \geq 1$ and every $1 \leq x \leq y$, $(1 + A/x)f_{t,m}(x) \leq (1 + A'/y)f_{t,m}(y)$.

Let H denote the a.a.s. event in the statement of Lemma 4, and let $A > 0$ be the constant in that same lemma. For each $t \in [n]$, define $W_t = Q \cap [t]$, so in particular we have $|W_t| \leq q$. In view of Claim 2, event H implies that, for every $t \in \mathbb{N}$ such that $j \leq t \leq n$ and $W_t \neq \emptyset$,

$$\begin{aligned} \sum_{w \in W_t} \deg(w, t) &\leq \left(1 + \frac{A}{m|W_t| + 1} \right) \left(2m\sqrt{|W_t|t} + \sqrt{8m|W_t|t} \log(et/|W_t|) \right) \\ &\leq \left(1 + \frac{A}{|W_t|} \right) f_{t,m}(|W_t|) \\ &\leq \left(1 + \frac{A'}{q} \right) f_{t,m}(q). \end{aligned}$$

(The last conclusion of the equation above is also true if $W_t = \emptyset$, so we can ignore the restriction we had on W_t . Moreover, let us mention that $|W_t| \leq t$ but it is not always true that $q \leq t$. This technical reason prevented us from defining function $f_{t,m}(x)$ as follows: $f_{t,m}(x) = 2m\sqrt{xt} + \sqrt{8mxt} \log(et/x)$.) For each $t \in [n]$ such that $t \geq j + 1$, let H_t be the event in $G_m^n = \pi(G_{m_1, m_2}^n)$ that, for every $j \leq s \leq t - 1$, we have

$$\sum_{w \in W_s} \deg(w, s) \leq \left(1 + \frac{A'}{q} \right) f_{s,m}(q).$$

By construction, $H_{j+1} \supseteq H_{j+2} \supseteq \dots \supseteq H_n \supseteq H$. For each $t \in [n]$, let E_t be the event that every edge of $\pi_\sigma(G_{m_1, m_2}^n)$ stemming from vertex t attaches this vertex to some vertex in Q . Our goal is to bound $\Pr\left(\bigcap_{t \in R} E_t \cap H\right)$. By labelling the r vertices in R as $t_1 < t_2 < \dots < t_r$ and since $H_{j+1} \supseteq H_{j+2} \supseteq \dots \supseteq H_n \supseteq H$, we conclude that

$$\Pr\left(\bigcap_{t \in R} E_t \cap H\right) \leq \Pr\left(\bigcap_{i=1}^r (E_{t_i} \cap H_{t_{i+1}})\right) \leq \prod_{i=1}^r \Pr(E_{t_i} \mid \widehat{H}_{t_i}), \quad (11)$$

where $\widehat{H}_{t_i} := H_{t_i} \cap \bigcap_{k=1}^{i-1} E_{t_k}$. Observe that event \widehat{H}_{t_i} implies that

$$\sum_{w \in Q \cap [t_i - 1]} \deg(w, t_i - 1) \leq \left(1 + \frac{A'}{q} \right) f_{t_i - 1, m}(q),$$

and (crucially!) only exposes information concerning edges created before time t_i , so the probability of E_{t_i} conditional upon \widehat{H}_{t_i} is at most

$$\begin{aligned} \left(\frac{\left(1 + \frac{A'}{q}\right) f_{t_i-1, m}(q) + 2m}{m(t_i - 1)} \right)^{m\sigma} &\leq \left(\left(1 + \frac{A'}{q}\right) \left(2\sqrt{\frac{q}{t_i-1}} + \sqrt{\frac{8q}{m(t_i-1)}} \log \left(e \left(1 + \frac{t_i-1}{q}\right)\right) \right) + \frac{2}{t_i-1} \right)^{m\sigma} \\ &\leq \left(\left(1 + \frac{A'}{q}\right) \left(2\sqrt{\frac{q}{j}} + \sqrt{\frac{8q}{mj}} \log \left(e \left(1 + \frac{j}{q}\right)\right) \right) + \frac{2}{j} \right)^{m\sigma}, \end{aligned} \quad (12)$$

where we used that $t_i - 1 \geq j$ (since $t_i \in R$) and the fact that $\frac{\log(e(1+x))}{\sqrt{x}}$ is decreasing with respect to x in $(0, +\infty)$ (which follows from elementary analysis). Hence, setting $C = A' + 1$, we obtain

$$\Pr \left(E_{t_i} \mid \widehat{H}_{t_i} \right) \leq \left(\left(1 + \frac{C}{q}\right) \left(2\sqrt{\frac{q}{j}} + \sqrt{\frac{8q}{mj}} \log \left(e \left(1 + \frac{j}{q}\right)\right) \right) \right)^{m\sigma}, \quad (13)$$

where the case $q \leq j$ follows immediately from (12) (since $2/j \leq \frac{1}{q} 2\sqrt{q/j}$), and the case $q \geq j$ is trivially true (since the right-hand side of (13) is greater than 1). Combining (11) and (13) together and using the fact that H holds a.a.s., we conclude that

$$\begin{aligned} \Pr \left(\bigcap_{t \in R} E_t \mid H \right) &= (1 + o(1)) \Pr \left(\bigcap_{t \in R} E_t \cap H \right) \\ &\leq \left(\left(1 + \frac{C}{q} + o(1)\right) \left(2\sqrt{\frac{q}{j}} + \sqrt{\frac{8q}{mj}} \log \left(e \left(1 + \frac{j}{q}\right)\right) \right) \right)^{m\sigma r}, \end{aligned}$$

which yields the statement of the lemma under the assumption that Claim 2 is valid.

Finally, we proceed to prove Claim 2. We will only sketch the main steps in the argument and leave the details to the reader. We may increase A if needed and assume it is a sufficiently large constant (independent of t and m). Take $A' = A^5$. For each $t, m, x \geq 1$, define $g_{t,m}(x) = (1 + A/x)f_{t,m}(x)$. Elementary (but rather tedious) computations show that

$$\frac{d}{dx} f_{t,m}(x) \geq 0 \quad \text{for all } x \geq 1 \quad \text{and} \quad \frac{d}{dx} g_{t,m}(x) \geq 0 \quad \text{for all } x \geq 7A \geq \frac{\sqrt{m/2+1}+3\frac{t}{t+x}}{\sqrt{m/2+1}-\frac{t}{t+x}} A.$$

Fix any $t, m \geq 1$ and any $1 \leq x \leq y$. If $x \geq 7A$, then trivially

$$g_{t,m}(x) \leq g_{t,m}(y) \leq (1 + A'/y)f_{t,m}(y),$$

as desired. Also, if $y \leq A^4$,

$$g_{t,m}(x) \leq (1 + A)f_{t,m}(x) \leq (1 + A)f_{t,m}(y) \leq (1 + A^5/y)f_{t,m}(y).$$

Otherwise, suppose $x \leq 7A$ and $y \geq A^4$. Then, assuming that A is a large enough constant, we get

$$g_{t,m}(x) \leq (1 + A)f_{t,m}(x) \leq (1 + A)f_{t,m}(7A) \leq (1 + A)\sqrt{7A} \left(2m\sqrt{t} + \sqrt{8mt} \log(et) \right) \leq f_{t,m}(A^4),$$

where the last step follows from simple computations by considering separately the cases $t \leq A^5$ and $t \geq A^5$. Hence,

$$g_{t,m}(x) \leq f_{t,m}(A^4) \leq f_{t,m}(y) \leq (1 + A'/y)f_{t,m}(y),$$

the proof of the claim is completed, and so the proof is finished. \square

3 Upper bound

3.1 Expansion properties

Let us start with investigating some properties of $G_{m-\text{old}}^n$ and G_m^n that will turn out to be important in determining the existence of perfect matchings and Hamiltonian cycles. All the results of this section are stated for $G_{m-\text{old}}^n$ and $\pi_1(G_{2m,m'}^n)$ where $m \in \mathbb{N}$ and $m' \in \mathbb{Z}_{\geq 0}$, so they also apply to G_{2m}^n (which can be obtained from $\pi_1(G_{2m,m'}^n)$ by setting $m' = 0$). First we will need the following technical lemma.

