

# Security through Distrusting

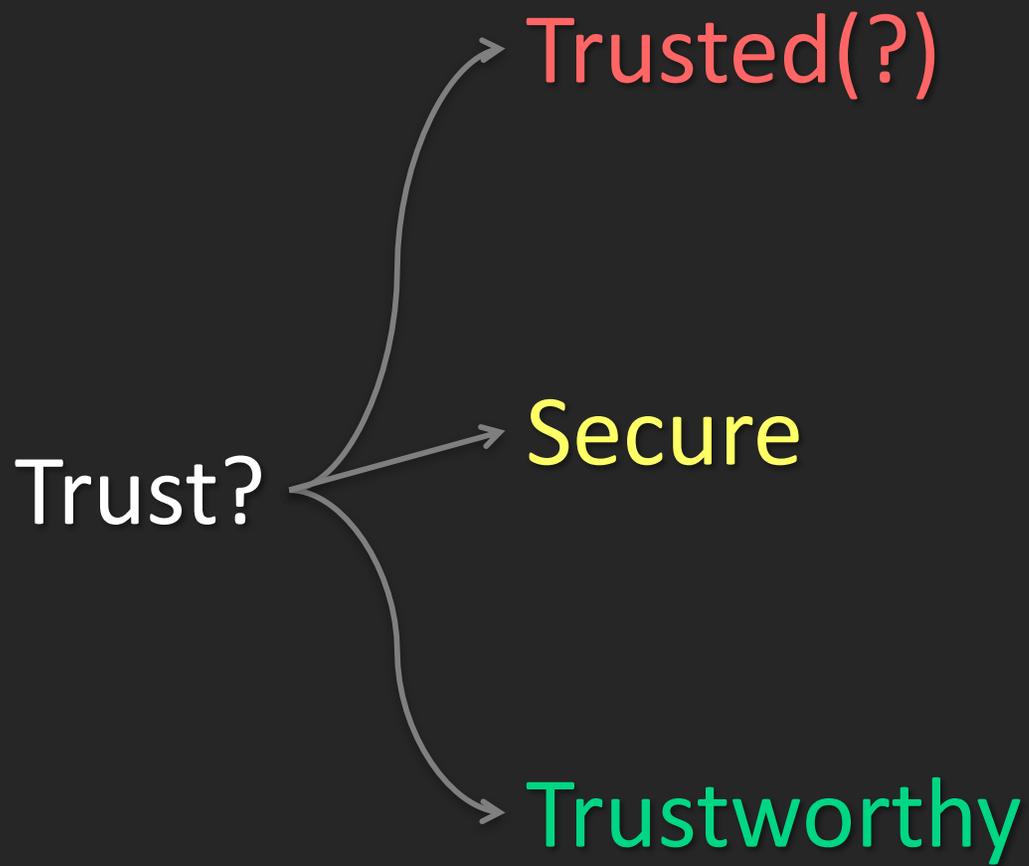
Joanna Rutkowska

Invisible Things Lab & [Qubes OS](#) Project

Black Hat EU, London, UK, December 7, 2017

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Trust consider harmful!

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# Security through Distrusting examples

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# Example #1: Pesky microphones

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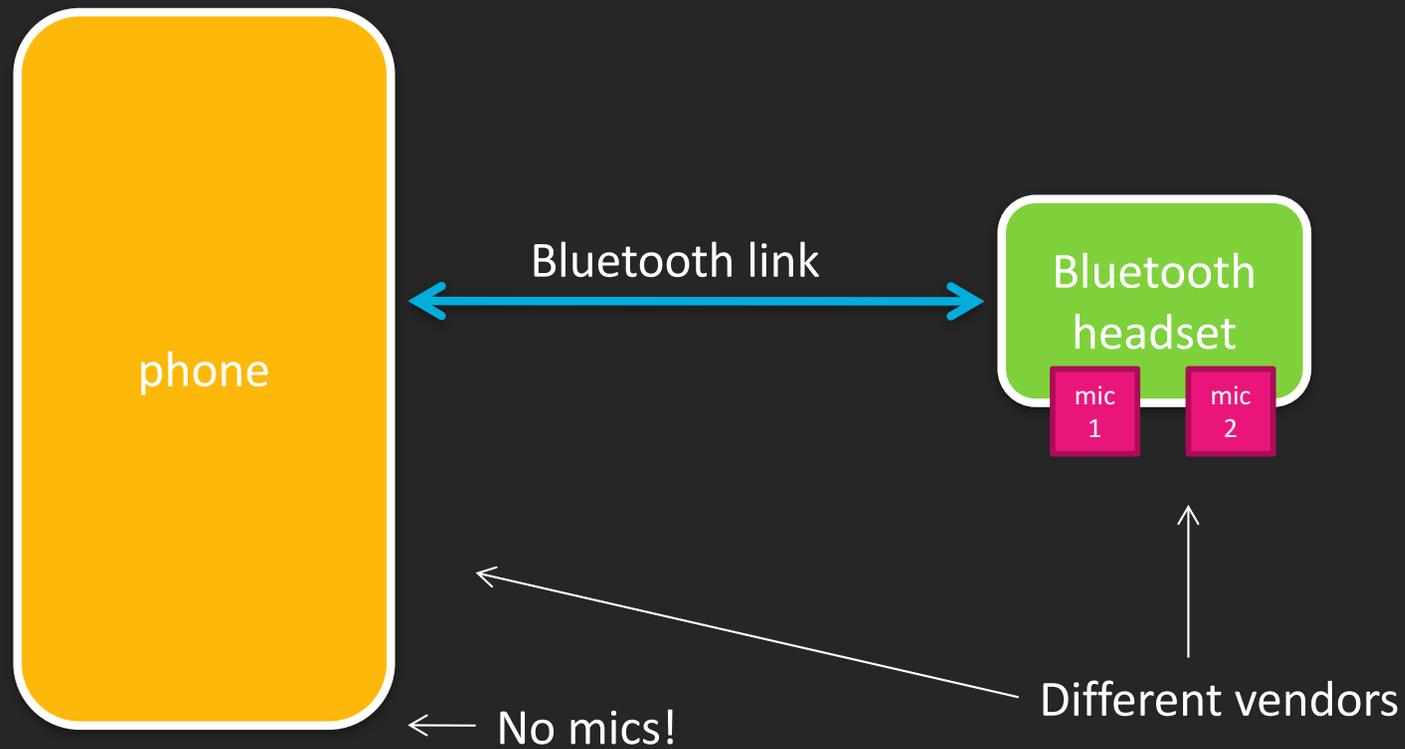




- Mics sniff our activities, including keystrokes, etc.
- Mics are difficult to neutralize
- Mics naturally “cross” security boundaries



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- Mics are difficult to neutralize
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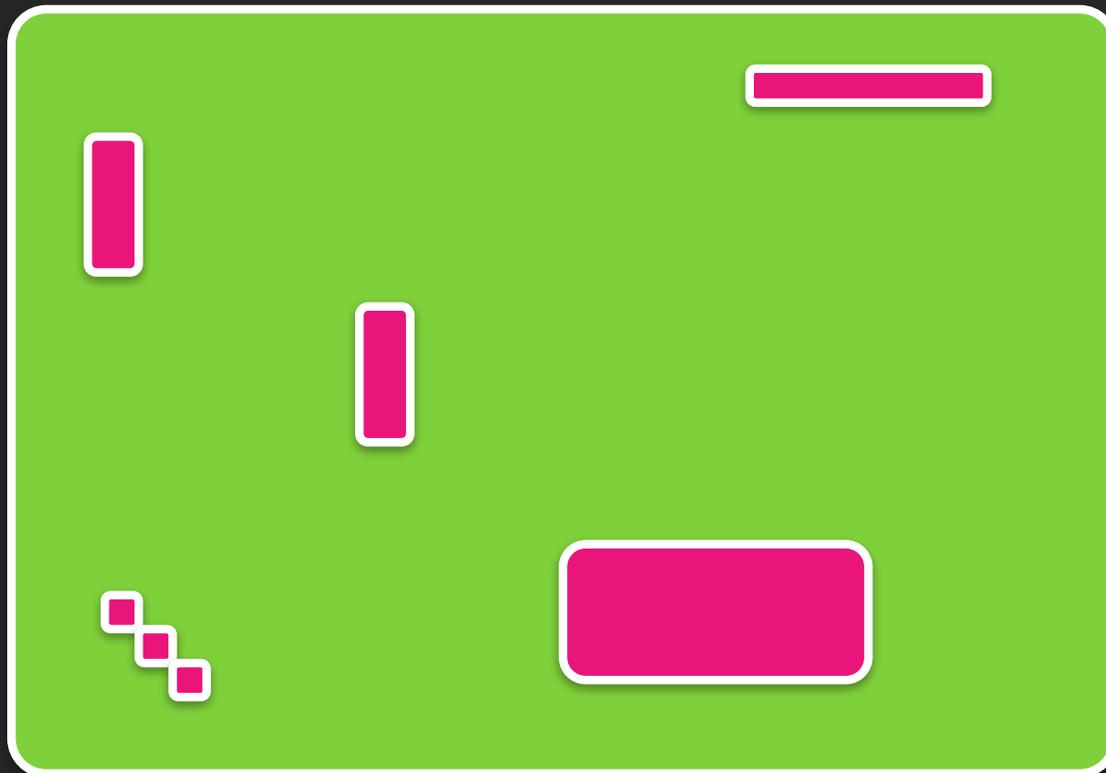
# Example #2: Stateless laptop

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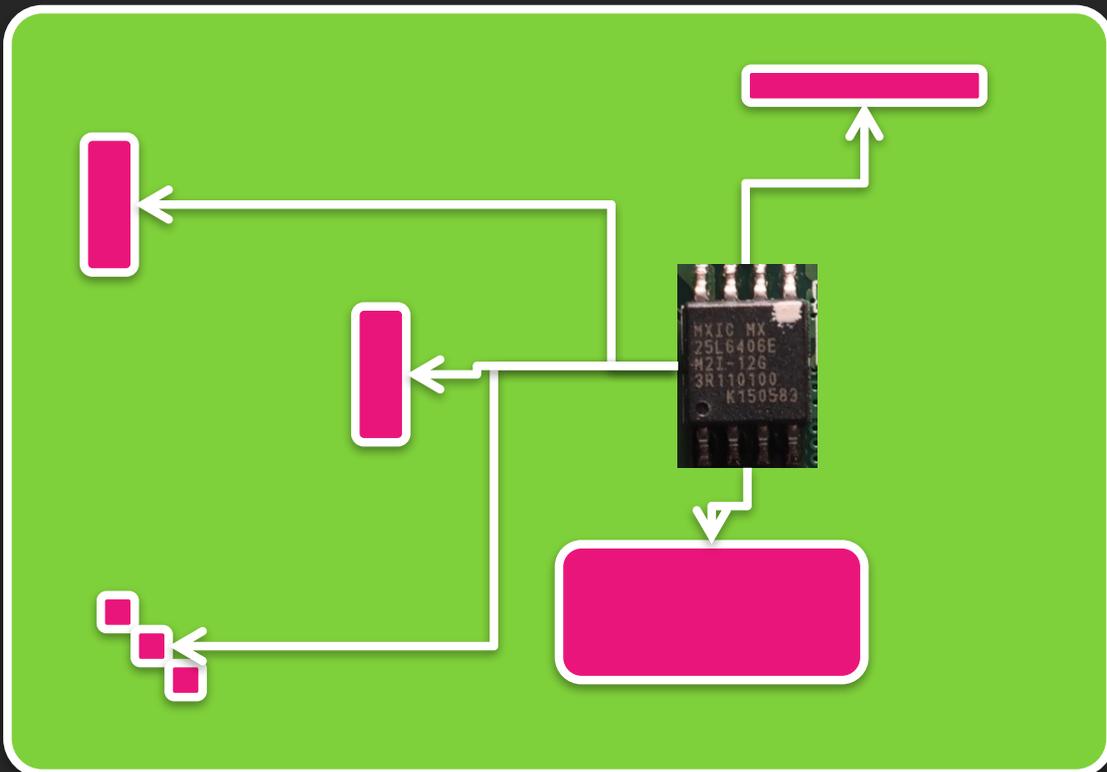


