arm64e

An ABI for Pointer Authentication

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What is arm64e?

- arm64e is an ABI for pointer authentication on ARMv8.3
- ARMv8.3 is an AArch64 extension provided by the Apple A12 and later (e.g. iPhone XR/XS, released September 2018)
- Used for all system software on those devices
- Not ABI stable yet — still looking for ways to strengthen it

What is Pointer Authentication?

- Security mitigation technique
- Provides control flow integrity (CFI), limited data integrity
- Basic idea: sign and authenticate pointers to prevent attackers from escalating memory corruption bugs

Memory Corruption

- Many exploits start with memory corruption bugs
- e.g. buffer overflows, use-after-free
- Ideally, these bugs wouldn’t exist
  - Safe languages, safe practices, static analysis, thorough code review
- Practically, mitigation still has an important place

Exploitation

- Limited memory corruption is not usually the goal of an attack
- Attacker wants to access sensitive information, make specific system calls, exfiltrate data over network, etc.
- Escalating an attack often requires corrupting control flow

Code Payloads

- Attacker wants to run some custom code
- Can’t just write new instructions in modern systems

```
MOV   X0, #0x8              ; first argument: client socket descriptor
MOV   X1,  #0x1F0174ED0     ; second argument: address of password file in memory
MOV   X2, #8096             ; third argument: length
BL    _write
```

Gadgets

- Instead, attacker finds gadgets: bits and pieces of existing functions that collectively do what the attacker wants

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Gadgets

- Instead, attacker finds gadgets: bits and pieces of existing functions that collectively do what the attacker wants

```
_getBitsInByte:
MOV   X0, #0x8                ; return number of bits in a byte
RET

_readPasswordHeader:
MOV   X17, #0x1F0174ED0       ; put address of password file in scratch register
LDR   X0, X17                 ; load from it (leaving address in register)
RET

; next we need a gadget that will move x17 into x1
; etc.
```

ROP/JOP

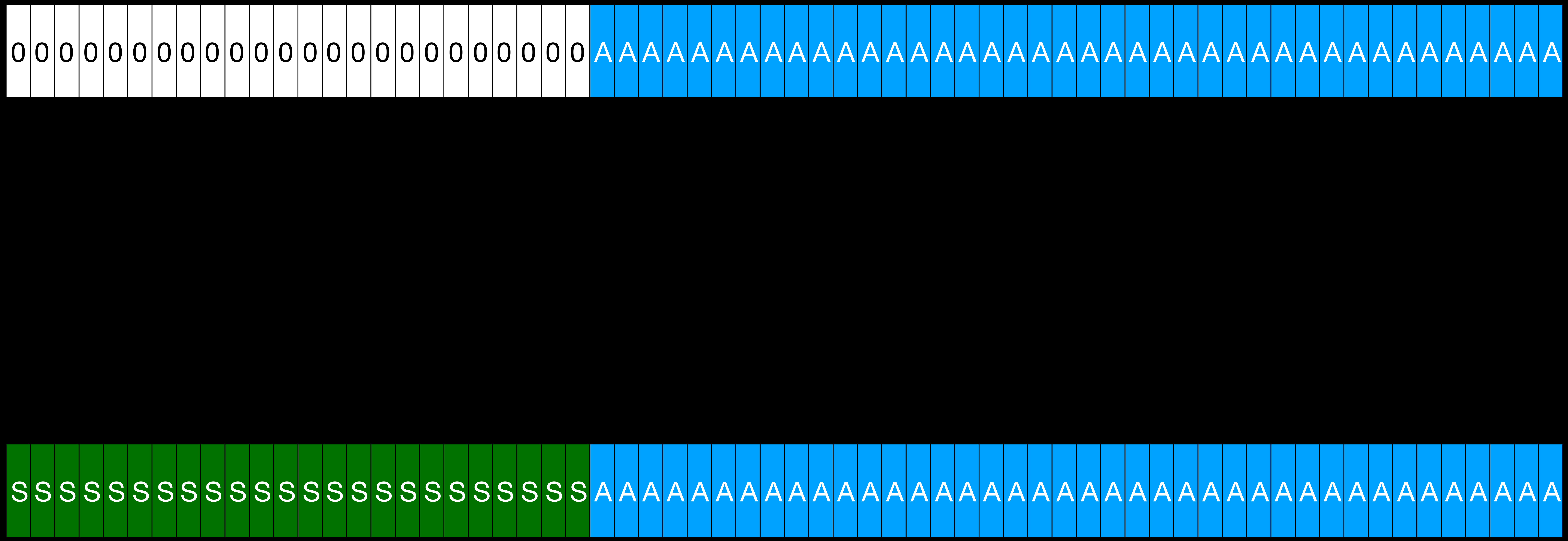
- Attacker must call all of these gadgets in the right sequence
- Use memory corruption to redirect indirect branches to gadgets
  - Redirecting returns: return-oriented programming (ROP)
  - Redirecting calls: jump-oriented programming (JOP)

Pointer Authentication

- Goal: prevent this from working by breaking attempts to redirect
- Add a signature to every code pointer
  - (and some select data pointers)
- Always authenticate signature before doing an indirect branch
  - (and some select loads)

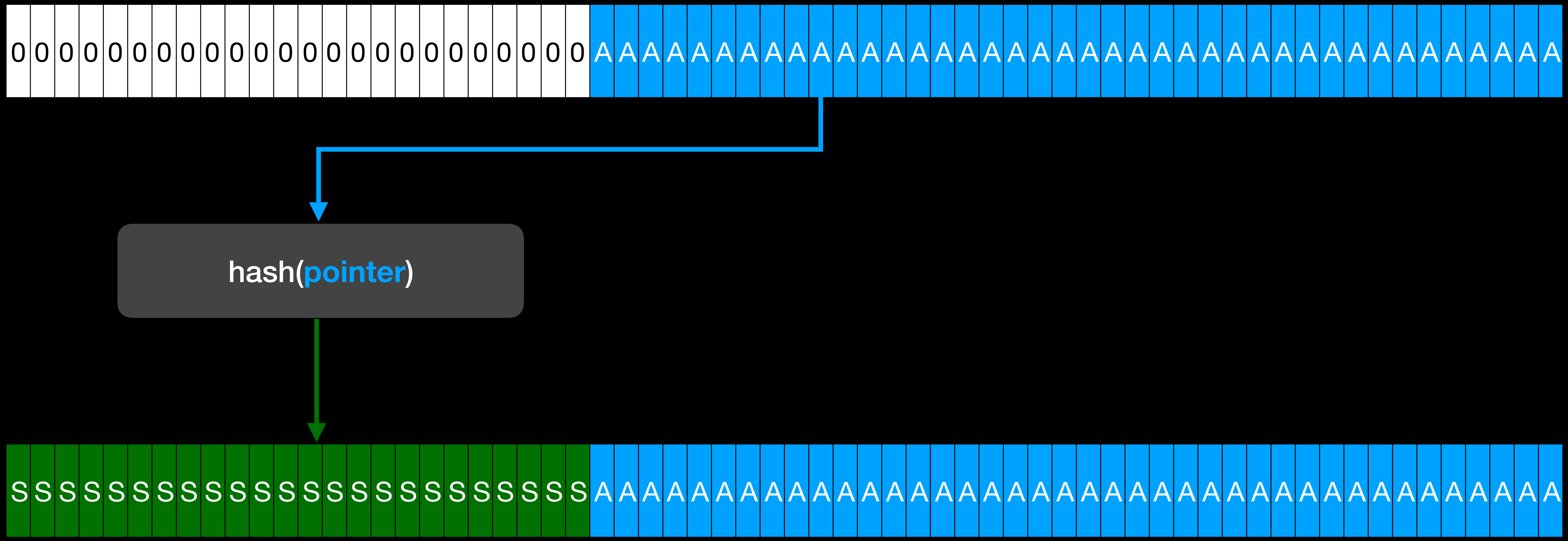
ARMv8.3 Pointer Signatures

- Signature is stored in unused high bits of a 64-bit pointer (~25 bits today)