Lemma 7. *Fix any constants $0 < x, y, d, \alpha < 1$, $m \in \mathbb{N}$, and $m' \in \mathbb{Z}_{\geq 0}$ satisfying*

$$(1-d)\varphi(-x)m-d > 0 \quad \text{and} \quad \alpha < \left(\frac{0.99y}{e}\right)^{1/((1-d)\varphi(-x)m-d)},$$

where φ is defined in (3). Then, the following holds a.a.s. for $G_{m-\text{old}}^n$ and for $\pi_1(G_{2m,m'}^n)$. For any integer k with $1 \leq k \leq \alpha n$ and $j = \lfloor k(n/k)^d \rfloor$, there are at most yk vertices in $[j]$ with fewer than $(1-x)m \log(n/j)$ neighbours in $[n] \setminus [j]$.

Proof. Let us start with showing the desired property for $G_{m-\text{old}}^n$. Pick a constant $A < \varphi(-x)$ but sufficiently close to $\varphi(-x)$ so that the following properties are satisfied:

$$(1-d)Am-d > 0 \quad \text{and} \quad \alpha < \left(\frac{0.99y}{e}\right)^{1/((1-d)Am-d)}. \quad (14)$$

Let us concentrate on any k in the range of consideration; that is, $1 \leq k \leq \alpha n$. Let X_v be the number of neighbours in $[n] \setminus [j]$ of a given vertex $v \in [j]$. By (6) (and the claim below), X_v is a sum of independent Bernoulli random variables with parameter

$$1 - \left(1 - \frac{1}{t-1}\right)^m \sim \frac{m}{t}$$

for $t \in [n] \setminus [j]$ (note that $j \gg 1$). Hence, we get that

$$\mathbf{E}[X_v] \sim \sum_{t=j+1}^n \frac{m}{t} \sim m \log(n/j).$$

By the generalization of the Chernoff bound (2) we get that the probability of v having less than $(1-x)m \log(n/j)$ neighbours in $[n] \setminus [j]$ is at most

$$\exp\left(- (1+o(1))m \log(n/j)\varphi(-x)\right) < (j/n)^{Am},$$

since $A < \varphi(-x)$. We will call such a vertex *bad*. Using the fact that the events that two or more vertices are bad are negatively correlated, the probability of having at least yk bad vertices in $[j]$ is at most

$$\binom{j}{\lceil yk \rceil} (j/n)^{Am \lceil yk \rceil} \leq \left(\frac{ej(j/n)^{Am}}{\lceil yk \rceil}\right)^{\lceil yk \rceil} \leq \left(\frac{e}{y} \left(\frac{k}{n}\right)^{(1-d)Am-d}\right)^{yk} =: a_k.$$

Note that, if $k + 1 \leq \varepsilon n$ for a small constant $\varepsilon > 0$,

$$\begin{aligned} \frac{a_{k+1}}{a_k} &= \left(\frac{k+1}{k}\right)^{yk((1-d)Am-d)} \left(\frac{e}{y} \left(\frac{k+1}{n}\right)^{(1-d)Am-d}\right)^y \\ &\leq \exp(y((1-d)Am-d)) \left(\frac{e}{y}\right)^y \varepsilon^{y((1-d)Am-d)} \\ &= \left(\frac{e}{y}\right)^y \exp(y((1-d)Am-d)(1+\log \varepsilon)) < 0.99, \end{aligned} \quad (15)$$

by the first inequality in (14) and provided that ε is sufficiently small. Hence, $\sum_{k=1}^{\lfloor \varepsilon n \rfloor} a_k = O(a_1) = o(1)$. On the other hand,

$$\sum_{k=\lfloor \varepsilon n \rfloor + 1}^{\lfloor \alpha n \rfloor} a_k \leq \sum_{k=\lfloor \varepsilon n \rfloor + 1}^{\lfloor \alpha n \rfloor} \left(\frac{e}{y} \alpha^{(1-d)Am-d}\right)^{yk} \leq \sum_{k=\lfloor \varepsilon n \rfloor + 1}^{\lfloor \alpha n \rfloor} 0.99^{yk} = O(0.99^{\varepsilon y n}) = o(1),$$

by the second inequality in (14). It follows that the probability that the desired property does not hold for some k is $o(1)$. The desired property holds for $G_{m\text{-old}}^n$.

Adjusting the proof for $\pi_1(G_{2m, m'}^n)$ is straightforward, by using Lemma 3 (with $\sigma = 1$, $m_1 = 2m$ and $m_2 = m'$) instead of (6) that we used above. The only difference is that now X_v is stochastically lower bounded by a sum of independent Bernoulli random variables with parameter

$$1 - \left(1 - \frac{1}{2t}\right)^{2m} \sim \frac{m}{t}$$

for $t \in [n] \setminus [j]$. It follows that $\mathbf{E}[X_v] \geq (1 + o(1))m \log(n/j)$, and the rest of the proof continues as before. \square

Next, we will use Lemma 7 to derive an expansion property of $G_{m\text{-old}}^n$ and $\pi_1(G_{2m, m'}^n)$ that will play a key role in the argument. We will show that for all sets of vertices K of moderate size, $N(K)$ is large. (Recall the definition of $N(K)$ from Section 2.1.)

Lemma 8. *Let $\ell \in \{1, 2\}$. Fix any constants $0 < x, y, z, d, \alpha < 1$ and $m \in \mathbb{N}$ satisfying*

$$y < z < 1 - \frac{\ell + 1}{dm}, \quad (16)$$

$$(1-d)\varphi(-x)m - d > 0, \quad (17)$$

and

$$\alpha < \begin{cases} \left(\frac{0.99y}{e}\right)^{\frac{1}{(1-d)\varphi(-x)m-d}} & (18a) \\ \exp\left(-\frac{\ell + 1 - z}{(1-x)(1-d)(z-y)}\right) & (18b) \\ \frac{1}{\ell + 1} - \frac{1}{d(1-z)m} & (18c) \\ \exp\left(-\frac{(1-z)m \log(\ell + 1) + \ell + 1 - \ell \log \ell}{d(1-z)m - \ell - 1}\right). & (18d) \end{cases}$$

Then, the following holds a.a.s. for $G_{m\text{-old}}^n$. Every set of vertices K with $1 \leq |K| \leq \alpha n$ satisfies $|N(K)| \geq \ell|K|$.

Moreover, suppose that we can replace (18c) and (18d) by the following stronger conditions

$$\alpha < \begin{cases} \frac{1}{\ell+1} - \frac{8/3}{d(1-z)m} & (18'c) \\ \exp\left(-\frac{(1-z)m \log((2+w)^2(\ell+1)) + \ell+1 - \ell \log \ell}{d(1-z)m - (\ell+1)}\right) & (18'd) \end{cases}$$

where

$$w = \sqrt{\frac{8}{m}} \log\left(e\left(1 + \frac{1}{(\ell+1)\alpha^d}\right)\right). \quad (19)$$

Then, the previous a.a.s. statement is also a.a.s. valid for $\pi_1(G_{2m,m'}^n)$ (for any $m' \in \mathbb{Z}_{\geq 0}$).

Note that conditions (16–19) can be trivially satisfied by taking any arbitrary $x, y, z, d \in (0, 1)$ and $w > 0$ such that $y < z$, and then choosing $m \in \mathbb{N}$ sufficiently large and $\alpha > 0$ sufficiently small. We leave the result with such flexibility as our goal will be to tune everything to get the constant m as small as possible.

Moreover, let us note that we assumed that $\ell \in \{1, 2\}$ as this is what will be used in our specific application of this lemma. However, it is straightforward to verify that the lemma holds for any real number $\ell > 0$ (by placing floors and ceilings in the appropriate places of the argument).