# Persistent laptop compromises...

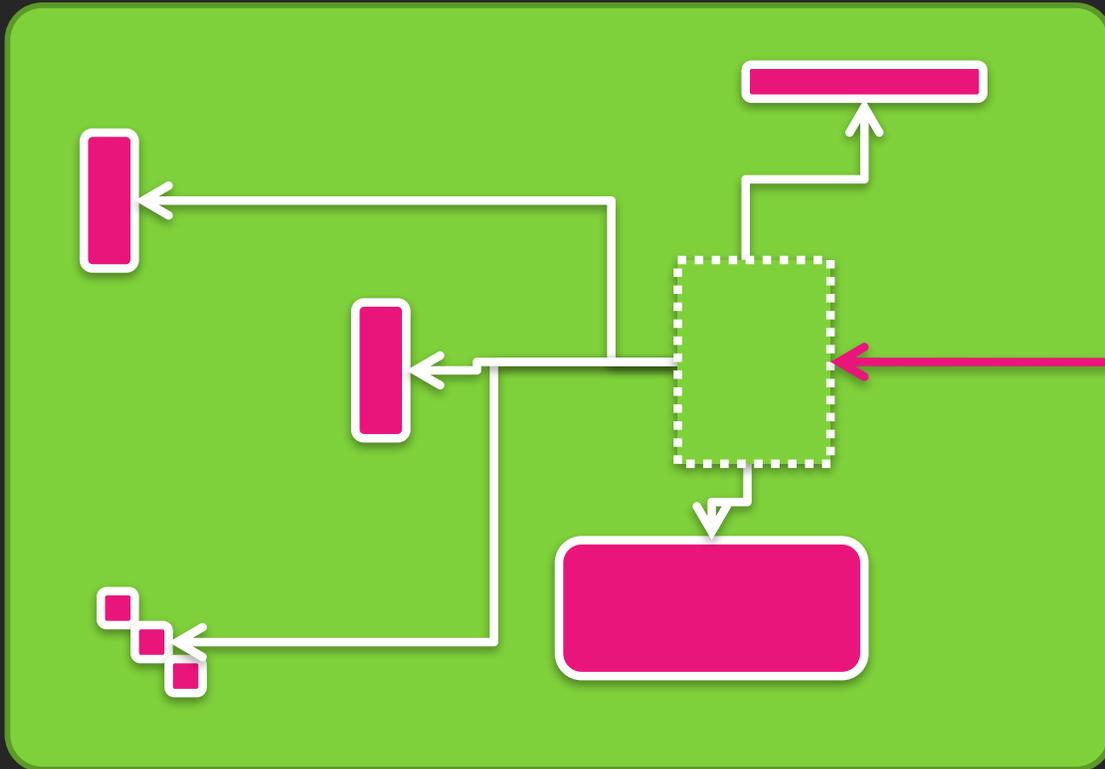
- Persist
- Store secrets
- PII



- Persist
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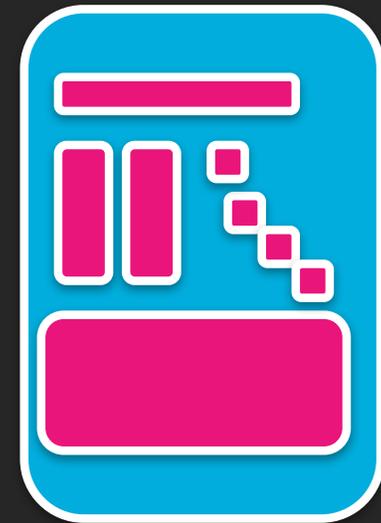


- Persist
- Store secrets
- PII



Stateless Hardware  
(persistent state eliminated)

- Firmware infections prevented
- No places to store stolen secrets
- Reliable way to verify firmware
- Reliable way to *choose* firmware
- Boot multiple environments
- Share laptops with others



Trusted Stick



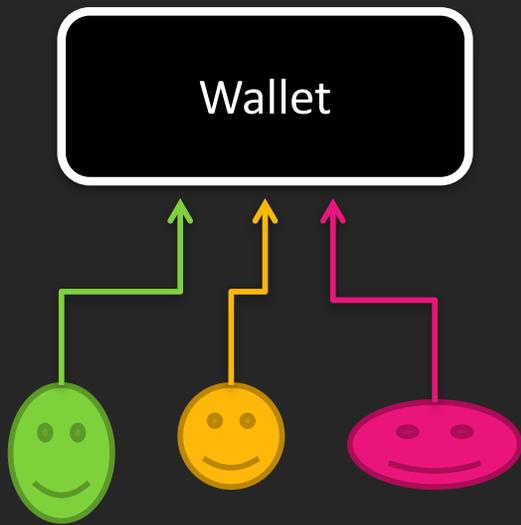
# Example #3: Multi-party signatures



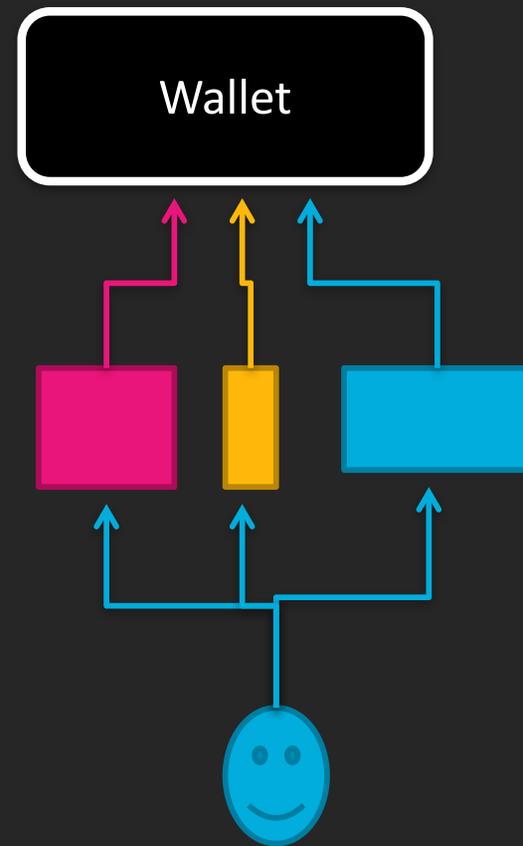
Photo via Peter Todd (@petertoddbtc)

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VS



Mutli-sig does not need to involve multiple users!



Also: not just Bitcoin wallets...

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# Example #4: Binary (multi) signing

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# Why care about binary (multi-) signing?

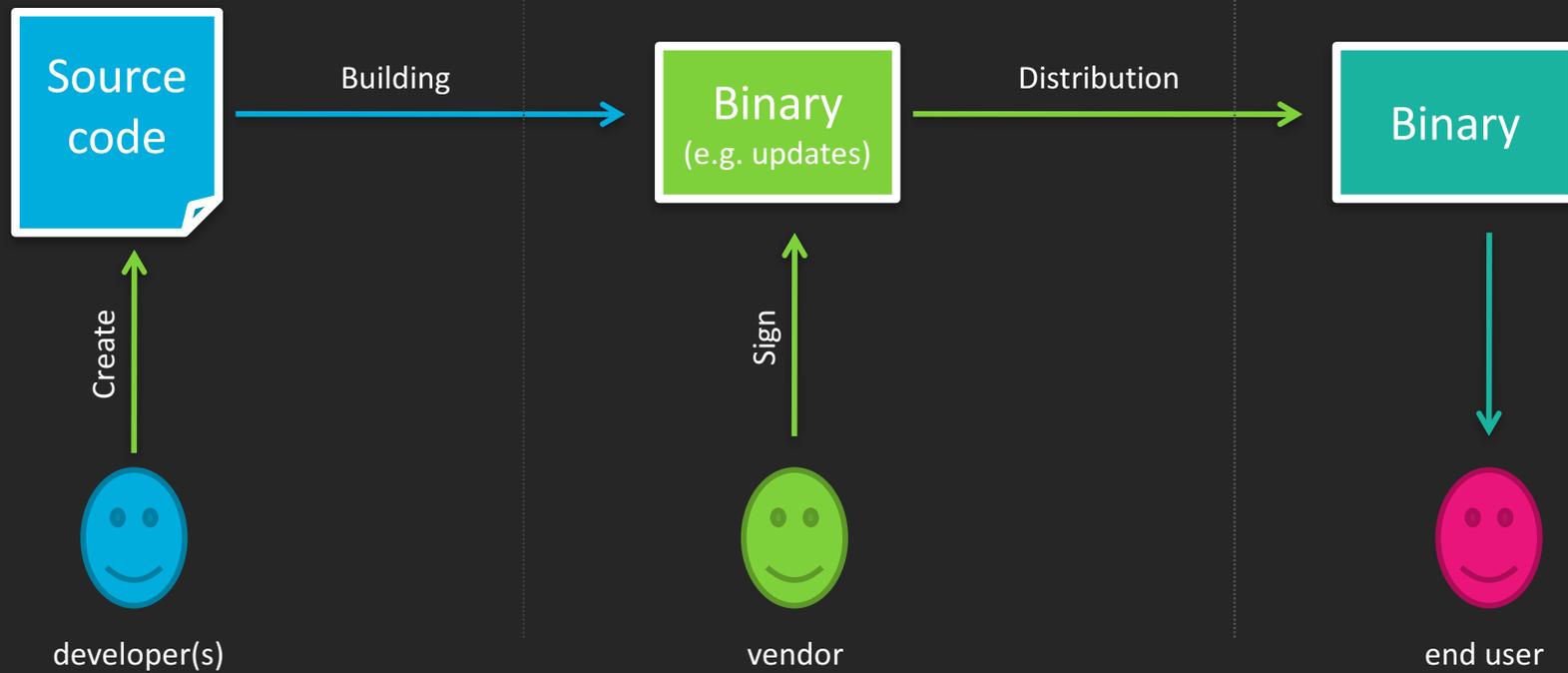
- OS installation images
- Applications
- Updates
- Firmware



Prime target for backdooring!