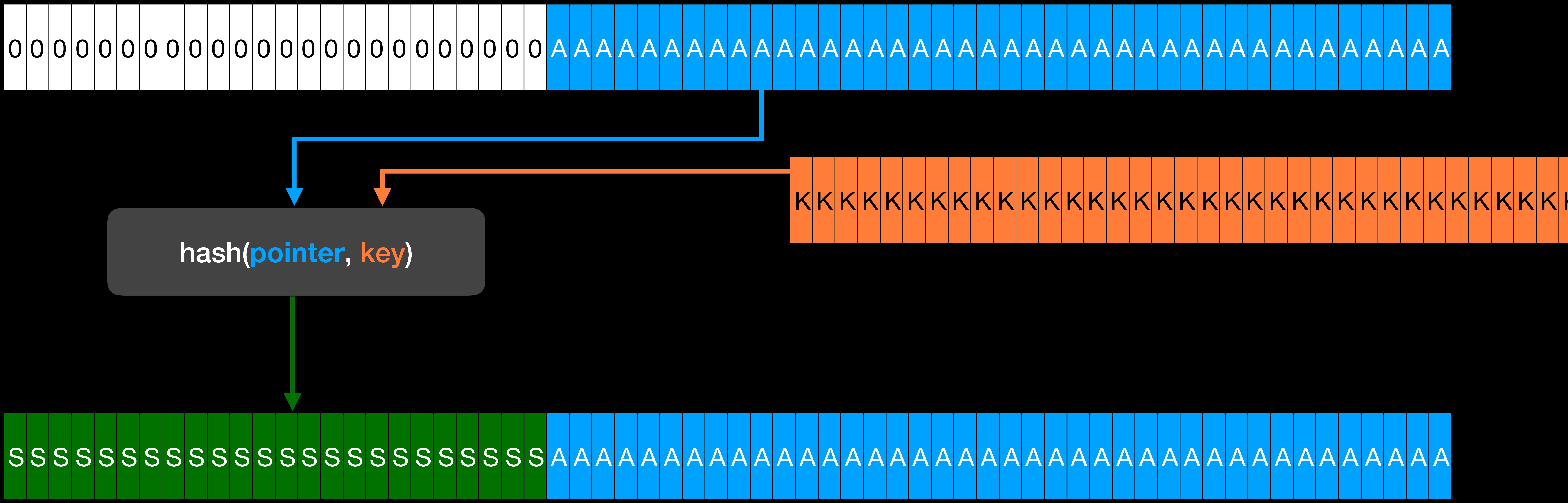
ARMv8.3 Pointer Signatures

- Computed by performing a cryptographic hash of the base pointer



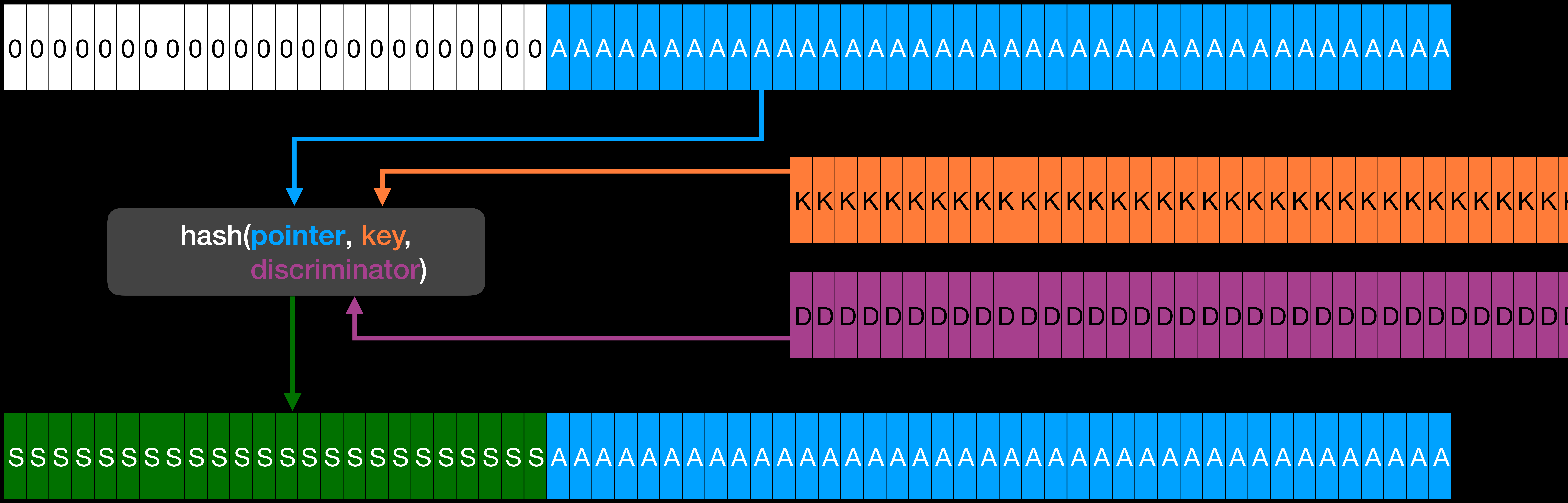
ARMv8.3 Pointer Signatures

- Hash also incorporates data from one of several secret 128-bit key registers, only directly readable by the kernel (a “pepper”)



ARMv8.3 Pointer Signatures

- Hash also incorporates a 64-bit discriminator (a “salt”)



Pointer Substitution

- Signing with secret key means attackers can’t forge signed pointers
- Attackers can still overwrite signed pointers with other signed pointers
- Means gadgets have to be whole functions, but apparently that’s not a serious hurdle

Discriminators

- Substitution only works if all the inputs to the hash are the same

```
hash(pointer, key, discriminator)
```

- Small number of keys, so it mostly comes down to discriminators

Discriminators

- Ideally, every different “purpose” would use a different discriminator
  - A pointer should only authenticate if a human programmer would say that the pointer was meant to be used there
- Pointer authentication mostly driven automatically by compiler
  - Limited by imperfect knowledge
  - Limited by language design

Language ABI

- Compiler automatically protects all indirect branches:
  - `return`
  - `switch`
  - symbol imports (GOT)
  - C function pointers
  - C++ virtual functions
  - etc.
- ABI rule specifies key and how to compute the discriminator

Discriminators in the ABI

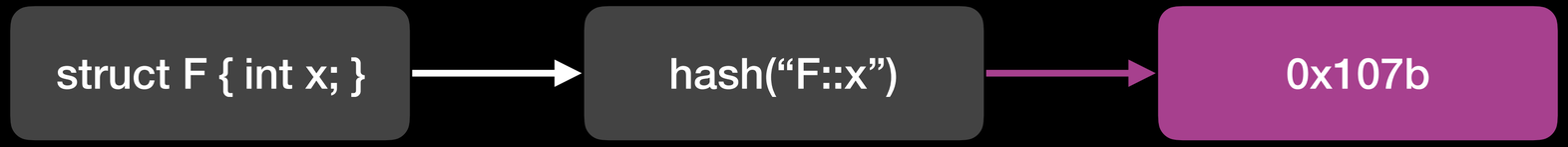
- ARMv8.3 allows discriminators to be arbitrary 64-bit values
- For practical reasons, discriminators used in language ABI are restricted
- Combination of two factors:
  - whether to use address diversity
  - choice of small constant discriminator

Address Diversity

- Discriminator includes storage address of pointer
- Same pointer stored in different places will have a different signature
- Copying requires re-signing, so attackers can’t replace pointers themselves, have to convince the program to do it for them
- Incompatible with `memcpy`, makes copies much more expensive

Constant Discriminators

- 16-bit constant integer
- Can be derived from declaration:



- Can be derived from type:

- Declaration is better, but can’t break abstract, type-based uses

Example: C++ Virtual Functions

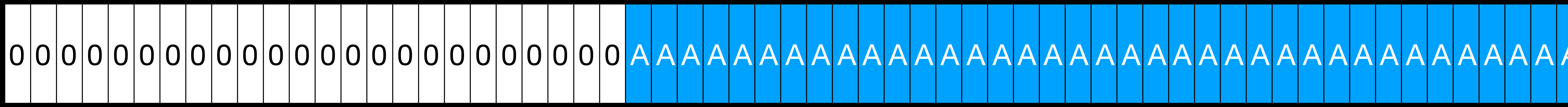
- No direct access to v-table in language, ODR provides strong guarantees
- Can sign virtual function pointers with address diversity
- Can use mangling of method declaration for constant discriminator
- Abstract uses (member function pointers) can be supported without weakening basic ABI
- V-table pointer in object also signed

Example: C Function Pointers

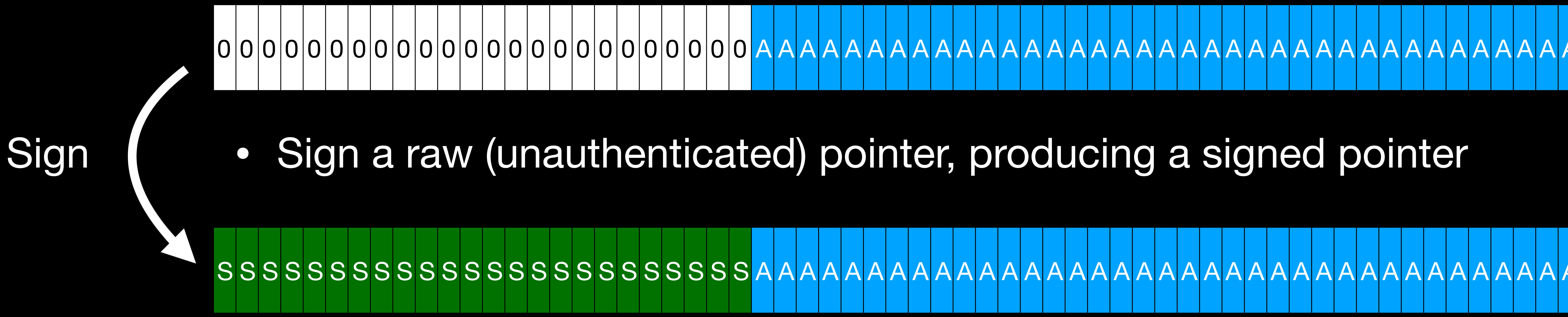
- Pointers must be copyable with `memcpy`, so no address diversity
- Can take address of function-pointer variables, so must use common discriminator for function-pointer type
- Lots of practical deployment challenges with discriminating by type
- Currently using a common discriminator of 0 for all C function pointers
- Clang provides language features to opt in to better discrimination

Generating Code for arm64e

Core Operations

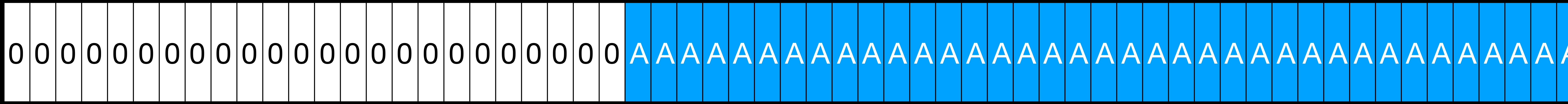


Core Operations



- Sign a raw (unauthenticated) pointer, producing a signed pointer

Core Operations



Sign

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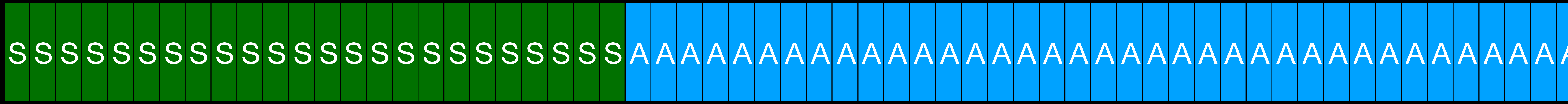
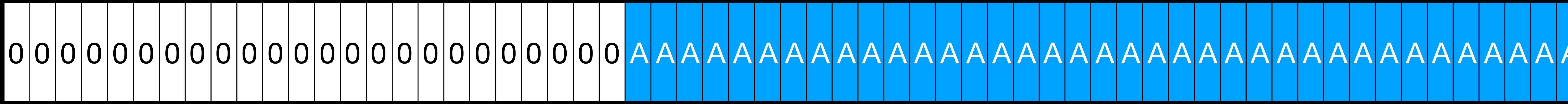


Auth

- Authenticate a signed pointer, producing a raw pointer
  - Verifies the signature, and strips it on success



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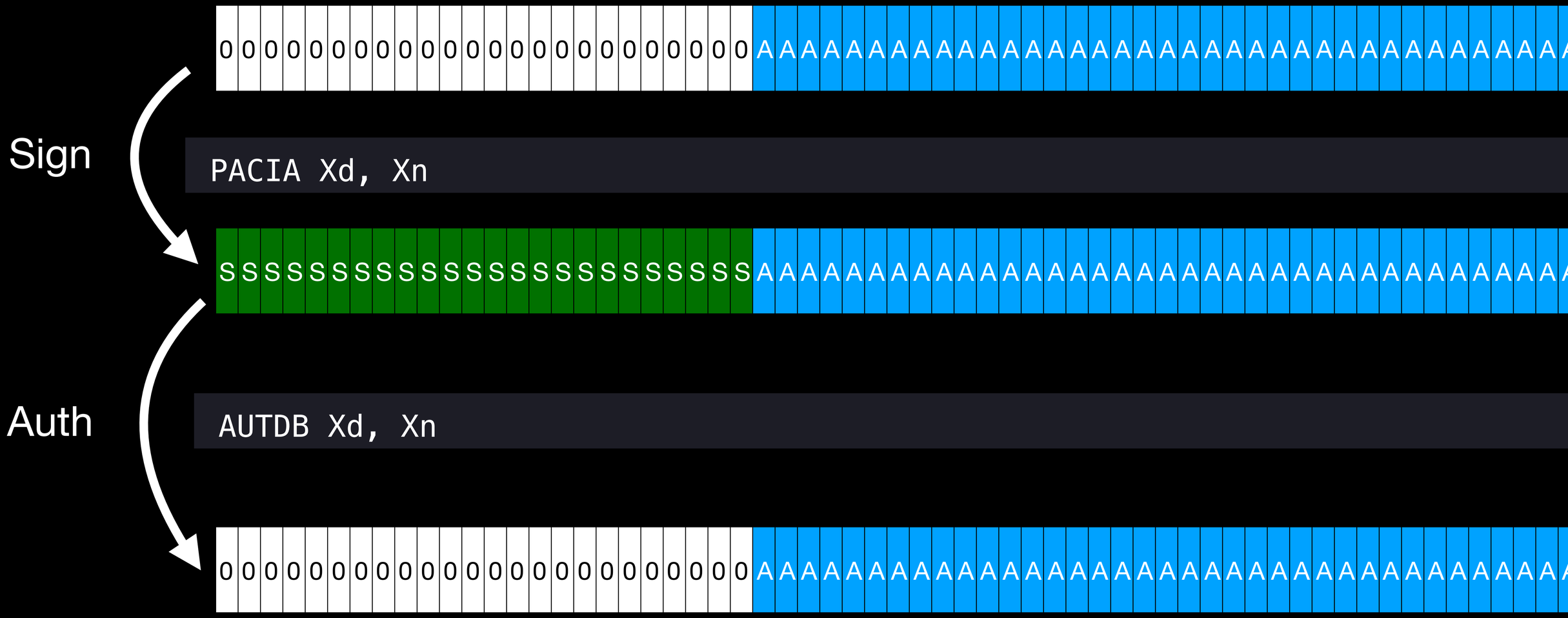
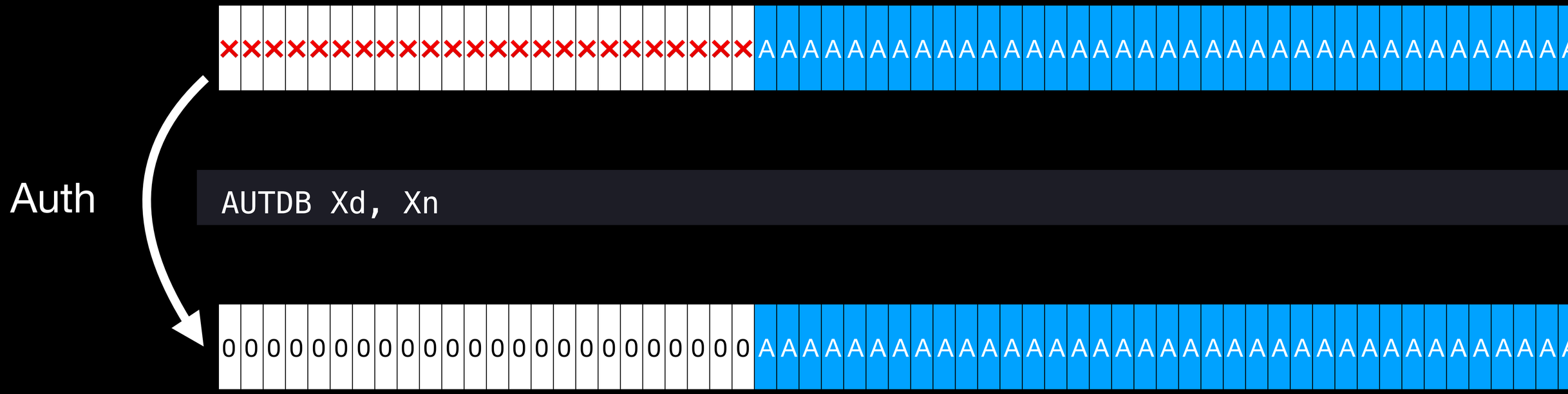