Proof. Let us start with showing the desired property for $G_{m\text{-old}}^n$. A set K of vertices of size k is of type 1 if it contains at least zk vertices in $[j]$, where $j = \lfloor k(n/k)^d \rfloor$. Otherwise, K is of type 2. First, we will focus on sets of type 1. Observe that constants x, y, d, α, m satisfy the requirements of Lemma 7 (by assumptions (17) and (18a)). Since our aim is to obtain a statement that holds a.a.s., we may assume that the conclusion of Lemma 7 holds, and proceed to prove the desired statement (deterministically) for any given set of type 1. Let K be a fixed set of type 1 of size $1 \leq k \leq \alpha n$. It follows that at least $(z-y)k$ vertices in $K \cap [j]$ have at least

$$(1-x)m \log(n/j) \geq (1-x)(1-d)m \log(1/\alpha) > \frac{\ell+1-z}{z-y}m$$

neighbours in $[n] \setminus [j]$ (by assumption (18b)). Denote this set of at least $(z-y)k$ vertices by K_0 , and let K'_0 be the set of neighbours of K_0 in $[n] \setminus [j]$. Looking at the degrees of the bipartite graph induced by the parts K_0 and K'_0 , we conclude that $|K'_0| \geq \frac{\ell+1-z}{z-y}|K_0| \geq (\ell+1-z)k$. Therefore, $|N(K)| \geq |K'_0 \setminus K| \geq |K'_0| - |K \setminus [j]| \geq (\ell+1-z)k - (1-z)k = \ell k$, as desired. This proves the claim for all sets of type 1.

Now, consider a set K of size $1 \leq k \leq \alpha n$ of type 2 (hence, K must contain at least $(1-z)k$ vertices in $[n] \setminus [j]$). Observe that if $|N(K)| < \ell k$, then there must exist a set $S = S(K)$ of size $\ell k - 1$ with $S \cap K = \emptyset$ such that all vertices in $K \setminus [j]$ have all of their neighbours in $K \cup S$. An important observation is that, when a vertex in $K \setminus [j]$ generated m edges to older vertices, all vertices in $[j]$ were available for possible destinations. Hence, the probability that one such vertex has all the neighbours in $K \cup S$ is at most

$$\left(\frac{(\ell+1)k-1}{j}\right)^m \leq \left(\frac{(\ell+1)k-1}{k(n/k)^d-1}\right)^m \leq \left((\ell+1)(k/n)^d\right)^m. \quad (20)$$

As a result, the expected number of sets K of type 2 and size k with $|N(K)| < \ell k$ is at most

$$\begin{aligned} b_k &:= \binom{n}{k} \binom{n-k}{\ell k} \left((\ell+1)(k/n)^d \right)^{(1-z)km} \\ &\leq \left(\frac{en}{\ell^{\ell/(\ell+1)} k} \right)^{(\ell+1)k} \left((\ell+1)(k/n)^d \right)^{(1-z)km} \\ &= \left(\left(\frac{e^{\ell+1}}{\ell^\ell} \right) (\ell+1)^{(1-z)m} (k/n)^{(1-z)dm-\ell-1} \right)^k =: b'_k. \end{aligned} \quad (21)$$

Note that (16) implies that the exponent of k/n in the expression above is positive. Hence, proceeding analogously as in (15), we can show that there exists a small enough constant $\varepsilon > 0$ such that $b'_{k+1}/b'_k < 0.99$ whenever $k+1 \leq \varepsilon n$. This implies that

$$b'_k \leq b'_1 (0.99)^{k-1} \quad \text{for } 1 \leq k \leq \varepsilon n, \quad (22)$$

and so $\sum_{k=1}^{\lfloor \varepsilon n \rfloor} b'_k = O(b'_1) = o(1)$. It follows that a.a.s. no set of type 2 and size at most εn fails to have the desired property. To deal with larger sets, it is more convenient to estimate b_k directly instead of b'_k . For $\varepsilon n \leq k \leq \alpha n$, we use Stirling's formula ($s! \sim \sqrt{2\pi s}(s/e)^s$, as $s \rightarrow \infty$) and obtain

$$b_k = \frac{\Theta(1/n) \left((\ell+1)(k/n)^d \right)^{(1-z)mk}}{(k/n)^k (\ell k/n)^{\ell k} (1 - (\ell+1)k/n)^{n-(\ell+1)k}} = e^{n(f(k/n)+o(1))},$$

where

$$f(\rho) = (1-z)m\rho \log \left((\ell+1)\rho^d \right) - \rho \log \rho - \ell\rho \log(\ell\rho) - (1 - (\ell+1)\rho) \log(1 - (\ell+1)\rho). \quad (23)$$

(Note that in the above estimate of b_k we implicitly used the fact that $n > (\ell+1)k$, which follows from (18c).) The second derivative of f is

$$f''(\rho) = \frac{md(1-z)(1 - (\ell+1)\rho) - (\ell+1)}{\rho(1 - (\ell+1)\rho)}, \quad (24)$$

which is strictly positive for $\rho \in (0, \alpha]$ (by (18c)), so f is concave up in that interval. Therefore,

$$b_k \leq e^{o(n)} \max \{ b_{\lfloor \varepsilon n \rfloor}, b_{\lfloor \alpha n \rfloor} \} = (1+o(1))^n \max \{ b_{\lfloor \varepsilon n \rfloor}, b_{\lfloor \alpha n \rfloor} \} \quad \text{for all } \varepsilon n \leq k \leq \alpha n.$$

From (22), we get that $b_{\lfloor \varepsilon n \rfloor} \leq b'_1 (0.99)^{\varepsilon n - 2} = o((0.99)^{\varepsilon n})$. Next, we proceed to bound $b_{\lfloor \alpha n \rfloor} = e^{n(f(\alpha)+o(1))}$. Observe that $0 < \alpha < 1/(\ell+1)$ by (18c) and, for every α in that range, elementary analysis shows that

$$-(1 - (\ell+1)\alpha) \log(1 - (\ell+1)\alpha) < (\ell+1)\alpha. \quad (25)$$

Hence,

$$f(\alpha) < (1-z)m\alpha \log \left((\ell+1)\alpha^d \right) - \alpha \log \alpha - \ell\alpha \log(\ell\alpha) + (\ell+1)\alpha < 0,$$

where the first inequality holds by (25) and the second one by (18d). Putting everything together, for each $\varepsilon n \leq k \leq \alpha n$,

$$b_k \leq (0.99 + o(1))^{\varepsilon n} + \left(e^{f(\alpha)} + o(1) \right)^n = o(1/n).$$

It follows that $\sum_{k=\lfloor \varepsilon n \rfloor}^{\lfloor \alpha n \rfloor} b_k = o(1)$, and so a.a.s. all sets of type 2 of size at most αn satisfy the desired property, and the proof is finished for $G_{m-\text{old}}^n$.

Adjusting the proof for $\pi_1(G_{2m,m'}^n)$ is more complicated than in the previous lemma. As Lemma 7 can be applied to $\pi_1(G_{2m,m'}^n)$, the proof for type 1 sets is not affected. For type 2 sets, we need to use Lemmas 4 and 6 in order to obtain an analogue of (21). Let H be the a.a.s. event in the statement of Lemma 4 (replacing m by $2m + m'$) with $\omega := \log \lfloor n^d \rfloor \rightarrow \infty$. As usual, as we aim for a statement that holds a.a.s., we may assume that event H holds. Consider any set K of size $1 \leq k \leq \alpha n$ of type 2 and any set S of size $\ell k - 1$ such that $S \cap K = \emptyset$. Put as before $j = \lfloor k(n/k)^d \rfloor$, so $j \geq \lfloor n^d \rfloor \rightarrow \infty$, as $n \rightarrow \infty$ (and also $\log j \geq \omega$). Setting $R = K \setminus [j]$ and $Q = K \cup S$ in Lemma 6 (with $\sigma = 1$, $m_1 = 2m$ and $m_2 = m'$), we conclude that, conditional upon H , the probability p that every vertex in $K \setminus [j]$ has all neighbours in $K \cup S$ with respect to graph $\pi_1(G_{2m,m'}^n)$ satisfies

$$\begin{aligned} p &\leq \left(\left(1 + \frac{C}{|Q|} + o(1) \right) \left(2\sqrt{\frac{|Q|}{j}} + \sqrt{\frac{8|Q|}{mj}} \log \left(e \left(1 + \frac{j}{|Q|} \right) \right) \right) \right)^{2m|R|} \\ &\leq \left(\left(1 + \frac{C}{k} + o(1) \right) 2\sqrt{(\ell+1)(k/n)^d} \left(1 + \sqrt{\frac{2}{m}} \log \left(e \left(1 + \frac{(n/k)^d}{\ell+1} \right) \right) \right) \right)^{2m|R|}, \end{aligned} \quad (26)$$