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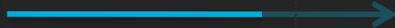


Source code

Create  
(... & sign!)



developer(s)



Binary  
(e.g. updates)

Sign



vendor



Distribution  
(https, signed binaries)

Binary

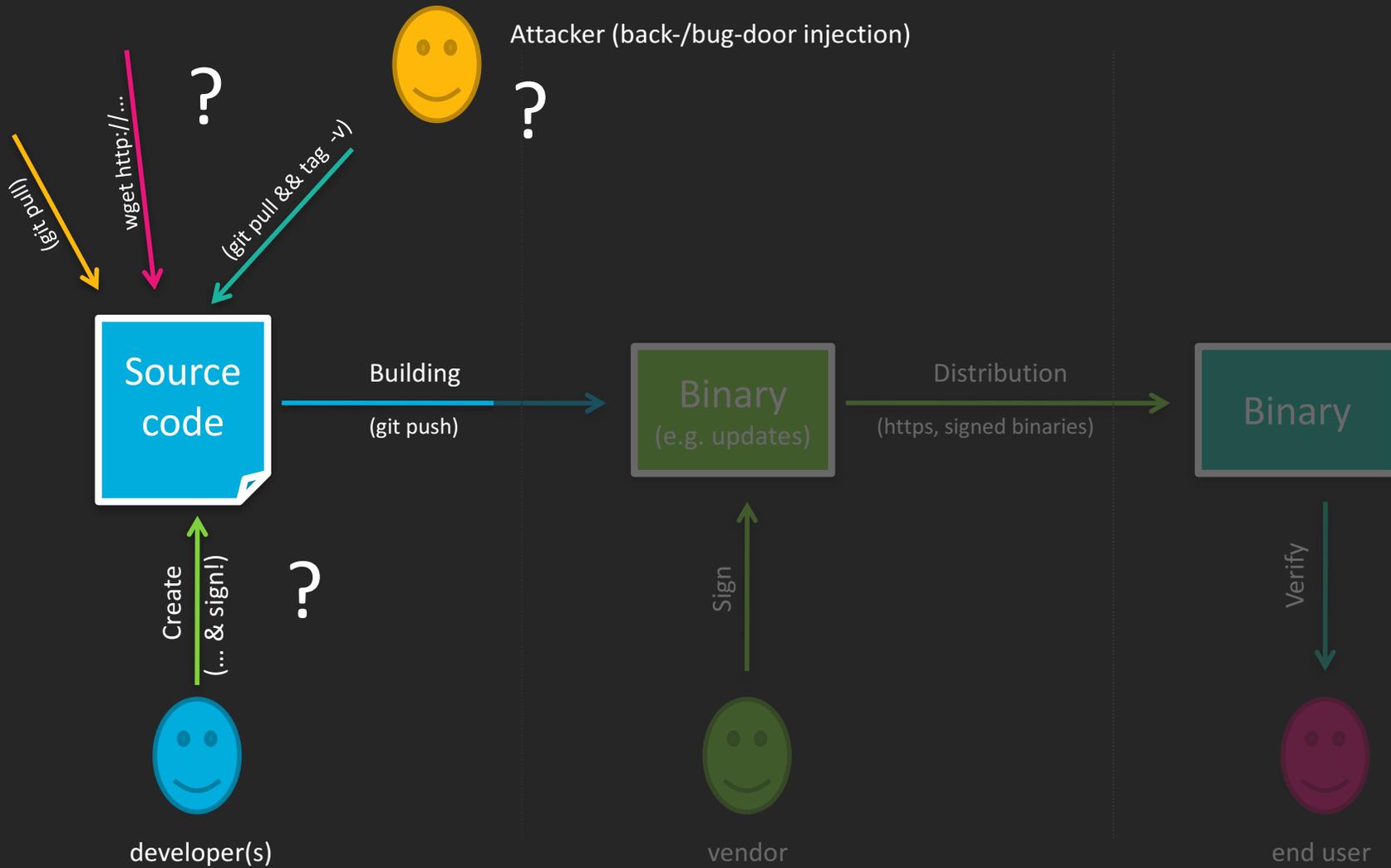
Verify



end user

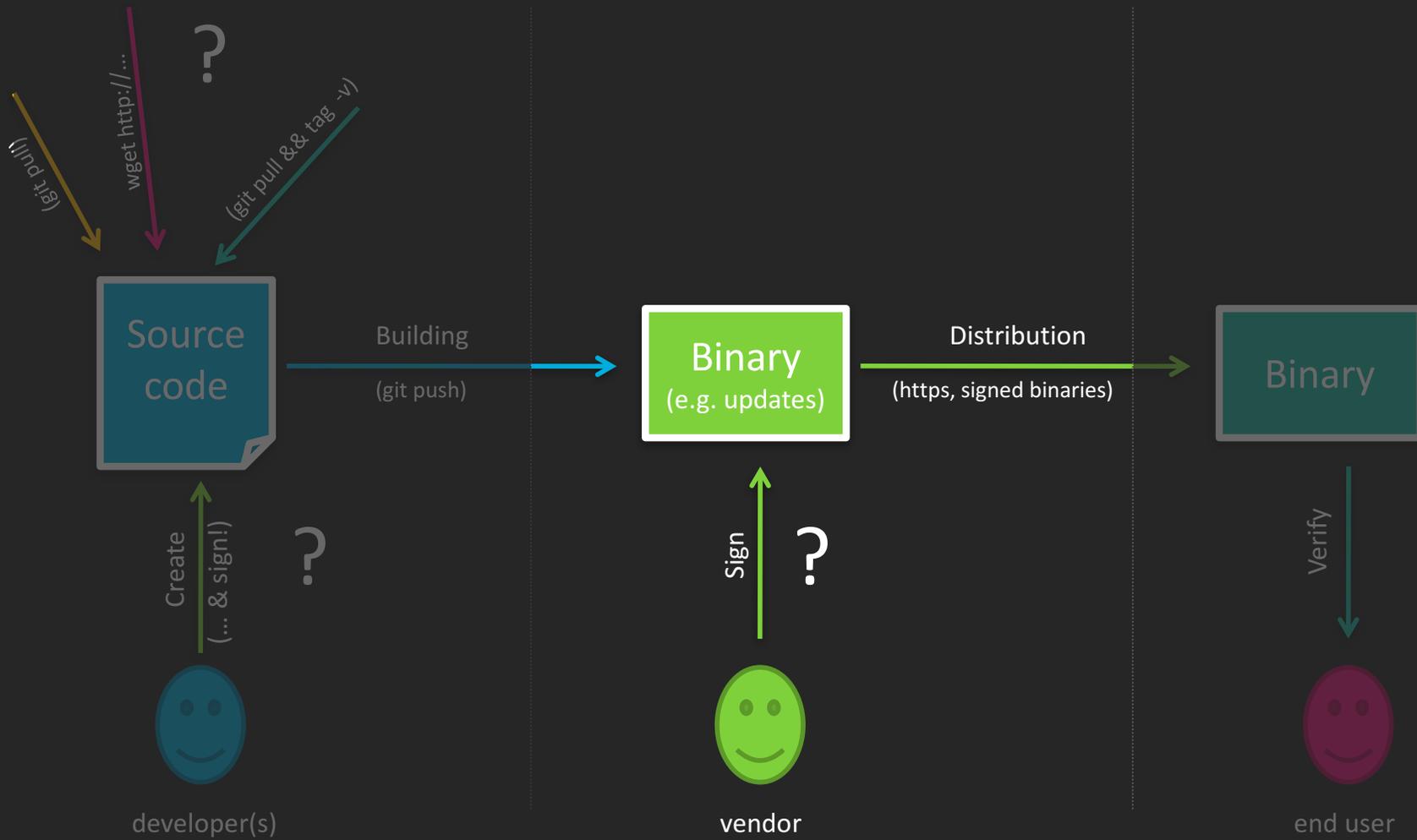


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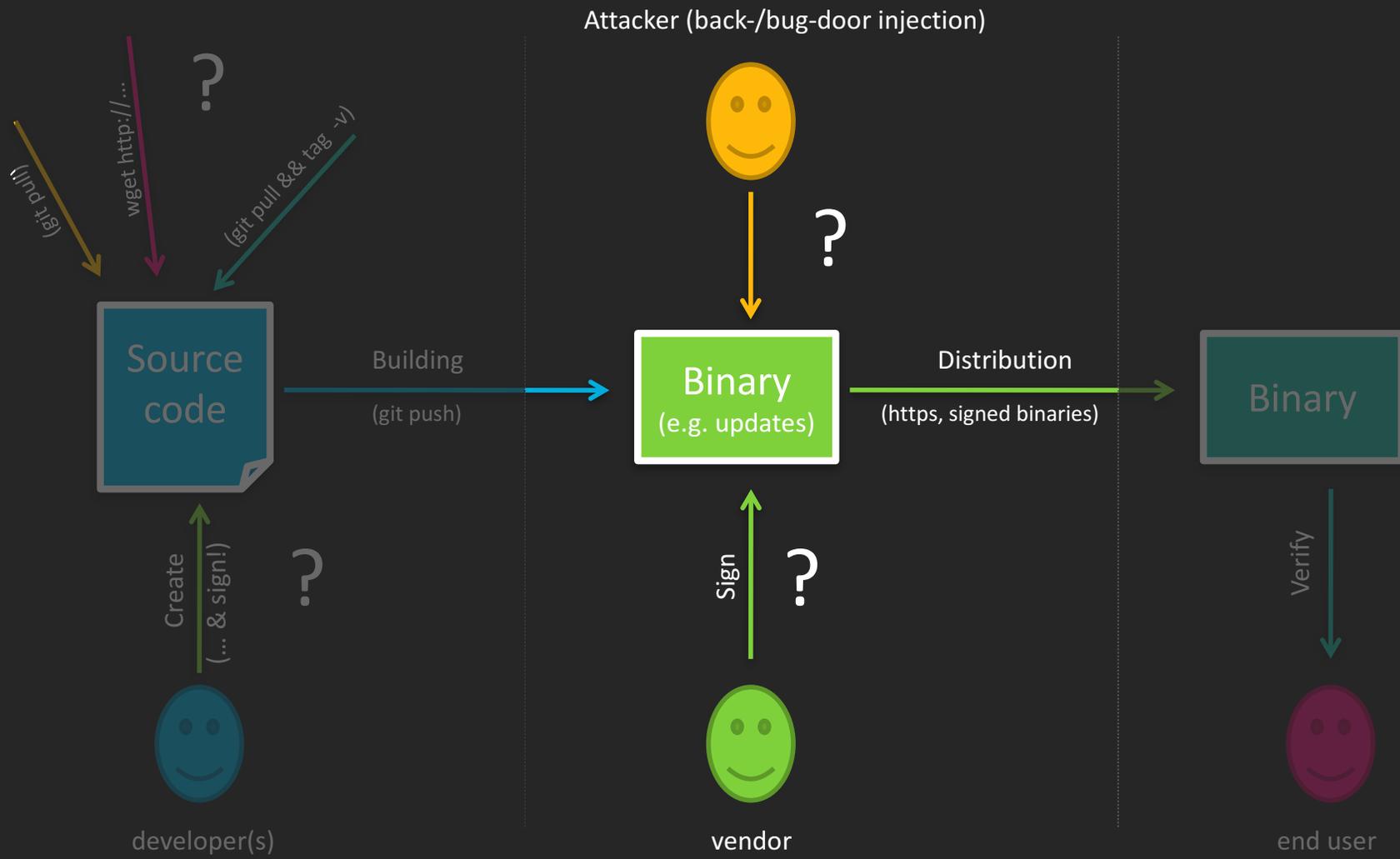
427F11FD 0FAA4B08 0123F01C DDF1A3E 36879494