Core Operations

- Pointers with an invalid signature can't be authenticated



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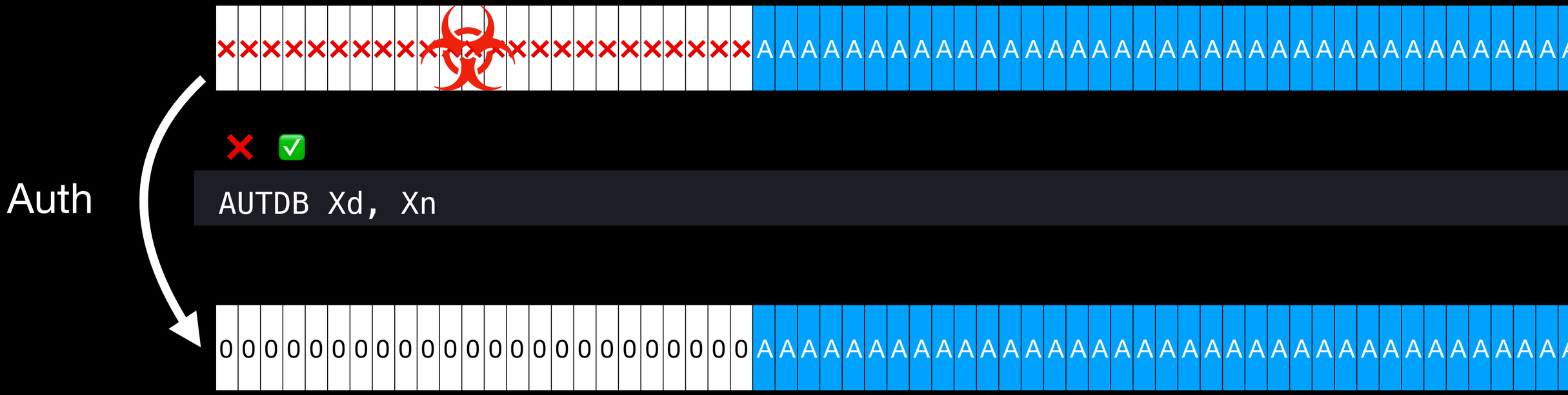

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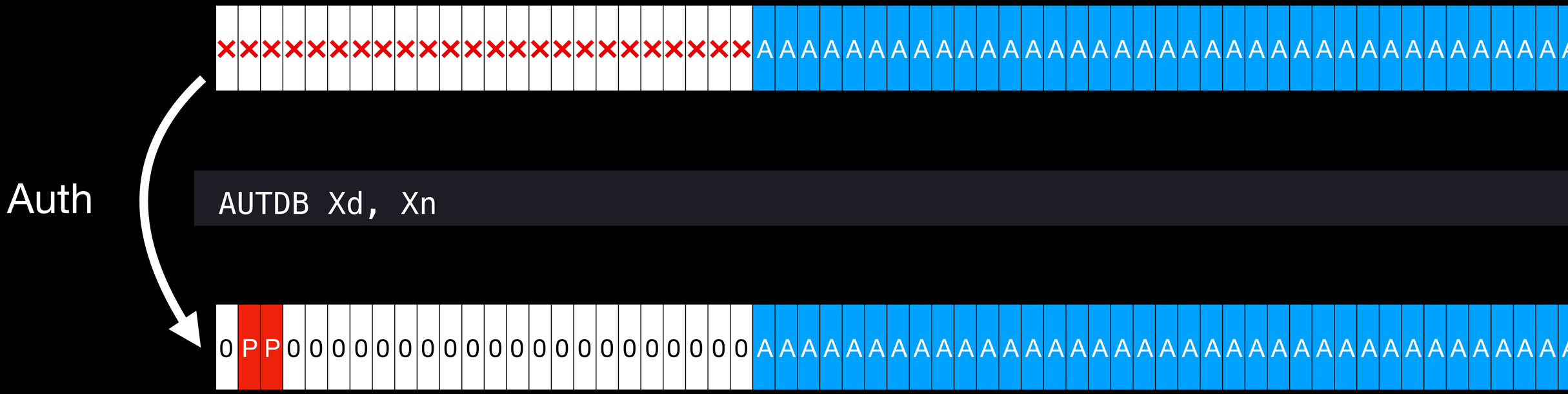

Security Requirements

- Auth: must prevent attackers from bypassing signature verification
- Sign: must prevent attackers from signing pointers they control
- Core operations deal with raw pointers
  - Raw pointers are vulnerable, because they aren't verified
  - Raw pointers shouldn't be exposed (spilled to memory, ...)

Security Guarantees

- It's hard to reason about arbitrary uses
  - No guarantees can be made (e.g., against spilling)
- But we can reason about certain critical, well-defined, uses
  - arm64e mainly uses auth/sign to implement Control Flow Integrity
  - We must guarantee integrity of pointers used in control flow

Important Use-cases

- Authenticate a pointer...
  - ...used as a branch/call target
  - ...that's immediately re-signed
- Sign a pointer...
  - ...to a constant, as a constant initializer
  - ...to a constant, in code

Important Use-cases

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```
(*funptr)();

obj->method();
```

Auth Operand Bundle: Call

- `"ptrauth"` operand bundle on indirect `calls`

```
call void %signed_callee() [ "ptrauth"(i32 0, i64 %disc) ]
```

- Guarantees integrity of the intermediate pointer
- On ARMv8.3, guarantees combined instruction codegen:

```
BLRAAZ Xd
```

Auth Operand Bundle: IndBr

- `indirectbr` is also indirect control flow
  - Let's give it a `"ptrauth"` operand bundle
- Tedious but straightforward patch

Auth Operand Bundle: Switch?

- Jump tables are created late
  - Jump table dispatch only exists in the backend
- We could sign the jump table entries
  - ...would require moving them from text to data
  - ...would prevent shrinking them for small offsets
- Too expensive

Jump Table Hardening

- Jump-table dispatch sequences are hardened using a custom sequence:

```
CMP   Xindex, #<jt size>
CSEL  Xindex, Xindex, XZR, ls          ; range-check the index
; we don't control the index: it could have been spilled across arbitrary blocks
; on index overflow, it's okay to pick any case: it's legitimate control flow
ADRP  Xjt, _JT0@PAGE
ADD   Xjt, _JT0@PAGEOFF                ; materialize the jump table address
LDRSW Xoffset, [Xjt, Xindex, lsl #2]   ; load the offset from the table
ADD   Xtarget, Xjt, Xoffset            ; compute the target
BR    Xtarget                          ; jump to it: no auth, because it's safe
```

Important Use-cases

- Authenticate a pointer...
  - ...used as a branch/call target
  - ...that's immediately re-signed
- Sign a pointer...
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  - ...to a constant, in code