where we used that $|Q|/j \leq \frac{(\ell+1)k-1}{k(n/k)^{d-1}} \leq (\ell+1)(k/n)^d$ and the fact that $\frac{\log(e(1+t))}{\sqrt{t}}$ is decreasing in $t \in (0, +\infty)$ (as observed below (12)). Since $|R| \geq (1-z)k$, replacing $|R|$ by $(1-z)k$ in (26) gives a valid bound (even in the case that the base of the power in the right-hand side of (26) is greater than 1, since p is a probability). Hence, we conclude that

$$p \leq \left(\left(1 + \frac{C}{k} + o(1) \right)^2 4(\ell+1)(k/n)^d \left(1 + \sqrt{\frac{2}{m}} \log \left(e \left(1 + \frac{(n/k)^d}{\ell+1} \right) \right) \right) \right)^{2(1-z)km} \quad (27)$$

$$\leq \left(C'(k/n)^d \log^2(n/k) \right)^{(1-z)km}, \quad (28)$$

for some constant $C' > 0$ that may depend on C , ℓ and α . The role of p (and its bounds (27) and (28)) will be very similar to that of $((\ell+1)(k/n)^d)^{(1-z)km}$ in the computations that we did for the $G_{m\text{-old}}^n$ model (see (20) and (21)). Indeed, for the $\pi_1(G_{2m,m'}^n)$ model, the expected number (conditional upon H) of sets K of type 2 and size k with $|N(K)| < \ell k$ is at most

$$\begin{aligned} b_k &:= \binom{n}{k} \binom{n-k}{\ell k} p \\ &\leq \left(\left(\frac{e^{\ell+1}}{\ell^\ell} \right) C'(k/n)^{(1-z)dm-\ell-1} \log^{2m(1-z)}(n/k) \right)^k =: b'_k. \end{aligned} \quad (29)$$

Our goal is to show that $\sum_{k=1}^{\lfloor \alpha n \rfloor} b_k = o(1)$. This, combined with the fact that H is a.a.s. true, will yield the desired result for $\pi_1(G_{2m,m'}^n)$, and finish the proof of the lemma. The argument to bound b_k (or b'_k) is analogous to the one we used for $G_{m\text{-old}}^n$, so we will only sketch the main differences. By inspecting the ratios b'_{k+1}/b'_k , we can easily check that (22) remains valid in the present context for $1 \leq k \leq \varepsilon n$ provided that $\varepsilon > 0$ is a sufficiently small constant given C' , m , ℓ , z and d . Therefore $\sum_{k=1}^{\lfloor \varepsilon n \rfloor} b'_k = O(b'_1) = o(1)$ as before. To analyze the case $\varepsilon n \leq k \leq \alpha n$, we plug the bound obtained in (27) into the definition of b_k in (29). Applying Stirling's formula to the resulting bound and performing elementary manipulations, we conclude that

$$b_k \leq e^{n(\hat{f}(k/n)+o(1))},$$

where $\hat{f}(\rho) = f(\rho) + g(\rho)$ with $f(\rho)$ defined in (23) and

$$g(\rho) := 2(1-z)m\rho \log \left(2 + \sqrt{\frac{8}{m}} \log \left(e \left(1 + \frac{1}{(\ell+1)\rho^d} \right) \right) \right).$$

The function \hat{f} will play the same role as f did in the argument for the for $G_{m\text{-old}}^n$ model. In order to adapt the previous proof to the current model and show that $\sum_{k=\lfloor \varepsilon n \rfloor}^{\lfloor \alpha n \rfloor} b_k = o(1)$, we only need to check that $\hat{f}(\alpha) < 0$ and $(\hat{f})''(\rho) > 0$ for all $\rho \in (0, \alpha]$. We proceed to verify these two claims. Firstly, from (19), we get that

$$g(\alpha) = (1-z)m\alpha \log(2+w)^2.$$

Combining this and (25) yields

$$\hat{f}(\alpha) < (1-z)m\alpha \log \left((2+w)^2(\ell+1)\alpha^d \right) - \alpha \log \alpha - \ell\alpha \log(\ell\alpha) + (\ell+1)\alpha < 0,$$

where the second inequality follows directly from (18'd). Secondly, we differentiate g twice and obtain

$$g''(\rho) = -\frac{2md(1-z)\sqrt{8/m}}{\rho(1+(\ell+1)\rho^d)(2+\lambda)} \left(1 + \frac{d\sqrt{8/m}}{(1+(\ell+1)\rho^d)(2+\lambda)} - \frac{d(\ell+1)\rho^d}{1+(\ell+1)\rho^d} \right),$$

where

$$\lambda = \lambda(\rho) = \sqrt{\frac{8}{m}} \log \left(e \left(1 + \frac{1}{(\ell+1)\rho^d} \right) \right).$$

Note that, for all $\rho \in (0, \alpha]$,

$$\frac{\sqrt{8/m}}{(1+(\ell+1)\rho^d)(2+\lambda)} = \frac{1}{(1+(\ell+1)\rho^d) \left(\sqrt{m/2} + \log \left(e \left(1 + \frac{1}{(\ell+1)\rho^d} \right) \right) \right)} \leq 1/4,$$

since $m \geq 1$ and $(1+t) \left(\sqrt{1/2} + \log(e(1+1/t)) \right) \geq 4.12 > 4$ for all $t \in (0, \infty)$ (by elementary analysis). Therefore,

$$\begin{aligned} g''(\rho) &\geq -\frac{2md(1-z)\sqrt{8/m}}{\rho(1+(\ell+1)\rho^d)(2+\lambda)} \left(1 + \frac{\sqrt{8/m}}{(1+(\ell+1)\rho^d)(2+\lambda)} \right) \\ &\geq -\frac{2md(1-z)}{\rho} (1/4)(1+1/4) \\ &\geq -(5/8) \frac{md(1-z)}{\rho}. \end{aligned}$$

Combining this bound with (24), we conclude that

$$(\hat{f})''(\rho) = f''(\rho) + g''(\rho) \geq \frac{(3/8)md(1-z)(1-(\ell+1)\rho) - (\ell+1)}{\rho(1-(\ell+1)\rho)},$$

which is positive for all $\rho \in (0, \alpha]$ by (18'c). This finishes the proof of the lemma for the $\pi_1(G_{2m, m'}^n)$ model. \square

Finally, we include some properties of large sets of vertices, whose size is not covered in the previous expansion lemma. We will show that (a.a.s. in $G_{m\text{-old}}^n$ or $\pi_1(G_{2m,m'}^n)$) large sets of vertices must induce some edges and disjoint pairs of large sets must have some edges across. These properties will guarantee the existence of a long path and a large matching that will be extended later to a Hamilton cycle and a perfect matching, respectively.

Lemma 9. *For any constant $m \geq 12$, let $\beta = \beta(m) \in (0, 1/2)$ be such that*

$$2\beta \log \beta + (1 - 2\beta) \log(1 - 2\beta) + \beta^2 m/4 = 0,$$

and for any constant $m \geq 1$, let $\gamma = \gamma(m) \in (0, 1)$ be such that

$$\gamma \log \gamma + (1 - \gamma) \log(1 - \gamma) + \gamma^2 m/2 = 0.$$

Fix $m, m' \in \mathbb{Z}_{\geq 0}$. Then, the following two properties hold a.a.s. for $G_{m\text{-old}}^n$ and for $\pi_1(G_{2m,m'}^n)$.