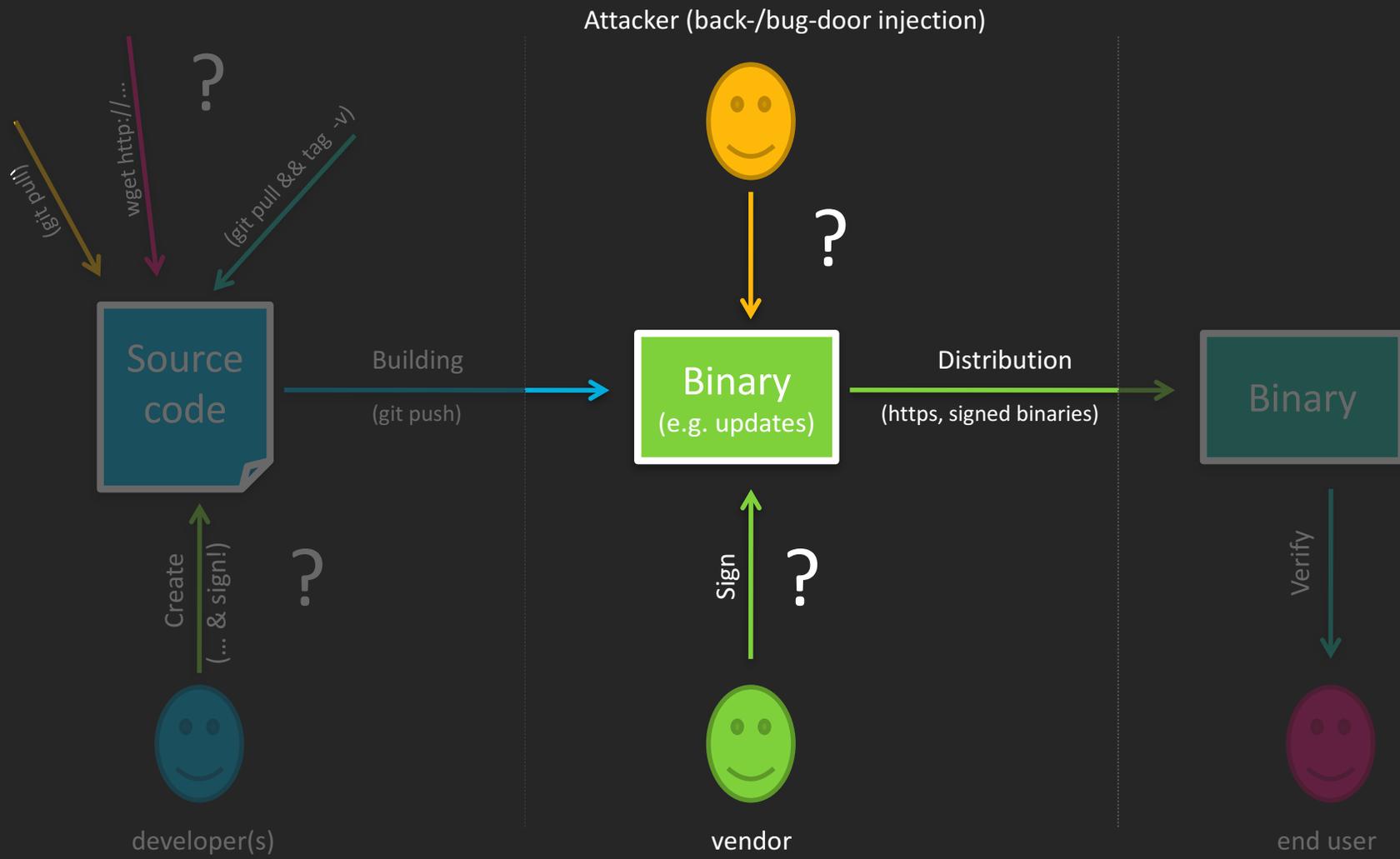
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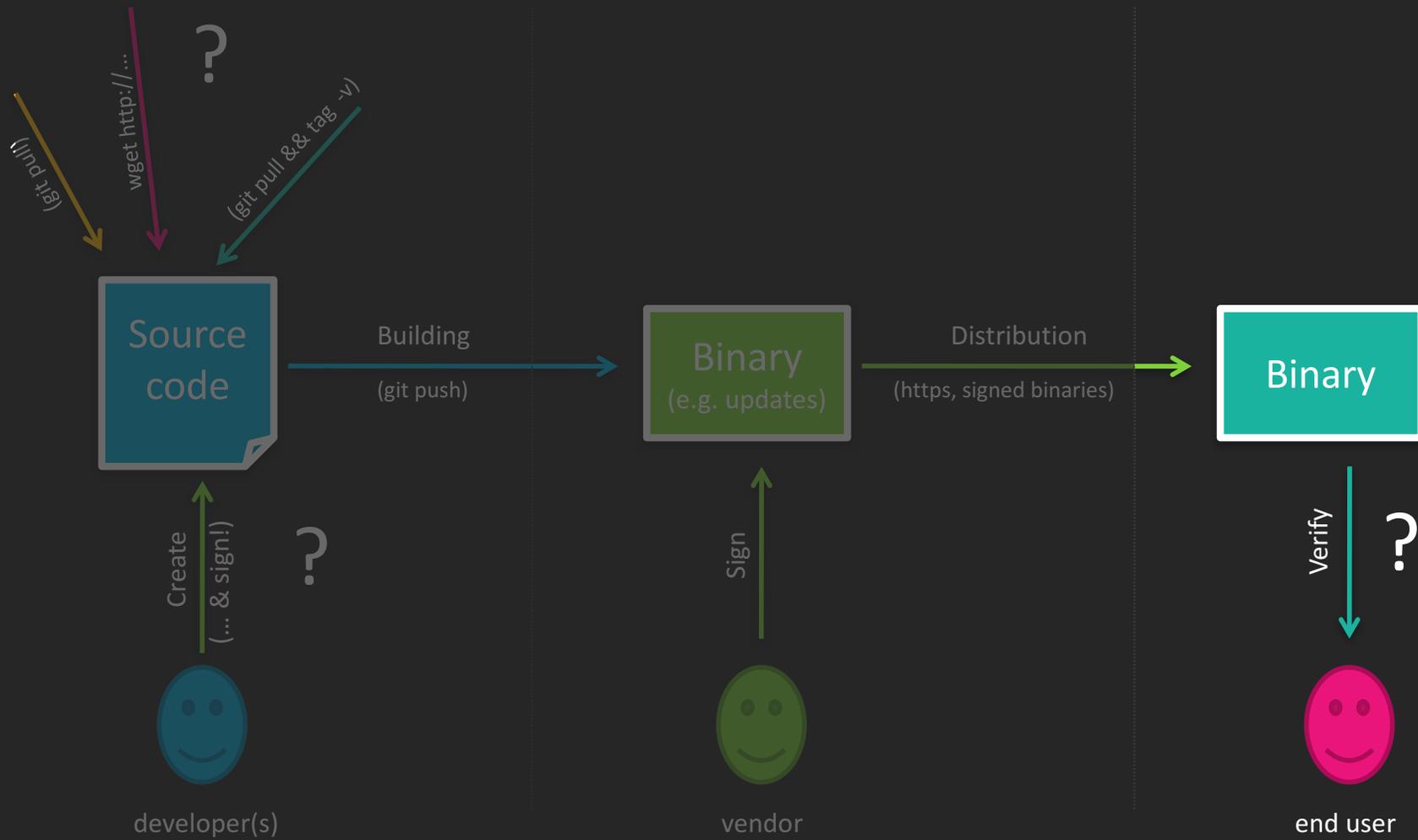
427F11FD 0FAA4B08 0123F01C DDFAL3E 36879494





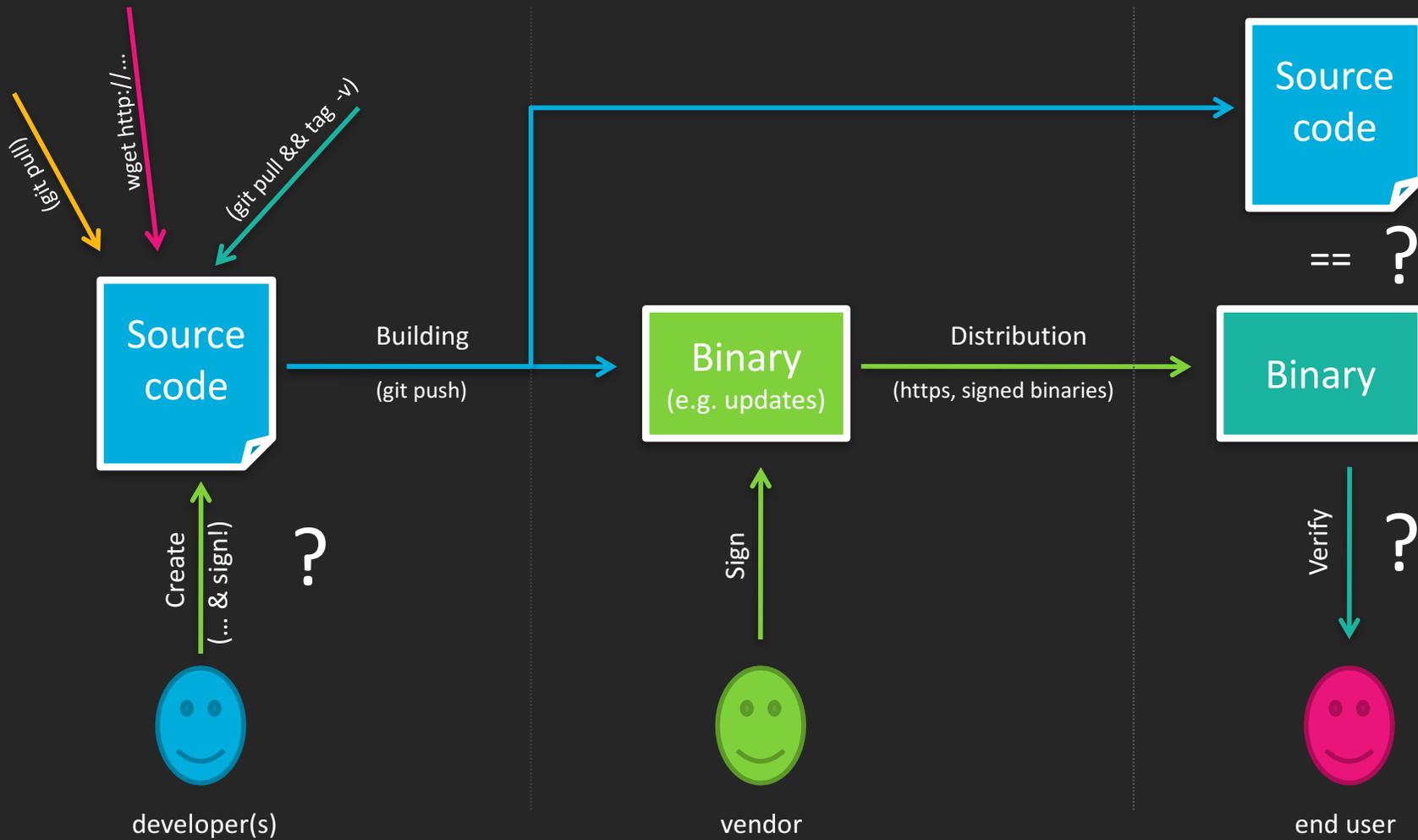
427F11FD 0FAA4B08 0123F01C DDFAL3E 36879494