```
void (*p)(char *);

return ((void)(*)(int *)) p;
```

Resign

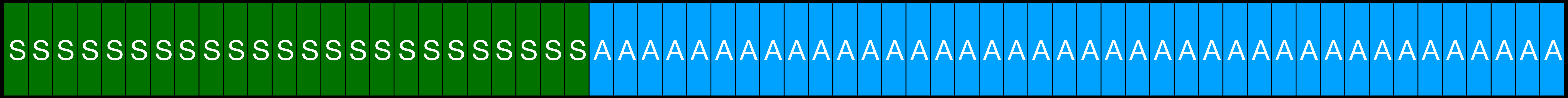
- Authenticate a pointer using key/discriminator A,
  and re-sign it using key/discriminator B

```
declare i64 @llvm.ptrauth.resign.i64(i64, i32, i64, i32, i64)
```

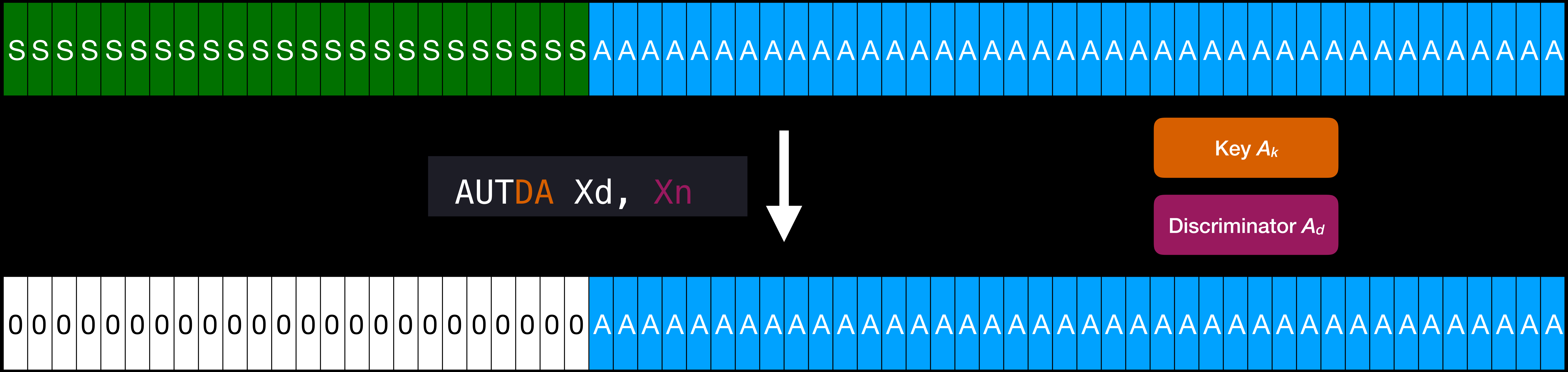
- Guarantees integrity of the intermediate pointer:

```
AUTDA X16, Xn
PACDB X16, Xm
```

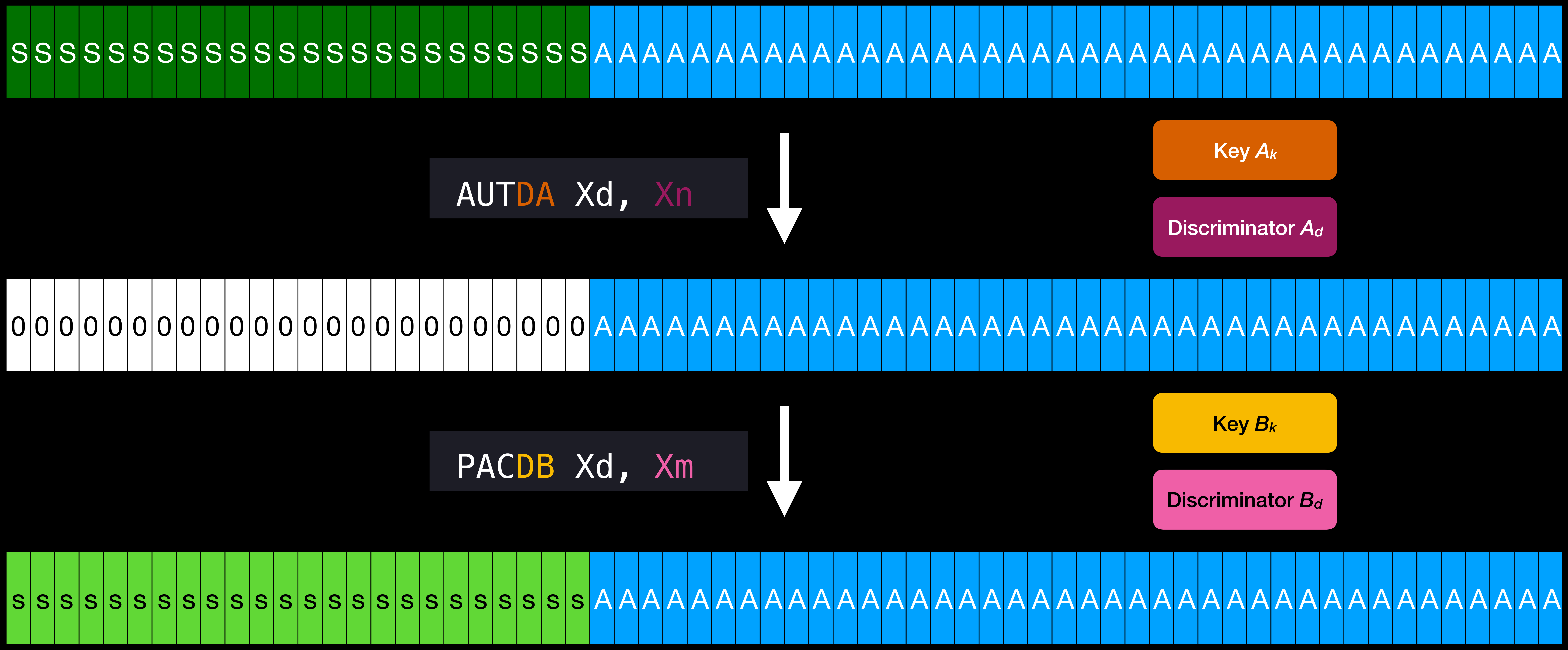
Resign


SSSSSSSSSSSSSSSSSSSSSSSSAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA

Resign


SSSSSSSSSSSSSSSSSSSSSSSSAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
AUTDA Xd, Xn
Key A_k
Discriminator A_d
000000000000000000000000AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA

Resign


SSSSSSSSSSSSSSSSSSSSSSSSAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
AUTDA Xd, Xn
Key A_k
Discriminator A_d
000000000000000000000000AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
PACDB Xd, Xm
Key B_k
Discriminator B_d
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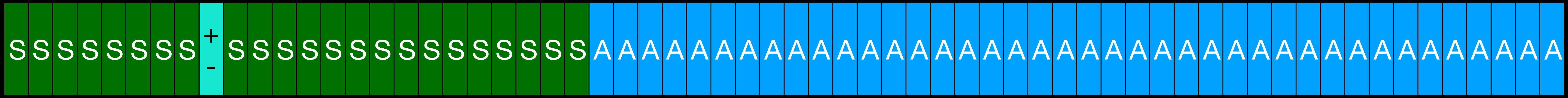
Resign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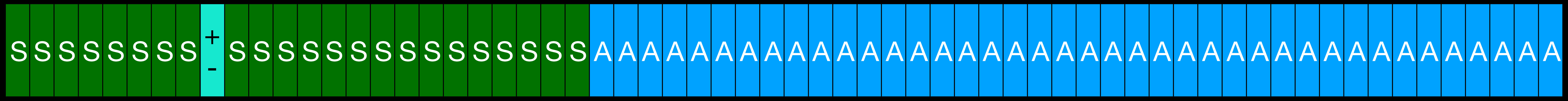
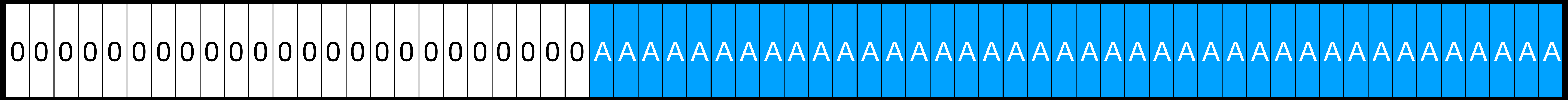