- (i) *If $m \geq 12$, then there is no pair of disjoint sets of vertices A, B , both of size at least βn , with no edges between A and B . As a result, the length of a longest path is at least $n - 2\lceil \beta n \rceil \sim (1 - 2\beta)n$.*
- (ii) *If $m \geq 1$, then no set C of size at least γn forms an independent set. As a result, the size of a maximum matching is at least $(1 - \gamma)n/2$.*

Before we prove the lemma, let us make a few observations. First, the lower bound of 12 on m is needed for $\beta = \beta(m)$ to be well defined. Second, note that the function $f(x) = -2x \log x - (1 - 2x) \log(1 - 2x)$ is maximized at $f(1/3) = \log 3$. Hence, $\beta \leq \sqrt{4 \log 3/m}$ and, in particular, β tends to zero as $m \rightarrow \infty$. In fact, one can show that $\beta \sim 8 \log m/m$ as $m \rightarrow \infty$. Similarly, we get that $\gamma \leq \sqrt{2 \log 2/m}$, and $\gamma \sim 2 \log m/m$ as $m \rightarrow \infty$. Finally, let us mention that the second part of part (i) uses simple ideas that proved to be extremely useful in many current applications (see, for example, [16, 17, 15, 26, 31]). Such techniques were used for the first time in [4, 5] (see the recent book [25] that covers several tools including this one, or another recent book on random graphs [21, Chapter 6.3]).

Proof. Let us start with the $G_{m\text{-old}}^n$ model and part (i). Consider any pair of disjoint sets A and B , both of size $\lceil \beta n \rceil$. Let U be the set of $(|A| + |B|)/2 = \lceil \beta n \rceil$ oldest vertices in $A \cup B$ and let $U' = (A \cup B) \setminus U$ contain the youngest half. Without loss of generality, A contains at least $\lceil \beta n/2 \rceil$ vertices in U and B contains at least $\lceil \beta n/2 \rceil$ vertices in U' . Therefore, the probability that there are no edges between $A \cap U$ and $B \cap U'$ is at most

$$(1 - \lceil \beta n/2 \rceil/n)^{m \lceil \beta n/2 \rceil} \leq (1 - \beta/2)^{m \beta n/2} \leq \exp(-\beta^2 mn/4) \tag{30}$$

by (6). Hence, using Stirling's formula ($s! \sim \sqrt{2\pi s}(s/e)^s$, as $s \rightarrow \infty$), the expected number of pairs of sets A, B that do not have the desired property is at most

$$\binom{n}{\lceil \beta n \rceil} \binom{n - \lceil \beta n \rceil}{\lceil \beta n \rceil} \exp(-\beta^2 mn/4) = o\left(\beta^{-2\beta n} (1 - 2\beta)^{-(1-2\beta)n} \exp(-\beta^2 mn/4)\right) = o(1),$$

by the definition of β . The first claim of part (i) follows by Markov's inequality.

The second claim of part (i) follows from the first claim (deterministically). For a contradiction, suppose that the first claim holds and there is no path of length $h = n - 2\lceil \beta n \rceil$ (or equivalently all paths contain at most h vertices). We perform the following algorithm and construct a path P . Let v_1 be an arbitrary vertex. Initially, let $P = (v_1)$, $U = V \setminus \{v_1\}$, and $W = \emptyset$. Then, if there

is an edge from v_1 to U (say from v_1 to $v_2 \in U$), we extend the path to $P = (v_1, v_2)$ and remove v_2 from U . We continue extending the path P this way for as long as possible. Since there is no path of length $h \leq n - 1$, we must reach the point of the process in which P cannot be extended, that is, there is a path from v_1 to v_k ($k \leq h$) and there is no edge from v_k to U . This time, v_k is moved to W and we try to continue extending the path from v_{k-1} , reaching another critical point in which another vertex will be moved to W , etc. If P is reduced to a single vertex v_1 and no edge to U is found, we move v_1 to W and simply re-start the process from another vertex from U , again arbitrarily chosen. An obvious but important observation is that during this algorithm there is never an edge between U and W . Moreover, in each step of the process, the size of U decreases by 1 or the size of W increases by 1. Finally, the number of vertices of the path P is always at most $h = n - 2\lceil\beta n\rceil$ by assumption. Hence, at some point of the process both U and W must have size at least $\lceil\beta n\rceil$. But this contradicts the first claim of part (i) and so this part is finished.

Now, let us move to part (ii). Consider any set C of size $k = \lceil\gamma n\rceil$. Again using (6), the probability C forms an independent set is at most

$$\prod_{i=1}^{k-1} \left(1 - \frac{i}{n}\right)^m \leq \exp\left(-\frac{m}{n} \sum_{i=1}^{k-1} i\right) = \exp\left(-\frac{mk(k-1)}{2n}\right) = O\left(\exp(-\gamma^2 mn/2)\right). \quad (31)$$

The expected number of sets that do not have the desired property is of order at most

$$\binom{n}{\lceil\gamma n\rceil} \exp(-\gamma^2 mn/2) = o\left(\gamma^{-\gamma n} (1-\gamma)^{-(1-\gamma)n} \exp(-\gamma^2 mn/2)\right) = o(1),$$

by the definition of γ , and the first claim of part (ii) is proved.

The second claim of part (ii) is now a trivial (deterministic) implication. Indeed, if there is no matching of size at least $(1-\gamma)n/2$, then any maximum (or even maximal) matching leaves at least γn vertices that are not matched that form an independent set. This would contradict the first claim and so the proof of the theorem for $G_{m\text{-old}}^n$ is finished. Adjusting the proof for $\pi_1(G_{2m,m'}^n)$ is trivial, and only requires using Lemma 3 (with $\sigma = 1$, $m_1 = 2m$ and $m_2 = m'$) instead of (6) in order to obtain the analogues of (30) and (31). \square

Tuning the constants in the previous lemmas requires some patience but is straightforward. (Maple or some other software might be helpful.) For example, we get the following set of constants for the $G_{m\text{-old}}^n$ model:

Lemma 10.

- (a) *The following constants satisfy conditions (16–18) in Lemma 8: $m = 120$, $\ell = 1$, $\alpha = 0.0538$, $x = 0.22791$, $y = 0.020063$, $z = 0.851649$ and $d = 0.387967$. Moreover, for $m = 120$ we have $\gamma = \gamma(m) \leq 0.06238$ in Lemma 9.*
- (b) *The following constants also satisfy conditions (16–18) in Lemma 8: $m = 2,900$, $\ell = 2$, $\alpha = 0.032003$, $x = 0.048929$, $y = 0.003625$, $z = 0.965269$ and $d = 0.353628$. Moreover, for $m = 2,900$ we have $\beta = \beta(m) \leq 0.014414$ in Lemma 9.*

Similarly, for $\pi_1(G_{2m,m'}^n)$, we get:

Lemma 11.

- (c) The following constants satisfy conditions (16–19) in Lemma 8: $m = 500$, $\ell = 1$, $\alpha = 0.016801$, $x = 0.149159$, $y = 0.008856$, $z = 0.905885$ and $d = 0.649188$. Moreover, for $m = 500$ we have $\gamma = \gamma(m) \leq 0.019675$ in Lemma 9.
- (d) The following constants also satisfy conditions (16–19) in Lemma 8: $m = 14,000$, $\ell = 2$, $\alpha = 0.008874$, $x = 0.026228$, $y = 0.001272$, $z = 0.980855$ and $d = 0.551906$. Moreover, for $m = 14,000$ we have $\beta = \beta(m) \leq 0.003760$ in Lemma 9.

3.2 Perfect matchings

In this subsection, we deal with perfect matchings. Let us start with the following deterministic result that can be found, for example, in [21, Lemma 6.3]. Let $G = (V, E)$ be any graph. Let $A = A(G) \subseteq V$ be the set of vertices that are isolated by some maximum matching. For $v \in A$, let

$$B(v) = \{w \in V \setminus \{v\} : \text{there exists a maximum matching that isolates both } v \text{ and } w\}.$$

Observe that $B(v) \subseteq A$, and moreover $w \in B(v)$ implies that $v \in B(w)$. The set $B(v)$ is very important for understanding whether a few additional random edges, with v as one of their endpoints, have a chance to increase the size of a maximum matching. Indeed, clearly, the size increases by one if an edge between v and some vertex in $B(v)$ is added to the graph. Moreover, the following lemma holds. (We provide the proof for completeness.)