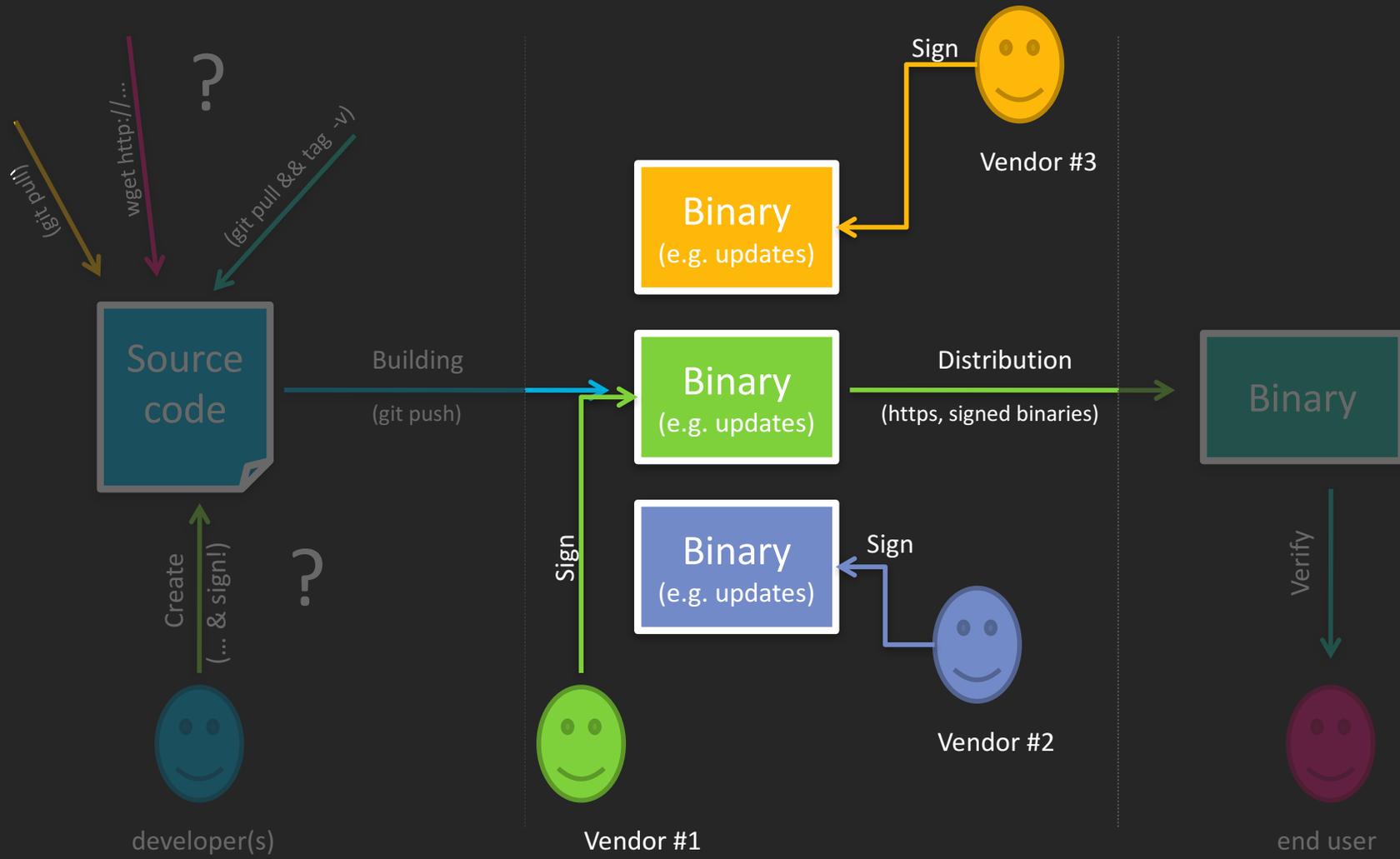




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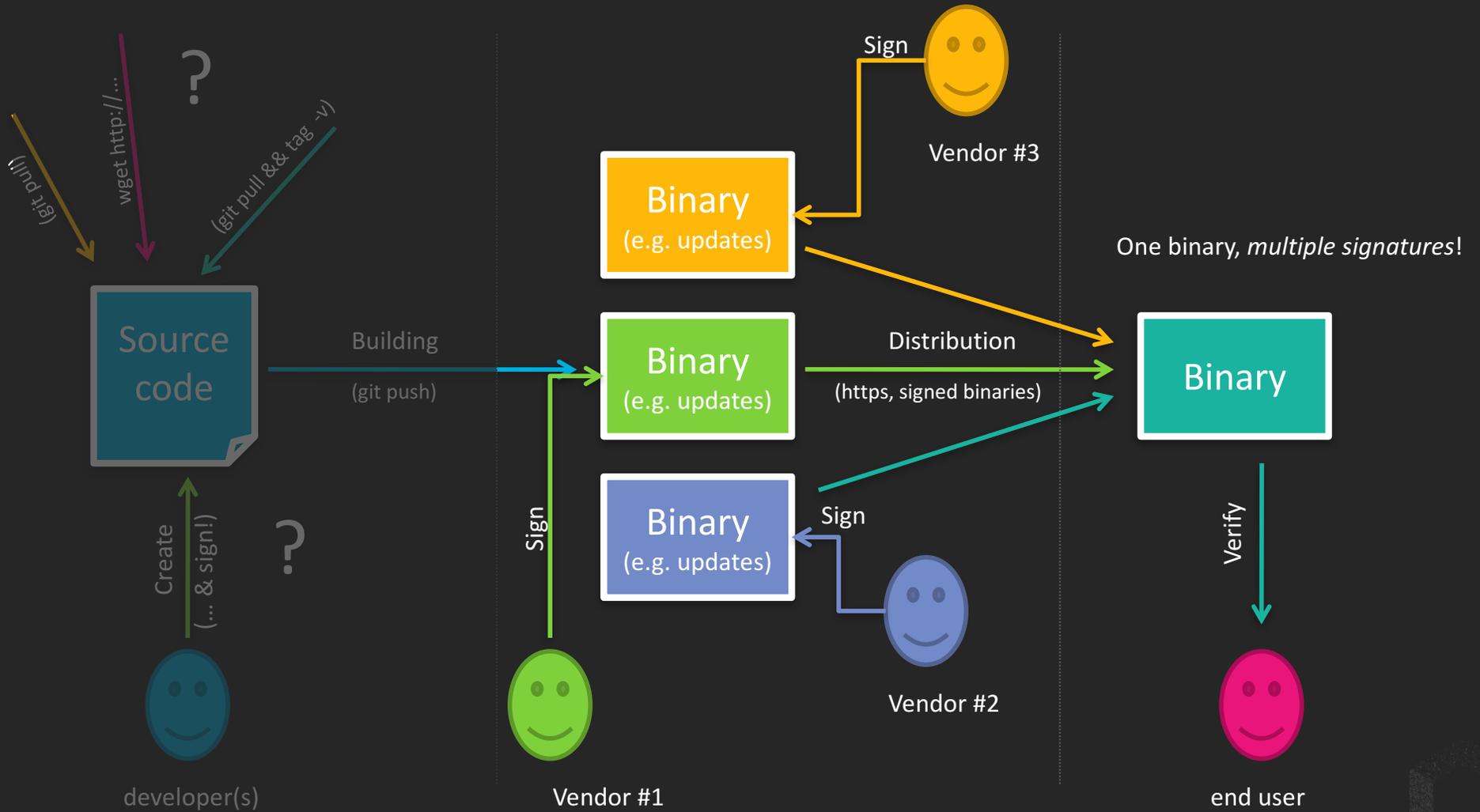






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# Multi-signed binaries

- Signed by people from different countries
- Different organizations (vendor & auditing)
- Signed by different *machines*
  - In the same organization
  - In different organization



<https://reproducible-builds.org>

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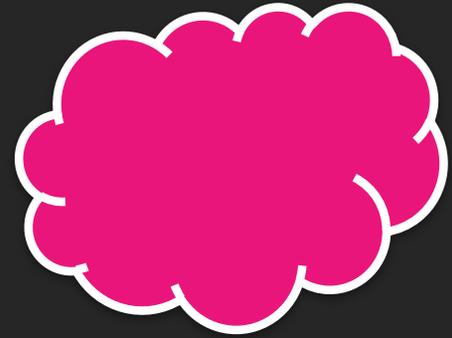
# Example 5: Preventing data leaks

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Some software  
(buggy/backdoored or  
otherwise compromised)

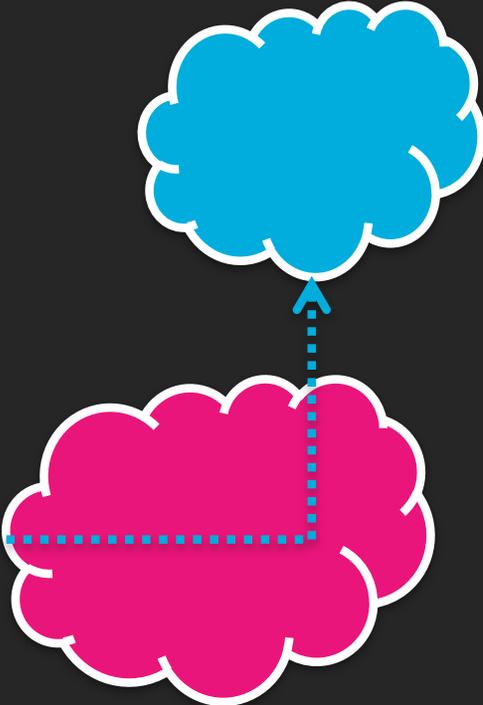
Your data...!