```
AUTDA Xd, Xn
```

- Broadcast a "selector" bit, used for the kernel address-space



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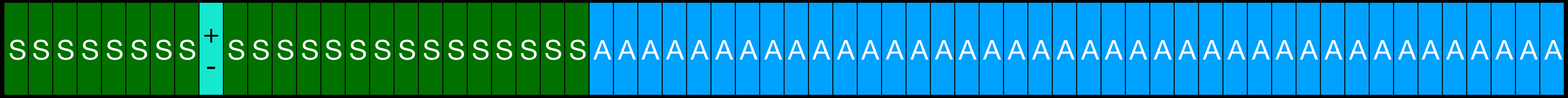


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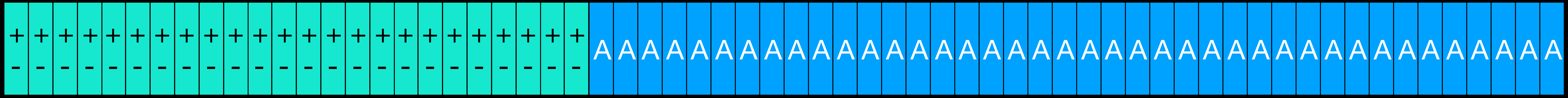

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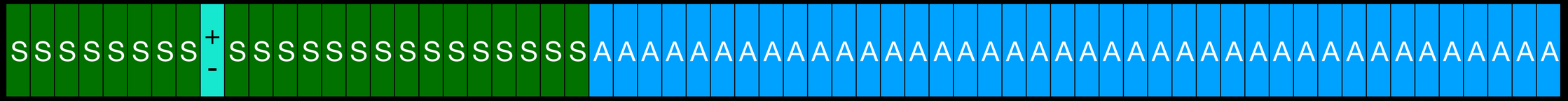

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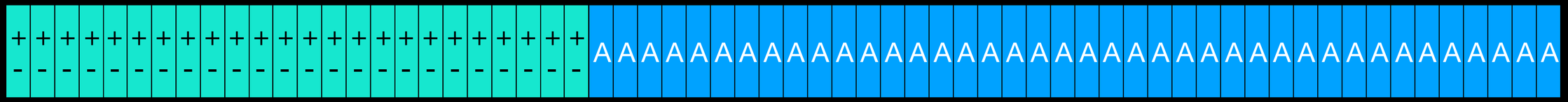
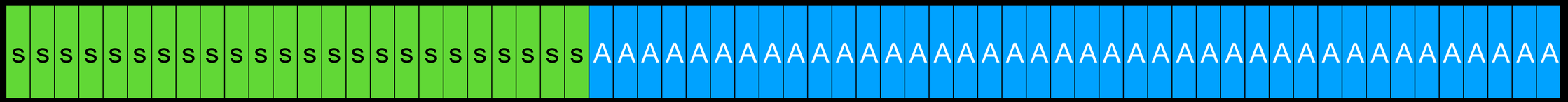


`PACDB Xd, Xm`

- Truncates address-space bits into the selector bit



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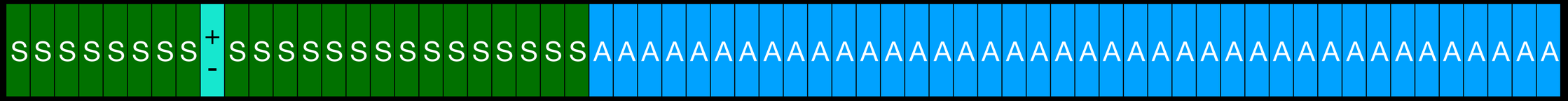

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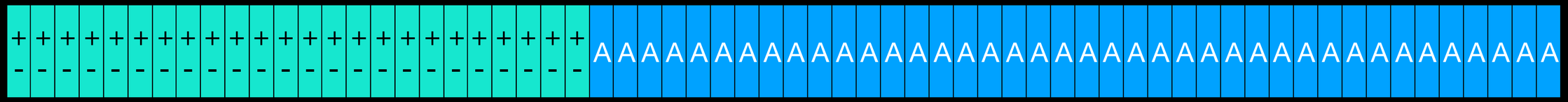

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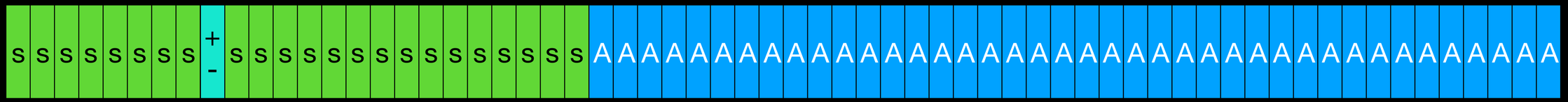
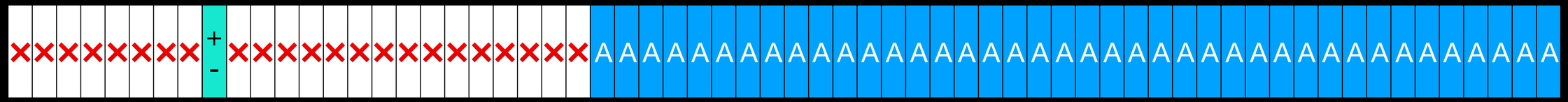


Resign Failures



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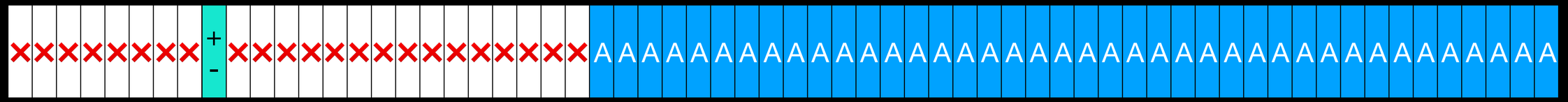
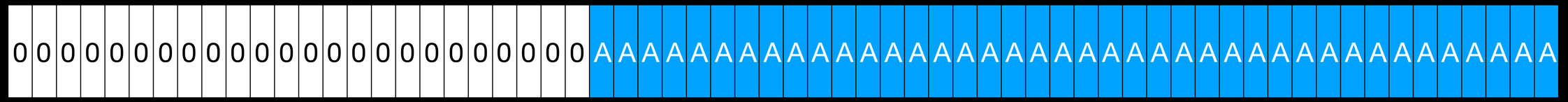