Lemma 12 ([21]). *Let G be a graph without a perfect matching. If $v \in A(G)$ and $B(v) \neq \emptyset$, then*

$$|N(B(v))| < |B(v)|.$$

Note that the condition $B(v) \neq \emptyset$ in the lemma has the sole purpose of excluding the case that G has odd order n and a maximum matching of size $(n - 1)/2$. Since our definition of perfect matching in this paper includes this situation, then the aforementioned condition is redundant.

Proof. Let v be a vertex that is isolated by a maximum matching M . Let $S_0 = V(G) - (V(M) \cup \{v\})$, and observe that $S_0 \neq \emptyset$ by our assumption on $B(v)$. Let S_1 be the set of vertices reachable from S_0 by a nontrivial even-length M -alternating path, and note that $S_0 \cap S_1 = \emptyset$. Furthermore, if $w \in B(v)$, let M' be a maximum matching isolating both v and w . If $w \notin S_0$, then in $M \Delta M'$ there is a path from S_0 to w , hence $w \in S_1$. Therefore $B(v) = S_0 \cup S_1$.

Let $x \in N_G(S_1)$ and let $y \in S_1$ be a neighbour of x . Let P be a nontrivial even-length M -alternating path from S_0 to y . Then $x \in V(M)$, as otherwise P followed by yx would be an M -augmenting path (contradicting the maximality of M). So let y' satisfy $xy' \in M$. We will show that $y' \in S_1$; this implies that M defines an injection $N_G(S_1) \rightarrow S_1$ by $x \mapsto y'$, and hence $|N_G(S_1)| \leq |S_1|$. Note that if x or y' lie on P , then xy' is an edge of P , so since $x \notin S_1$ we may truncate P to end with xy' and we are done. Otherwise, we may extend P by xy' .

Finally, note that $N_G(S_0) \subseteq S_1 \cup N_G(S_1)$ since each $z \in N_G(S_0)$ is adjacent in M to some $w \in S_1$. Thus

$$N_G(B(v)) = N_G(S_0 \cup S_1) = (N_G(S_0) \cup N_G(S_1)) \setminus (S_0 \cup S_1) \subseteq N_G(S_1) \setminus S_0$$

and so

$$|B(v)| = |S_0| + |S_1| > |S_1| \geq |N_G(S_1)| > |N_G(B(v))|. \quad \square$$

Now, we are ready to prove the main result for perfect matchings. Recall that we say that a graph on n vertices has a perfect matching if it has a matching of size $\lfloor n/2 \rfloor$.

Theorem 13. *Let $m \geq 159$. Then a.a.s. $G_{m\text{-old}}^n$ has a perfect matching.*

Proof. We are going to use the two-round exposure technique, as discussed in Section 2.3. Recall that the graph $G_{m\text{-old}}^n$ can be seen as a union of two independent graphs: $G_{m_1\text{-old}}^n$ and $G_{m_2\text{-old}}^n$ as long as $m = m_1 + m_2$. In our situation, $m_1 = 120$ and $m_2 = m - m_1 \geq 39$. It follows from Lemma 8, Lemma 9(ii) and Lemma 10(a) that the properties stated there hold a.a.s. for $G_{m_1\text{-old}}^n$ with $\ell = 1$, $\alpha = 0.0538$ and $\gamma = \gamma(m_1) \leq 0.06238$. In particular, as we aim for an a.a.s. statement, we may expose the edges of $G_{m_1\text{-old}}^n$ first and assume that every set of vertices $K \subseteq [n]$ with $1 \leq |K| \leq \alpha n$ satisfies $|N(K)| \geq |K|$, and that $G_{m_1\text{-old}}^n$ has a matching that isolates at most γn vertices. It is an immediate but crucial observation that these two assumptions will still hold if we add extra edges to $G_{m_1\text{-old}}^n$.

Now, we will repeatedly expose some edges of $G_{m_2\text{-old}}^n$ (in some specific order) and add them to $G_{m_1\text{-old}}^n$. At each stage of the process, let G denote the current graph, consisting of $G_{m_1\text{-old}}^n$ together with the exposed edges from $G_{m_2\text{-old}}^n$. We will argue that, if G does not have a perfect matching, then there is a good chance that exposing more edges of $G_{m_2\text{-old}}^n$ increases the size of a maximum matching by 1. Consider $A = A(G)$ and suppose that G has no perfect matching. As all small sets expand well in $G \supseteq G_{m_1\text{-old}}^n$, it follows from Lemma 12 that $|B(v)| \geq \alpha n$ for any $v \in A$ (note that $B(v) \neq \emptyset$, as no perfect matching is found yet). Therefore, recalling that $A \supseteq B(v)$, we conclude that $|A| \geq \alpha n$.

Suppose that we are at the first stage of the process (that is, when $G = G_{m_1\text{-old}}^n$), and let us focus on the youngest vertex v in A . We expose all edges of $G_{m_2\text{-old}}^n$ from v to older vertices and add them to G . Note that if we find some edge between v and $B(v)$, then we increase the size of a maximum matching by one. As all vertices in $B(v)$ are older than v , the probability of increasing the size of a maximum matching is at least $1 - (1 - \alpha)^{m_2}$ (see (6)). We update G (including all exposed edges from $G_{m_2\text{-old}}^n$) and also the set $A = A(G)$, and continue the process moving to another vertex v , the youngest vertex in A that is not exposed yet with respect to $G_{m_2\text{-old}}^n$. As before $|B(v)| \geq \alpha n$ but, perhaps, one vertex from $B(v)$ (the one that is already exposed) is younger than v . Hence, the probability of increasing the size of a maximum matching is at least $1 - (1 - (\alpha n - 1)/n)^{m_2}$. In general, as long as we expose edges from t vertices and a perfect matching is not found, the probability of extending a maximum matching after exposing another vertex is at least $1 - (1 - (\alpha n - t)/n)^{m_2}$.

It remains to estimate the probability that the process exposes edges of $\lfloor \alpha n \rfloor$ vertices and no perfect matching is found. If this happens, then the random variable X that counts the number of times the process extends some maximum matching is smaller than $\gamma n/2$. Note that X is stochastically bounded from below by a random variable Y that is a sum of independent Bernoulli random variables with parameter $1 - (1 - (\alpha n - t)/n)^{m_2}$, for $t \in \{0, 1, \dots, \lfloor \alpha n \rfloor - 1\}$. Using the Euler–Maclaurin formula and the change of variable $x = t/n$, we can approximate $\mathbf{E}Y$ by an integral and obtain

$$\begin{aligned} \frac{\mathbf{E}[Y]}{n} &\sim \int_0^\alpha \left(1 - (1 - (\alpha - x))^{m_2}\right) dx \\ &= \left[x - \frac{(1 - \alpha + x)^{m_2+1}}{m_2 + 1} \right]_0^\alpha = \alpha - \frac{1}{m_2 + 1} + \frac{(1 - \alpha)^{m_2+1}}{m_2 + 1} > \frac{\gamma}{2}. \end{aligned} \quad (32)$$

It follows from the generalized Chernoff bound that a.a.s. the process does not fail and a perfect matching is found. \square

Adjusting the result to G_m^n is straightforward, given all the tools that we developed in Section 2.4.