Some software  
(buggy/backdoored or  
otherwise compromised)

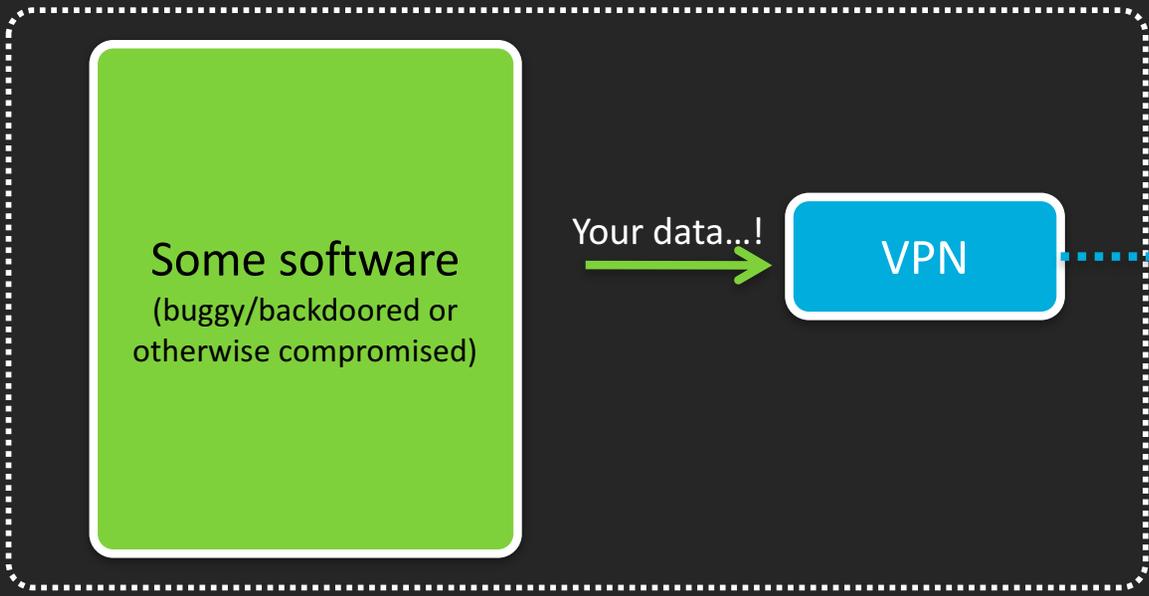
Your data...!

VPN



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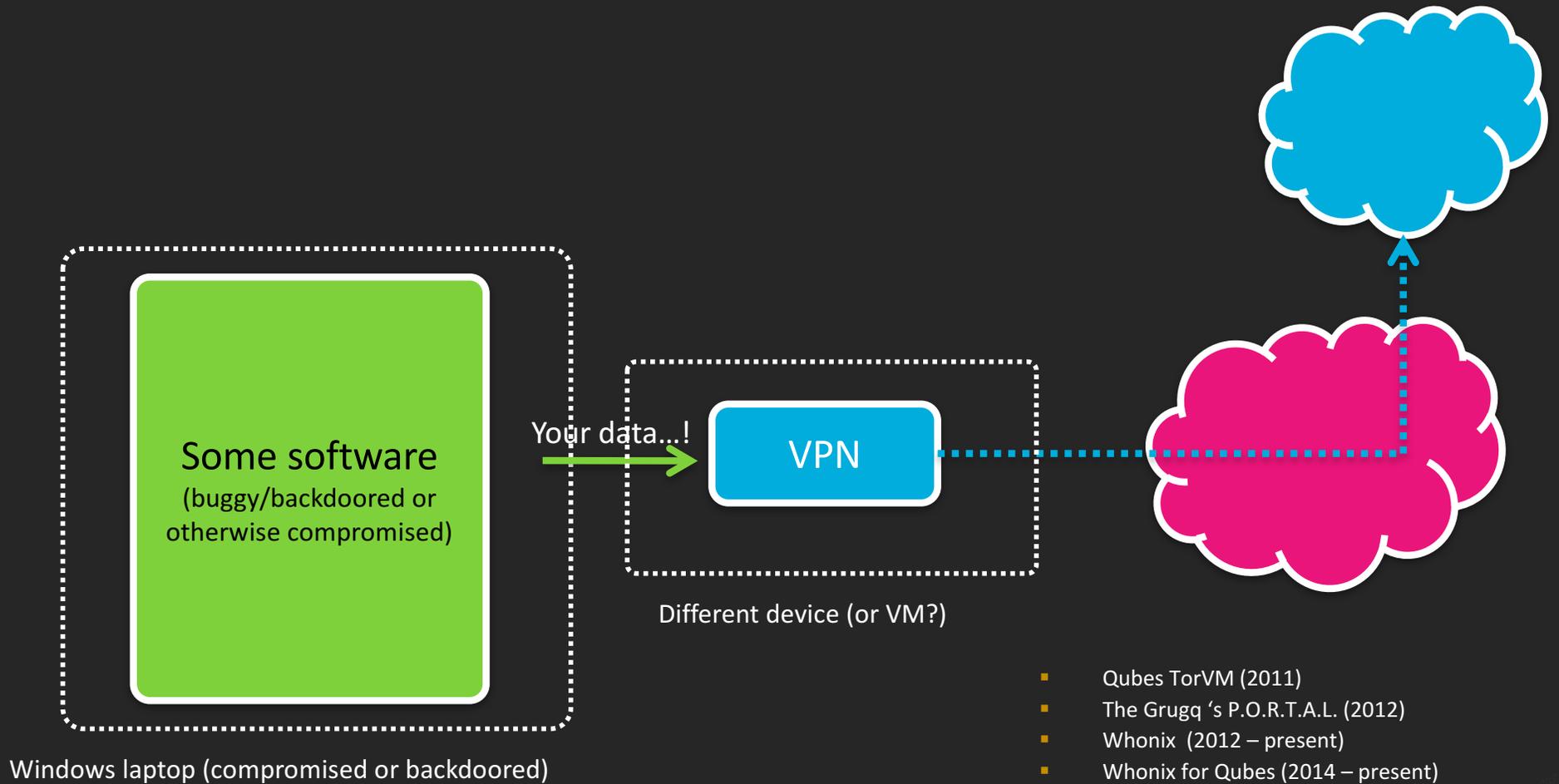




Windows laptop (compromised or backdoored)

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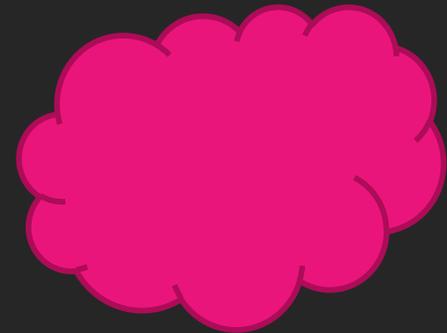
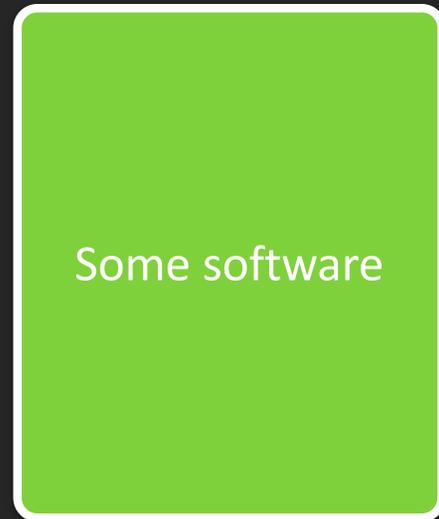




- Qubes TorVM (2011)
- The Grugq 's P.O.R.T.A.L. (2012)
- Whonix (2012 – present)
- Whonix for Qubes (2014 – present)
- Tor-enabled routers (multiple projects/products)



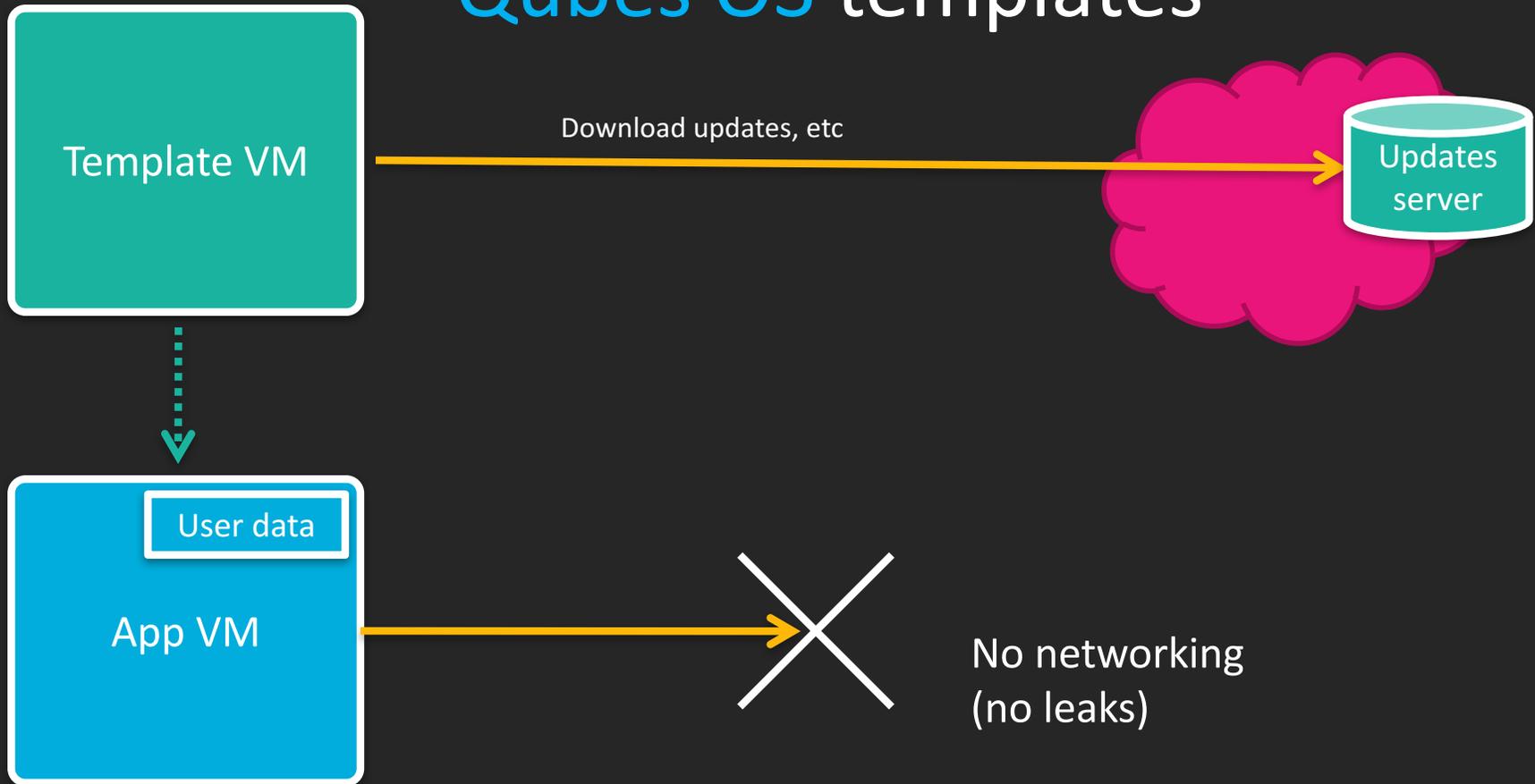
# Cut off networking?