```
AUTDA Xd, Xn
```

- AUT poisons result pointer, because the signature is invalid



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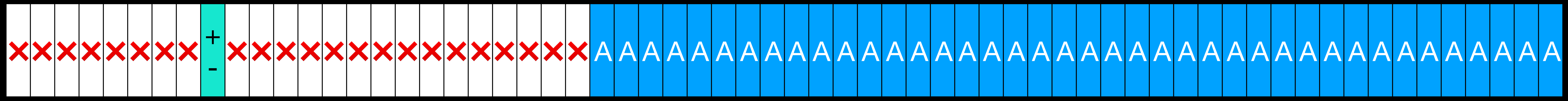

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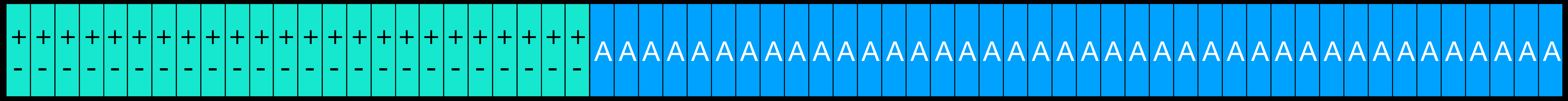

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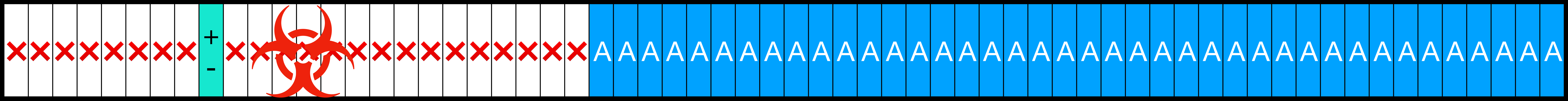

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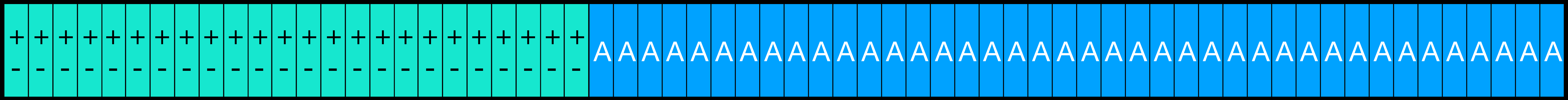

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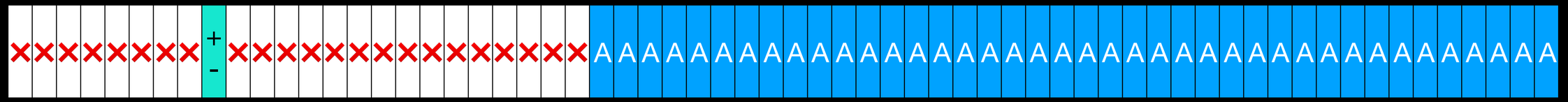


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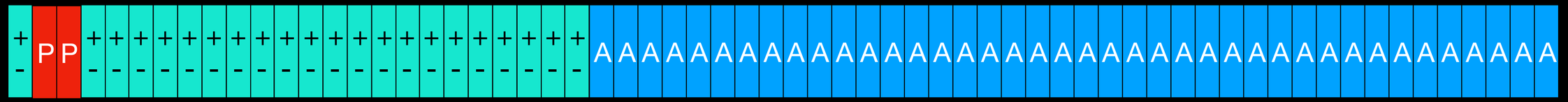

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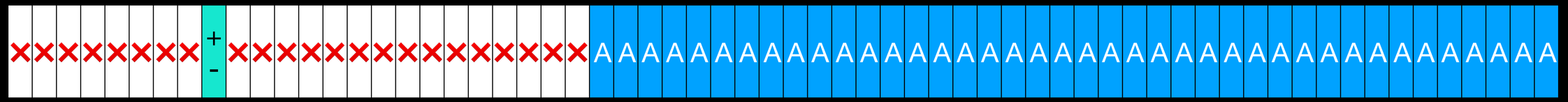


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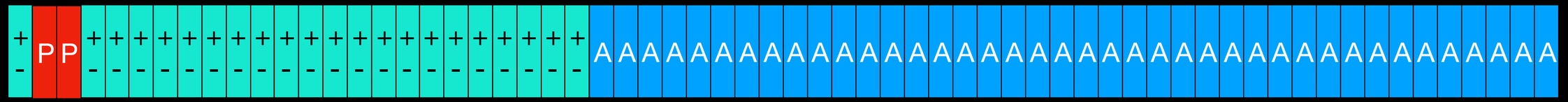
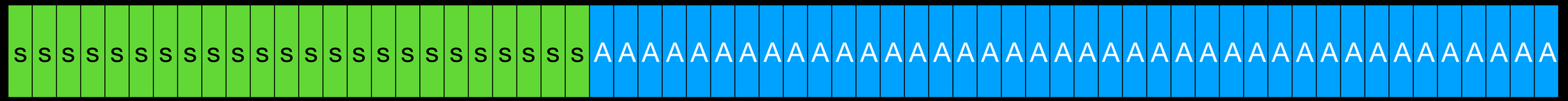



```
PACDB Xd, Xm
```

- PAC corrupts result pointer, because poison bits conflict with addrspace bits



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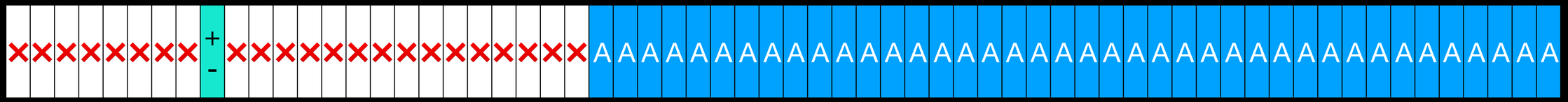


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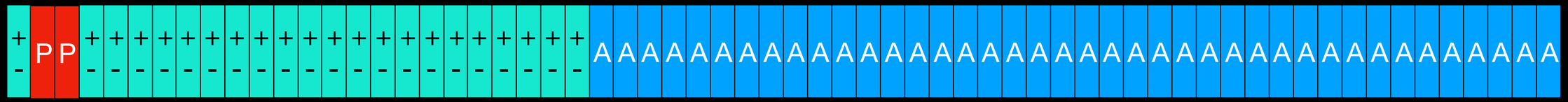
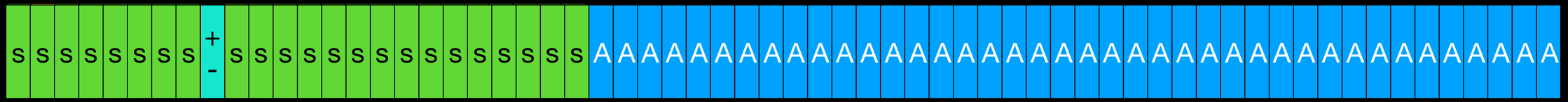
```

- AUT poisons result pointer, because the signature is invalid

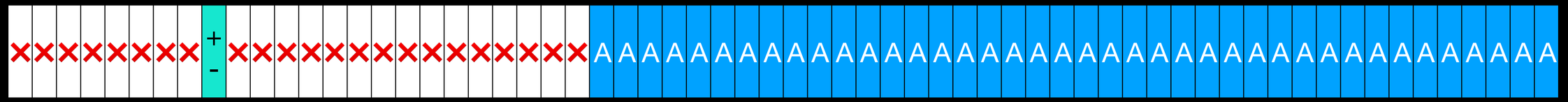



```
PACDB Xd, Xm
```

- PAC corrupts result pointer, because poison bits conflict with addrspace bits

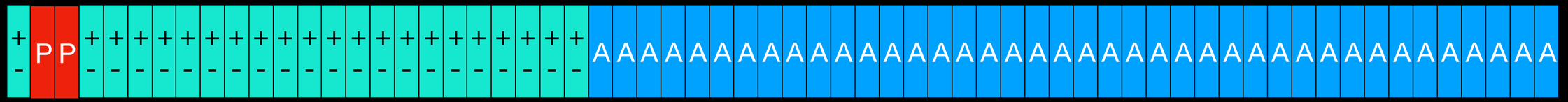


Resign Failures



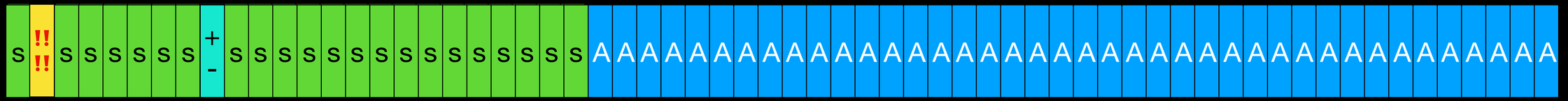

```
AUTDA Xd, Xn
```

- AUT poisons result pointer, because the signature is invalid

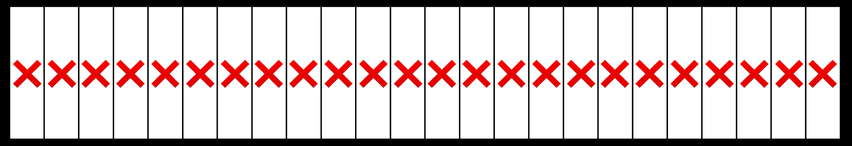



```
PACDB Xd, Xm
```

- PAC corrupts result pointer, because poison bits conflict with addrspace bits

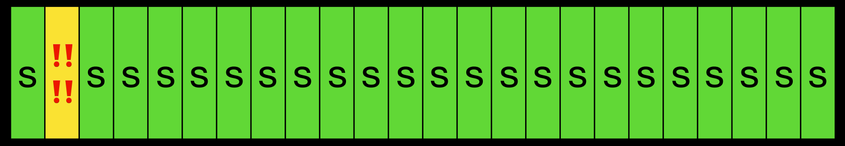


Resign Problem #1



```
AUTDA Xd, Xn
PACDB Xd, Xm
```

- Resigning an invalidly-signed pointer produces a validly-signed pointer!