Corollary 14. *Let $m \geq 1,260$. Then a.a.s. G_m^n has a perfect matching.*

Proof. The argument is analogous to that of Theorem 13, so we will only sketch the differences. Let $m_1 = 500$, and $m_2 = m - 2m_1 \geq 260$. It follows from Lemma 8, Lemma 9(ii) and Lemma 11(c) that the properties stated there hold a.a.s. for $\pi_1(G_{2m_1, m_2}^m)$ with $\ell = 1$, $\alpha = 0.016801$ and $\gamma = \gamma(m_1) \leq 0.019675$. In particular, $\pi_1(G_{2m_1, m_2}^m)$ has the desired properties as we claimed before for $G_{120\text{-old}}^m$ in the proof of Theorem 13. The second part of the argument has to be adjusted slightly, using Lemma 2 instead of (6). This time, X is stochastically bounded from below by a random variable Y that is a sum of independent Bernoulli random variables with parameter $1 - (1 - (\alpha n - t)/(2n))^{m_2}$, $t \in \{0, 1, \dots, \lfloor \alpha n \rfloor - 1\}$. Now

$$\begin{aligned} \frac{\mathbf{E}[Y]}{n} &\sim \int_0^\alpha \left(1 - (1 - (\alpha - x)/2)^{m_2}\right) dx = \left[x - \frac{2(1 - \alpha/2 + x/2)^{m_2+1}}{m_2 + 1} \right]_0^\alpha \\ &= \alpha - \frac{2}{m_2 + 1} + \frac{2(1 - \alpha/2)^{m_2+1}}{m_2 + 1} > \frac{\gamma}{2}, \end{aligned} \tag{33}$$

and the rest of the proof continues unaltered. \square

3.3 Hamiltonian cycles

In this subsection, we deal with Hamiltonian cycles. The argument is very similar to the one we used for perfect matchings so we skip some details. As before, we start with a deterministic result that can be found, for example, in [21, Corollary 6.7]. This approach was an important breakthrough in finding Hamiltonian cycles in random graphs and came with the result of Pósa [32]. Let $G = (V, E)$ be any graph. Suppose that $P = (a, \dots, x, y, y', \dots, b', b)$ is a path and bx is an edge where $x \neq b$ is either a or an interior vertex of P . Then, the path $P' = (a, \dots, x, b, b', \dots, y', y)$ is said to be obtained from P by a *rotation* with vertex a fixed. Now, for a given path P with a as one of its endpoints, let $\text{END}(P, a)$ be the set of vertices v such that there exists a path from a to v that is obtained from P by a sequence of rotations with vertex a fixed. Then the following lemma holds. (And again we include the proof for completeness.)

Lemma 15 ([21]). *Let $G = (V, E)$ be a graph, P be any longest path of G , and a one of its endpoints. Then,*

$$|N(\text{END}(P, a))| < 2|\text{END}(P, a)|.$$

Proof. We claim that if $v \in V(P) \setminus \text{END}(P, a)$ and v is adjacent to some vertex w in $\text{END}(P, a)$, then there is $w' \in \text{END}(P, a)$ such that $vw' \in E(P)$. Indeed, consider the path P_w witnessing $w \in \text{END}(P, a)$. Consider x with $vx \in E(P)$; if $vx \notin E(P_w)$, then one of the rotations yielding P_w from P deleted the edge vx and hence $x \in \text{END}(P, a)$, so $w' = x$ satisfies the claim. Otherwise $N_P(v) = N_{P_w}(v)$, and performing a rotation with the edge vw in P_w shows that one of the vertices of $N_{P_w}(v)$ is in $\text{END}(P, a)$.

Now since P is a longest path, $N_G(\text{END}(P, a)) \subseteq V(P)$. Our claim implies that furthermore $N_G(\text{END}(P, a)) \subseteq N_P(\text{END}(P, a))$. Each vertex of $\text{END}(P, a)$ has at most two neighbours along P , but also the endpoint of P (other than a) has only one; hence $|N_P(\text{END}(P, a))| < 2|\text{END}(P, a)|$. \square

Suppose that G is connected but has no Hamiltonian cycle. Let $A = A(G) \subseteq V$ be the set of vertices that are endpoints of some longest path. (In particular, the length of a longest path could be $n = |V|$.) For $v \in A$, let

$$B(v) = \{w \in V \setminus \{v\} : w \in \text{END}(P, v) \text{ for some longest path } P \text{ of } G \text{ having } v \text{ as an endpoint}\}.$$

This time, the set $B(v)$ is important for understanding whether a few additional random edges, with v as one of their endpoints, have a chance to increase the length of a longest path (if there is no Hamiltonian path) or create a Hamiltonian cycle (if there is a Hamiltonian path). Indeed, if there is a Hamiltonian path, then adding an edge between v and some vertex in $B(v)$ creates a Hamiltonian cycle. Otherwise, some longest path of length $k < |V|$ creates a cycle of length k . Since G is connected, there is at least one edge joining a vertex from the cycle with some vertex outside and so a longer path can be created.

Now, we are ready to prove the main result for Hamiltonian cycles. Since the proof is very similar to the one of Theorem 13, we simply sketch it.

Theorem 16. *Let $m \geq 3,214$. Then a.a.s. $G_{m\text{-old}}^n$ has a Hamiltonian cycle.*

Proof. Again, we are going to use the two-round exposure: $G_{m\text{-old}}^n = G_{m_1\text{-old}}^n \cup G_{m_2\text{-old}}^n$, with $m_1 = 2,900$ and $m_2 = m - m_1 \geq 314$. It follows from Lemma 8, Lemma 9(i) and Lemma 10(b) that the properties stated there hold a.a.s. for $G_{m_1\text{-old}}^n$ with $\ell = 2$, $\alpha = 0.032003$ and $\beta = \beta(m_1) \leq 0.014414$. In particular, we may expose $G_{m_1\text{-old}}^n$ and assume that every set of vertices $K \subseteq [n]$ with $1 \leq |K| \leq \alpha n$ satisfies $|N(K)| \geq 2|K|$, and that $G_{m_1\text{-old}}^n$ has a path of length at least $n - 2\lceil \beta n \rceil$. Moreover, we assume that $G_{m_1\text{-old}}^n$ is connected by [9].

As in the proof of Theorem 13, we will sequentially pick vertices and expose the edges of $G_{m_2\text{-old}}^n$ that connect them to older vertices. During this process, we try to extend the length of a longest path and, once a Hamiltonian path is created, to close the desired cycle. Each time, if the current graph does not have a Hamiltonian cycle yet, we update the set A and all $B(v)$'s. As all small sets expand well, it follows from Lemma 15 that $|B(v)| \geq \alpha n$ for every $v \in A$, and so also $|A| \geq \alpha n$. At each step, we choose the youngest vertex v in A that has not been picked yet, and expose the m_2 edges of $G_{m_2\text{-old}}^n$ that go from v to older vertices. As before, as long as we have only exposed edges from t vertices and a Hamiltonian cycle has not been found, the probability of improving our current situation in the next step is at least $1 - (1 - (\alpha n - t)/n)^{m_2}$. From an analogous computation to the one leading to (32), the expected number of successful steps in the process is at least $(1 + o(1))cn$, where

$$c := \alpha - \frac{1}{m_2 + 1} + \frac{(1 - \alpha)^{m_2 + 1}}{m_2 + 1} > 2\beta.$$

Since it only takes $2\lceil \beta n \rceil \sim 2\beta n$ (or less) successful steps to discover a Hamiltonian cycle, we use the generalized Chernoff bound as before to conclude that a.a.s. a Hamiltonian cycle is found during the process. \square

Finally, we will extend the result to G_m^n for the last time, by adjusting the proof of Theorem 16 in the same spirit as Corollary 14 was obtained from Theorem 13. We omit details, and just state the main differences from the previous arguments.

Corollary 17. *Let $m \geq 29,500$. Then a.a.s. G_m^n has a Hamiltonian cycle.*

Proof. Let $m_1 = 14,000$, and $m_2 = m - 2m_1 \geq 1,500$. Take $\ell = 2$ and $\alpha = 0.008874$, and note that $\beta = \beta(m_1) \leq 0.003760$. Then, by Lemma 8, Lemma 9 and Lemma 11(d), $\pi_1(G_{2m_1, m_2}^n)$ has the same desired properties as we claimed for $G_{2,900\text{-old}}^n$ in the proof of Theorem 16. After adjusting (33) and verifying that

$$\alpha - \frac{2}{m_2 + 1} + \frac{2(1 - \alpha/2)^{m_2 + 1}}{m_2 + 1} > 2\beta,$$

the proof is finished. \square

4 Lower bound

In [11], a simplified model for crawling complex networks such as the web graph was proposed, which is a variation of the robot vacuum edge-cleaning process introduced earlier by Messinger and Nowakowski [28]. In particular, it was shown that a.s. the preferential attachment model G_2^n has the following property: the optimal robot crawler needs substantially more than n steps to visit all the vertices of the graph (see Theorem 10 in [11] and the discussion preceding it). This immediately implies (in a strong sense) that G_2^n is not Hamiltonian a.s. Moreover, these observations, combined with an additional argument presented below, show a stronger property, namely, that G_2^n has no perfect matching a.s.

