

# Linux Security in 10 Years



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# Security Strategy

## ☞ Raise TCO

- Total cost of 'own'ership (Dave Aitel)

## ☞ Aim for skilled attackers

- APT these days

## ☞ Create unpredictable & hostile environment

- ASLR
- Infoleak removal
- No RWX in memory or on disk

# Access Control != Security

- ⌘ Often, Access Control only plays a role post-exploitation
  - A “last line” of defense
- ⌘ Post-exploitation, an attacker wants permanence
- ⌘ Develop more complex exploit that plays along with published SELinux policies?
- ⌘ Attack that new perf\_counter system call completely unmonitored by SELinux?

# Access Control Won't Save You

⌘ Vmsplice

⌘ Aout loader

⌘ Tee/splice

⌘ Brk

⌘ Perf\_counter

⌘ Mremap

⌘ Move\_pages

⌘ Pipe

⌘ ELF loader

⌘ Etc...

# Kernel in the TCB

- ∞ Lots of attention paid to hardening userland
- ∞ Nearly no mainline attention to the kernel
- ∞ What will attackers target?
- ∞ Enlightenment
  - Disables SELinux, TOMOYO, IMA, AppArmor, all other LSMs
  - Grants full root, full capabilities, works in Xen
  - Upcoming LXC/OpenVZ support, since:
    - “If you are inside a user\_namespace your capabilities will only be good for manipulating other objects [...] that you have created after you entered the user namespace”

# Lessons From Last Year's Exploits

- ✎ Only public exploits produce a change in public perception of security
- ✎ Kernel security wasn't suddenly horrible in 2009, I simply showed how horrible it's always been
- ✎ Unlike with Tavis v. Microsoft, I received no threats from Linux vendors
  - Although...
- ✎ In the end, stronger SELinux protections, stronger mmap\_min\_addr, much higher user awareness

# Decade TODO List (for you) pt.1

## ∞ Remove infoleaks

- Symbol information
- Slabinfo
- PAX\_USERCOPY

## ∞ Remove RWX from kernel

## ∞ Protect sensitive data

- Constify function pointers!
- IDT/GDT/syscall table/etc
- Vsyscall shadow table (see sgrakkyu's remote SELinux-disabling exploit)

# Decade TODO List (for you) pt.2

- ☞ Protect against invalid userland memory accesses in general
- ☞ Make refcount overflows unexploitable
  - Currently equivalent to use-after-free
- ☞ `kmalloc(sizeof(somestruct) * attacker_len)`
  - See recent `ethtool get_rxnfc()` vulnerability
- ☞ Basically, secure the kernel! Your super fine-grained security systems will thank you



# Payoffs

## ☞ PAX\_UDEREF

- Found likely oldest Linux bug ever ( $\geq$  v0.01)
- vgaarb direct userland dereference
- NVIDIA direct userland dereference

## ☞ PAX\_KERNEXEC

- Enlightenment won't run (nor (all?) other memory-corruption based public exploits)

## ☞ PAX\_USERCOPY

- Found heap-based ~64kb infoleak

## ☞ PAX\_MEMORY\_SANITIZE

- Found use-after-free in CONFIG\_NO\_BOOTMEM

# Think Next Generation

- ✎ ASLR is a simple, useful technique
  - Ineffective in several cases (ones mainline doesn't handle properly already, and others)
  - Statistics-based security
- ✎ Deterministic control flow integrity
  - So long ret2libc/ROP/any other name
- ✎ The syscall table is protected – how about those page tables?

# Discussion

∞ Into the lion's den!