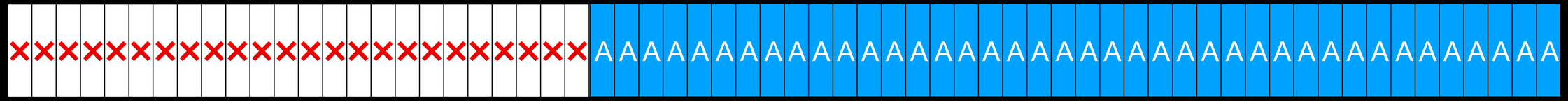


Resign Mitigation #1

- Resign sequence should check for AUT failure
- And return a pointer with no leaked signature bits

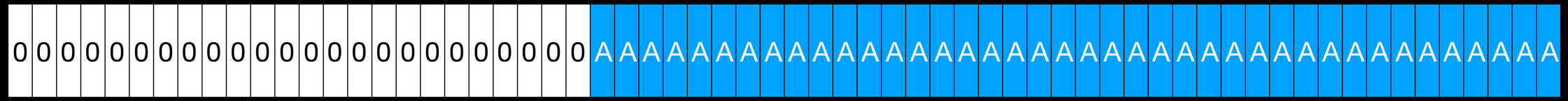
```
MOV   X17, X16              ; We'll need a copy of the pointer
AUTDA X16, X1               ; Authenticate it
XPACD X17                   ; But strip the signature from the copy
CMP   X16, X17              ; Compare the two
PACDB X16, X2               ; Sign the result
CSEL  X16, X16, X17, eq     ; On strip/auth mismatch: return the stripped value
```

Resign Problem #2




```
MOV   X17, X16
AUTDA X16, X1
XPACD X17
CMP   X16, X17
PACDB X16, X2
CSEL  X16, X16, X17, eq
```

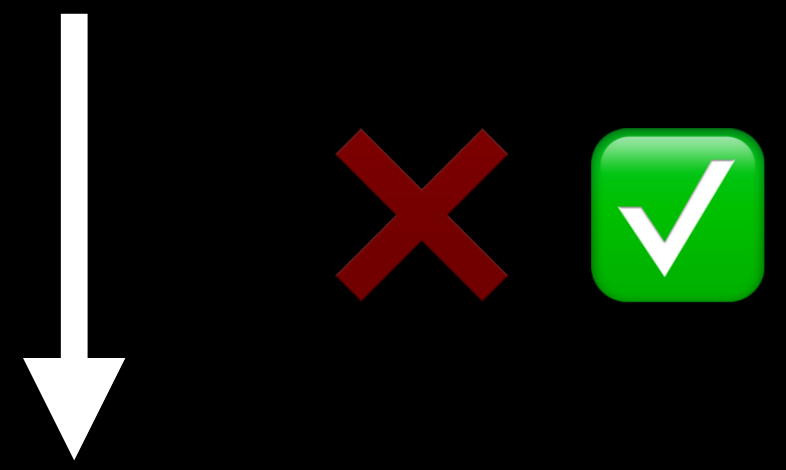
- Checked resign can be bruteforced
- If the result signature bits aren't all 0 (or all 1), the resign succeeded



Resign Problem #2

SSSSSSSSSSSSSSSSSSSSSSSSAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA

```
MOV   X17, X16
AUTDA X16, X1
XPACD X17
CMP   X16, X17
PACDB X16, X2
CSEL  X16, X16, X17, eq
```



- Checked resign can be bruteforced
- If the result signature bits aren't all 0 (or all 1), the resign succeeded

SSSSSSSSSSSSSSSSSSSSSSSSAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA

Resign Mitigation #2

- Resign sequence shouldn't be bruteforceable
  - Not a problem for most auths: the result is (really) used immediately
- Resign sequence should trap on auth failure

```
 MOV   X17, X16              ; We'll need a copy of the pointer
 AUTDA X16, X1               ; Authenticate it
 XPACD X17                   ; But strip the signature from the copy
 CMP   X16, X17              ; Compare the two
 B.EQ  Lsuccess              ; On success, move on
 BRK   #0xc472               ; On mismatch, trap!
Lsuccess:
 PACDB X16, X2               ; Sign the result
```

Important Use-cases

- Authenticate a pointer...
  - ...used as a branch/call target
  - ...that's immediately re-signed
- Sign a pointer...
  - ...to a constant, as a constant initializer
  - ...to a constant, in code

```
typedef void (*fnptr_t)(char *);

fnptr_t actions[] = { &f1, &f2 };
```

Signed Pointer Constant

- `llvm.ptrauth` Authenticated "wrapper" Global (ideally a ConstantExpr)

```
@f.ptrauth = private constant { i8*, i32, i64, i64 }
    { i8* bitcast (i8()* @f to i8*), i32 <key>, i64 <addr disc>, i64 <disc> },
    section "llvm.ptrauth"

@signed_f = constant i8()* bitcast ({ i8*, i32, i64, i64 }* @f.ptrauth to i8()*)
```

- Lowered to a new mach-o relocation:

```
_signed_f:
  .quad _f@AUTH(ia,1234,addr)
```

Important Use-cases

- Authenticate a pointer...
  - ...used as a branch/call target
  - ...that's immediately re-signed
- Sign a pointer...
  - ...to a constant, as a constant initializer
  - ...to a constant, in code

```
void f(char *);

return &f;
```

Signed Pointer Materialization

- `llvm.ptrauth` globals can be used in code too:

```
ret i8()* bitcast ({ i8*, i32, i64, i64 }* @f.ptrauth to i8()*)
```

- Which we lower to:

```
ADRP X16, _f@PAGE
ADD X16, X16, _f@PAGEOFF    ; materialize the pointer, the Darwin way
PACIA X16, Xn               ; sign it
```

Signed Pointer Materialization

```
ADRP X16, _f@PAGE
ADD X16, X16, _f@PAGEOFF     ; materialize the pointer
PACIA X16, Xn                ; sign it
```

- Prevent transforms from exposing the intermediate pointer
  - Backend uses combined ops (PtrAuthGA in ISel, pseudo in AArch64)
- Prevent OS exceptions from exposing the intermediate register value
  - The compiler always uses x16/x17 for "sensitive registers"
  - The kernel guarantees the integrity of x16/x17 on exceptions

Important Use-cases

- Authenticate a pointer...
  - ...used as a branch/call target
  - ...that's immediately re-signed
- Sign a pointer...
  - ...to a constant, as a constant initializer
  - ...to a constant, in code

arm64e

- An ABI for Pointer Authentication
- Extends arm64 language ABIs to provide CFI
  - Discriminator choice is constrained, but is the key to hardening
- Exposes interesting compiler problems
  - Integrity must be preserved throughout all transformations
- Not ABI stable yet — still looking for ways to strengthen it

arm64e

An ABI for Pointer Authentication

LLVM Developers' Meeting
October 22nd, 2019

John McCall
Ahmed Bougacha