Here we briefly recall the the main ingredients of the argument in [11] without reference to the robot crawling process. Fix some constant $c \in (0, 1)$ and consider G_2^n . We say that a vertex t is *old* if $t \leq cn$ (and *young* otherwise). Moreover, a vertex is called *lonely* if it is never chosen as the recipient of an edge stemming from another vertex added later on in the process. (In particular, a lonely vertex must have degree 2.) Let A_n be the number of vertices in G_2^n that are young, lonely, with two old neighbours; let B_n be the number of vertices that are young, lonely, with exactly one old neighbour; let C_n be the number of vertices that are old and lonely; and finally, let D_n count the number of vertices that are old but not lonely. Pick any young vertex $t > cn$. It follows from Lemma 5 that the probability that exactly i edges stemming from t are attached to old vertices is asymptotic to

$$\binom{2}{i} \sqrt{cn/t}^i (1 - \sqrt{cn/t})^{2-i}.$$

Moreover, the probability that t is lonely is equal to

$$\prod_{i=t+1}^n \left(1 - \frac{2}{4i} + O(i^{-2})\right)^2 = \exp\left(-\sum_{i=t+1}^n i^{-1} + O(t^{-1})\right) \sim t/n,$$

and this is still valid conditional on any event that involves only edges stemming from vertices in $[t]$. In particular, the probability that t is lonely and has exactly one old neighbour is asymptotic to $2\sqrt{cn/t}(1 - \sqrt{cn/t})t/n$, and thus

$$\mathbf{E}B_n \sim n \int_c^1 2\sqrt{c/x}(1 - \sqrt{c/x})x \, dx = n \left(\frac{4\sqrt{c}}{3} - 2c + \frac{2c^2}{3}\right).$$

A similar calculation gives that $\mathbf{E}(B_n(B_n - 1)) \sim (\mathbf{E}B_n)^2$ (since, for each pair of different young vertices t_1 and t_2 , the events $\{t_1 \text{ contributes to } B_n\}$ and $\{t_2 \text{ contributes to } B_n\}$ are asymptotically independent). This implies that $\mathbf{Var}B_n = o((\mathbf{E}B_n)^2)$, and hence by Chebyshev inequality we get that a.s. $B_n \sim \mathbf{E}B_n$. Proceeding analogously for A_n , C_n and D_n , we conclude that a.s.:

$$\begin{aligned} A_n &\sim n \int_c^1 (\sqrt{c/x})^2 x \, dx = n(c - c^2), \\ B_n &\sim n \int_c^1 2\sqrt{c/x}(1 - \sqrt{c/x})x \, dx = n \left(\frac{4\sqrt{c}}{3} - 2c + \frac{2c^2}{3}\right), \\ C_n &\sim n \int_0^c x \, dx = n\frac{c^2}{2}, \\ D_n &\sim n \left(c - \frac{c^2}{2}\right). \end{aligned}$$

Consider the graph H (on $n - D_n$ vertices) obtained by deleting the old but not lonely vertices. The C_n old and lonely vertices as well as the A_n young and lonely vertices with two old neighbours are isolated in H , and the B_n young and lonely vertices with one old neighbour have degree 1 in H . Since

$$\frac{B_n}{2n} + \frac{A_n}{n} + \frac{C_n}{n} \sim \frac{1}{2} \left(\frac{4\sqrt{c}}{3} - 2c + \frac{2c^2}{3} \right) + c - c^2 + \frac{c^2}{2} > c - \frac{c^2}{2} \sim \frac{D_n}{n}$$

for any choice of $c \in (0, 1)$, the graph a.a.s. does not have a Hamiltonian cycle. In fact, this argument *almost* shows that a.a.s. there is no perfect matching. Indeed, the number of odd components in H is at least the number of isolated vertices in H , which in turn is a.a.s. equal to $A_n + C_n \sim (c - c^2/2)n$, and so it coincides with an asymptotic value of D_n , the number of vertices removed from G_2^n to get H . Hence, it *almost* violates Tutte's condition for existence of a perfect matching.

In order to show that the necessary (and sufficient) Tutte's condition does not hold for, say, $c = 1/4$, we are going to show that a.a.s. there are $\Omega(n)$ components of size 3 in H (and so there are more odd components than vertices removed). As the constant hidden in the $\Omega()$ notation does not affect the argument, we are not going to optimize it here. Let us partition the vertex set $V = [n]$ into 4 subsets $V_1 = \lfloor [n/4] \rfloor$, $V_2 = \lfloor [n/2] \rfloor \setminus \lfloor [n/4] \rfloor$, $V_3 = \lfloor [3n/4] \rfloor \setminus \lfloor [n/2] \rfloor$, and $V_4 = [n] \setminus \lfloor [3n/4] \rfloor$. (In particular, V_1 is the set of old vertices.) We say that a triple (v_2, v_3, v_4) forms a *cherry* if $v_2 \in V_2$, $v_3 \in V_3$ are the two older neighbours of a lonely vertex $v_4 \in V_4$, and v_4 is the only younger neighbour of each of v_2 and v_3 . Moreover, we say that a cherry (v_2, v_3, v_4) is *sweet* if v_2 and v_3 both have two neighbours in V_1 (these neighbours are older than them; the third, common and younger, neighbour is v_4 , of course). Let S_n denote the number of sweet cherries.

We proceed to bound from below the probability that a given tuple (v_2, v_3, v_4) forms a sweet cherry. First note that the probability that v_2 has two neighbours in V_1 is at least $(1 + o(1))(1/2)^2$. Conditional upon that, the probability that no vertex i ($v_2 < i < v_3$) is adjacent to v_2 is at least $\prod_{i=v_2+1}^{v_3-1} \left(1 - \frac{2}{4(i-1)}\right)^2$. Conditional upon these two events, the probability that v_3 has two neighbours in V_1 is at least $(1 + o(1))(1/3)^2$. Conditional on all the previous events, the probability that no vertex i ($v_3 < i < v_4$) is adjacent to v_2 or v_3 is at least $\prod_{i=v_3+1}^{v_4-1} \left(1 - \frac{4}{4(i-1)}\right)^2$. Conditional on all the above, the probability that v_4 is adjacent to both v_2 and v_3 is $(2 + o(1)) \left(\frac{2}{4n}\right)^2$. Finally, conditional upon all the previous events, the probability that no vertex i ($v_4 < i \leq n$) is adjacent to v_2 or v_3 is at least $\prod_{i=v_4+1}^n \left(1 - \frac{6}{4(i-1)}\right)^2$. Putting everything together, the probability that (v_2, v_3, v_4) is a sweet cherry is at least

$$(1 + o(1)) 2 \left(\frac{1}{6}\right)^2 \left(\frac{1}{2n}\right)^2 \prod_{i=v_2+1}^n \left(1 - \frac{3}{2(i-1)}\right)^2 = \Omega(1) \left(\frac{1}{n}\right)^2 \left(\frac{v_2}{n}\right)^3 = \Omega(1/n^2).$$

Since the number of possible tuples (v_2, v_3, v_4) is $\Theta(n^3)$, we conclude that $\mathbf{E}S_n = \Omega(n)$. Moreover, given two different tuples (v_2, v_3, v_4) and (v'_2, v'_3, v'_4) , easy calculations show that the events $\{(v_2, v_3, v_4) \text{ is a sweet cherry}\}$ and $\{(v'_2, v'_3, v'_4) \text{ is a sweet cherry}\}$ are asymptotically independent for $v_2 \neq v'_2$, $v_3 \neq v'_3$ and $v_4 \neq v'_4$ or disjoint otherwise. Then, we can argue as we did before for B_n to conclude that $\mathbf{Var}S_n = o((\mathbf{E}S_n)^2)$ and a.a.s. $S_n \sim \mathbf{E}S_n$. As a result, there are a.a.s. $\Omega(n)$ sweet cherries (inducing components of size 3 in H) and so G_2^n a.a.s. has no perfect matching.

Unfortunately, we do not know how to show that a.a.s. $G_{2\text{-old}}^n$ has no perfect matching (perhaps it does have one, as some experiments for small values of n might suggest). On the other hand, showing that a.a.s. it has no Hamiltonian cycle is easy: a.a.s. there are three lonely vertices that have a common neighbour.

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