Not very useful...



# Qubes OS templates



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# Example #6: Compartmentalization

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# “Classic” compartmentalization...

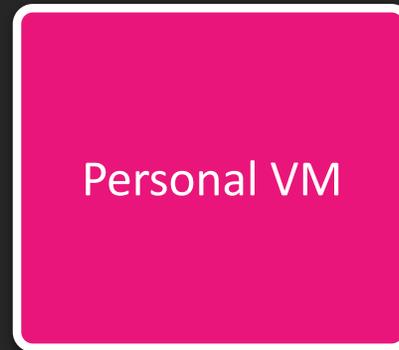
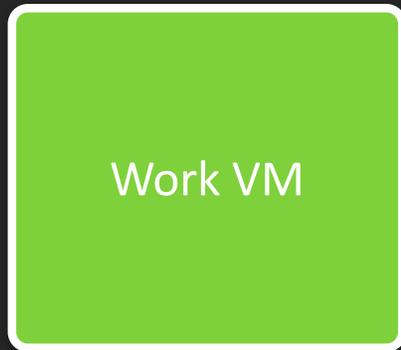
Work VM

Personal VM

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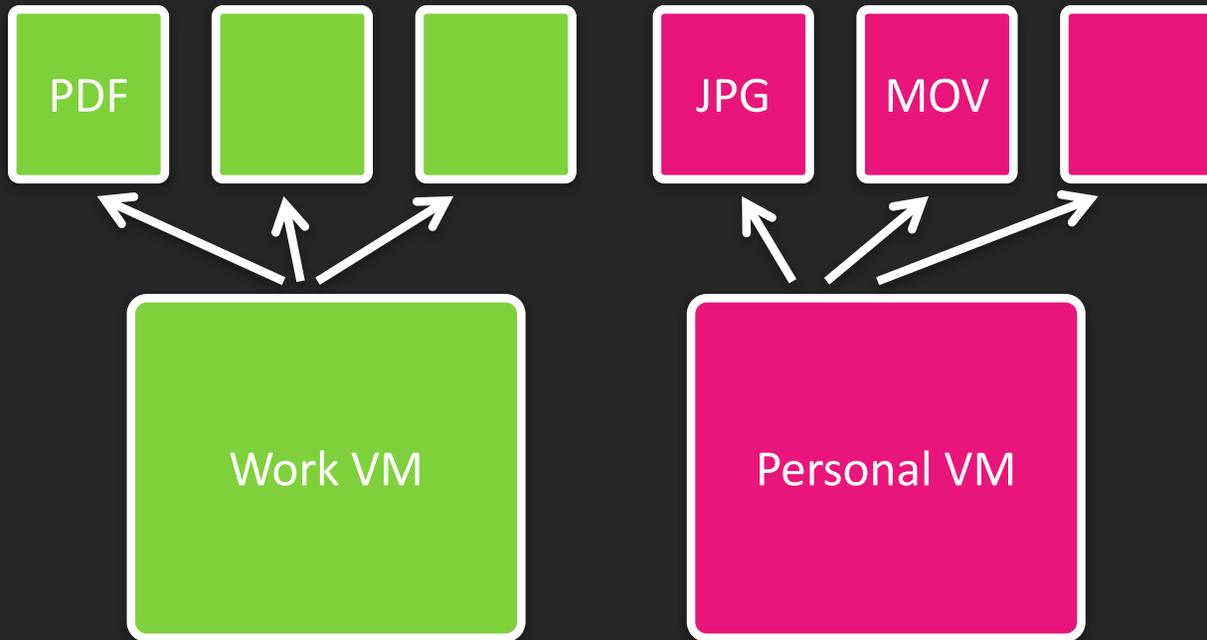


# “Classic” compartmentalization...



...not very useful!





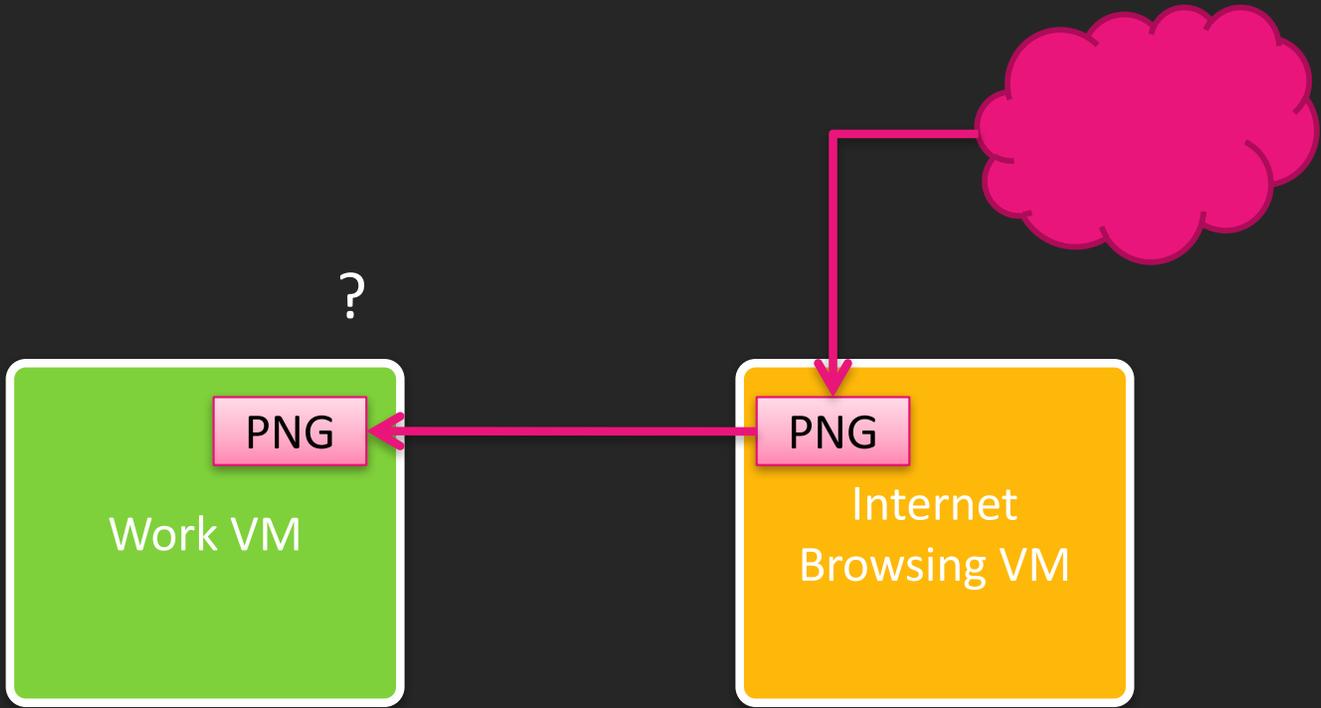
...more useful...



# Inter-compartments data transfers

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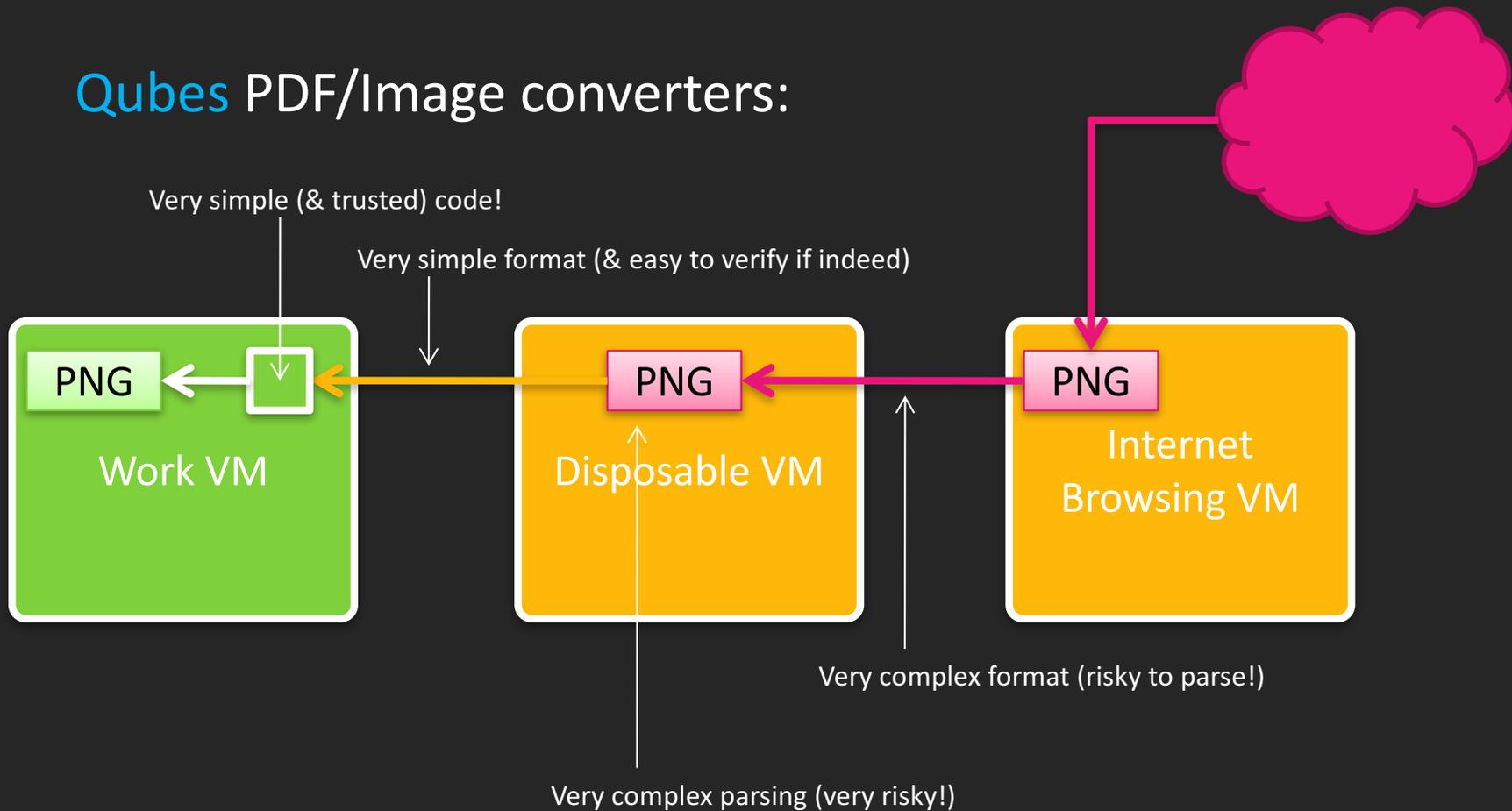




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# Qubes PDF/Image converters:



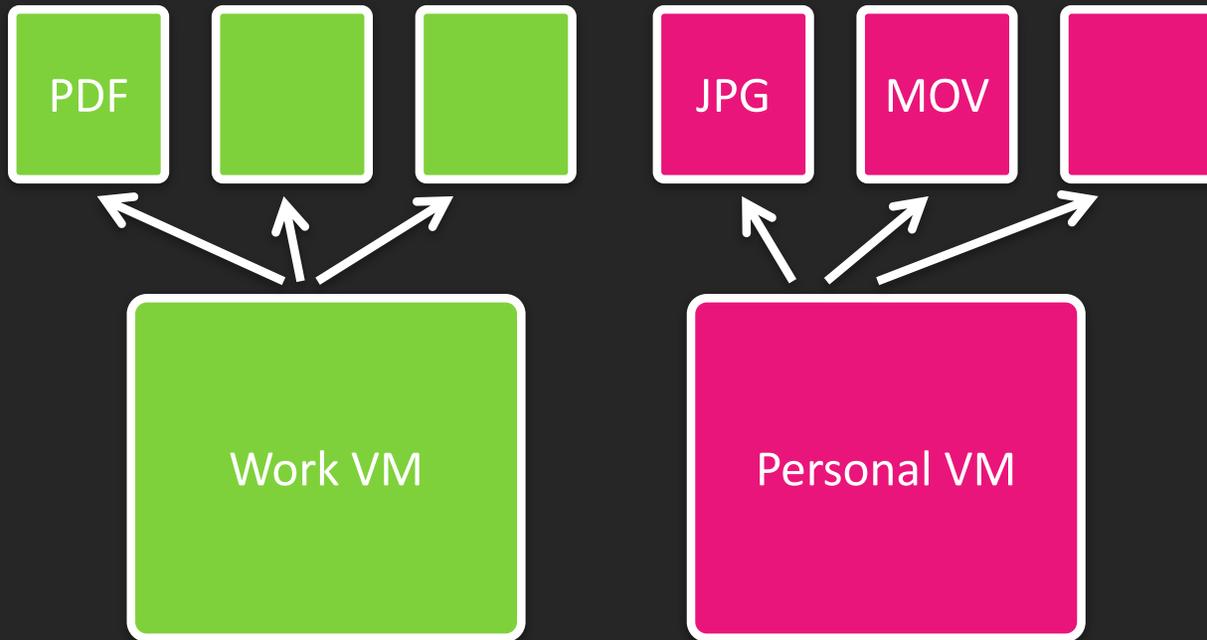
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App sandboxing is just part of the story...

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Networking stacks

USB & Bluetooth stacks

Graphics & UI

VPNs & firewalling

Corporate management

Root of trust (admin)



*Isolation* is just part of the story!

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### [varia] Software Guard Extensions Programming Reference

34 of 156 Software Guard Extensions Programming Refer... 88.66%

ENCLAVE OPERATION

After AEX has completed, the logical processor is no longer in enclave mode and the exiting event is processed normally. Any new events that occur after the AEX has completed are treated as having occurred outside the enclave (e.g. a #PF in dispatching to an interrupt handler).

#### 3.2.3 Resuming Execution after AEX

After system software has serviced the event that caused the logical processor to exit an enclave, the logical processor can re-start execution using ERESUME. ERESUME restores registers and returns control to where execution was interrupted.

If the cause of the exit was an exception or a fault and was not resolved, the event will be triggered again if the enclave is re-entered using ERESUME. For example, if an enclave performs a divide by 0 operation, executing ERESUME will cause the enclave to attempt to re-execute the faulting instruction and result in another divide by 0 exception. In order to handle an exception that occurred inside the enclave, software can enter the enclave at a different location and invoke the exception handler within the enclave by executing the EENTER instruction. The exception handler within the enclave can attempt to resolve the faulting condition or simply return and indicate to software that the enclave should be terminated (e.g. using EEXIT).

##### 3.2.3.1 ERESUME Interaction

ERESUME restores registers depending on the mode of the enclave (32 or 64 bit).

- In 32-bit mode (IA32\_EFER.LMA = 0 || CS.L = 0), the low 32-bits of the legacy registers (EAX, EBX, ECX, EDX, ESP, EBP, ESI, EDI, EIP and EFLAGS) are restored from the thread's GPR area of the current SSA frame. Neither the upper 32 bits of the legacy registers nor the 64-bit registers (R8 ...R15) are loaded.
- In 64-bit mode (IA32\_EFER.LMA = 1 && CS.L = 1), all 64 bits of the general processor registers (RAX, RBX, RCX, RDX, RSP, RBP, RSI, RDI, R8 ...R15, RIP and RFLAGS) are loaded.

Extended features specified by SECS.ATTRIBUTES.XFRM are restored from the XSAVE area of the current SSA frame. The layout of the x87 area depends on the current values of IA32\_EFER.LMA and CS.L:

- IA32\_EFER.LMA = 0 || CS.L = 0
  - 32-bit load in the same format that XSAVE/FXSAVE uses with these values.
- IA32\_EFER.LMA = 1 && CS.L = 1
  - 64-bit load in the same format that XSAVE/FXSAVE uses with these values plus REX.W = 1

### 3.3 CALLING ENCLAVE PROCEDURES

#### 3.3.1 Calling Convention

In standard call conventions subroutine parameters are generally pushed onto the stack. The called routine, being aware of its own stack layout, knows how to find parameters based on compile-time-computable offsets from the SP or BP register (depending on runtime conventions used by the compiler).

Because of the stack switch when calling an enclave, stack-located parameters cannot be found in this manner. Entering the enclave requires a modified parameter passing convention.

For example, the caller might push parameters onto the untrusted stack and then pass a pointer to those parameters in RAX to the enclave software. The exact choice of calling conventions is up to the writer of the edge routines; be those routines hand-coded or compiler generated.

#### 3.3.2 Register Preservation

As with most systems, it is the responsibility of the callee to preserve all registers except that used for returning a value. This is consistent with conventional usage and tends to optimize the number of register save/restore operations.

3-4 Ref. # 329298-001

### [Dom0] Qubes VM Manager

System VM View About

Name	State	NetVM	CPU Graph	MEM
dom0	●	n/a		2598 MB
sys-net	●	n/a		301 MB
sys-firewall	●	sys-net		301 MB
varia	●	sys-firewall		979 MB
work-web	●	sys-firewall		1173 MB
work-mutt	●	sys-firewall		604 MB
keys-iti-email	●	---		478 MB
work	●	sys-firewall		607 MB
personal	●	sys-firewall		750 MB

```
[work] user@work:~$
```

Biggest challenge for **Qubes OS** is  
how to do *desktop integration* (seamless UX)  
without compromising isolation!

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# Example #7: Almighty admins?

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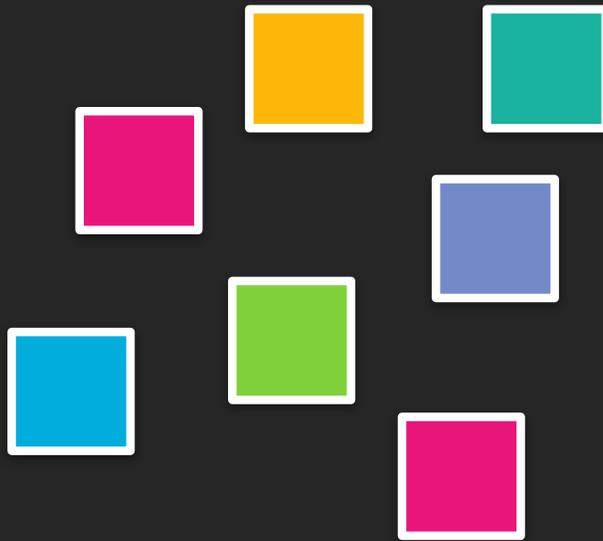
Admins can steal all our data :(

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Admin



- Access to their data
- Can't modify policies
- Can't modify software/VM images



User

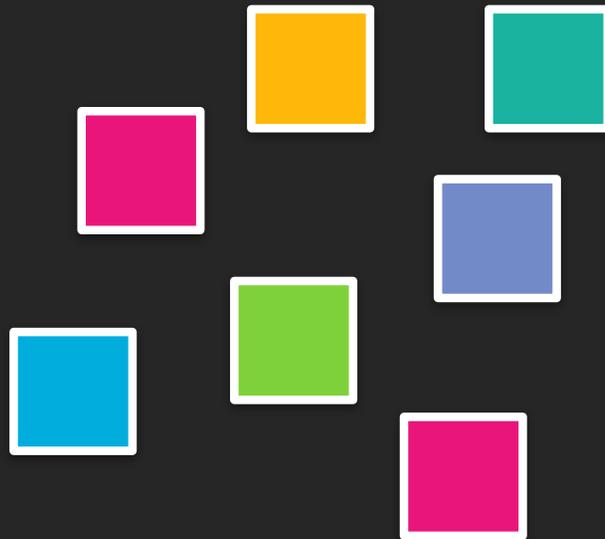




Admin



Can access and do *everything* she wants!



- Access to their data
- Can't modify policies
- Can't modify software/VM images



User



427F11FD 0FAA4B08 0123F01C DDF A1A3E 36879494



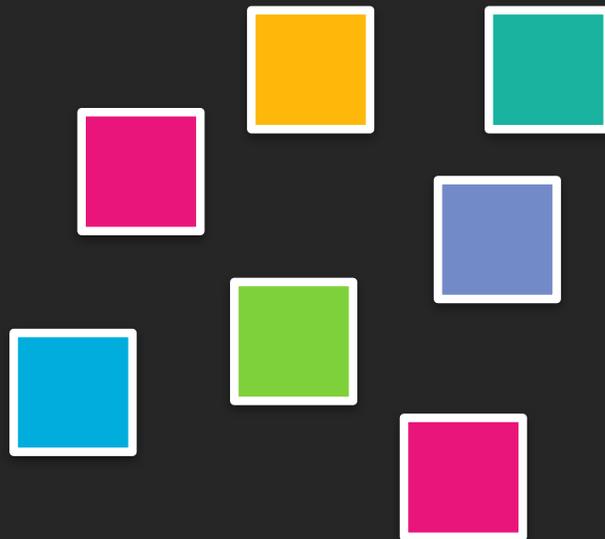
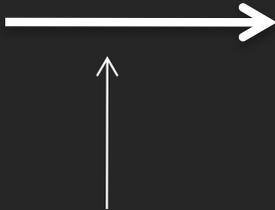
Hmm...

# What we want instead:



Admin

- No access to user data
- Can modify policies
- Can install software/VM images



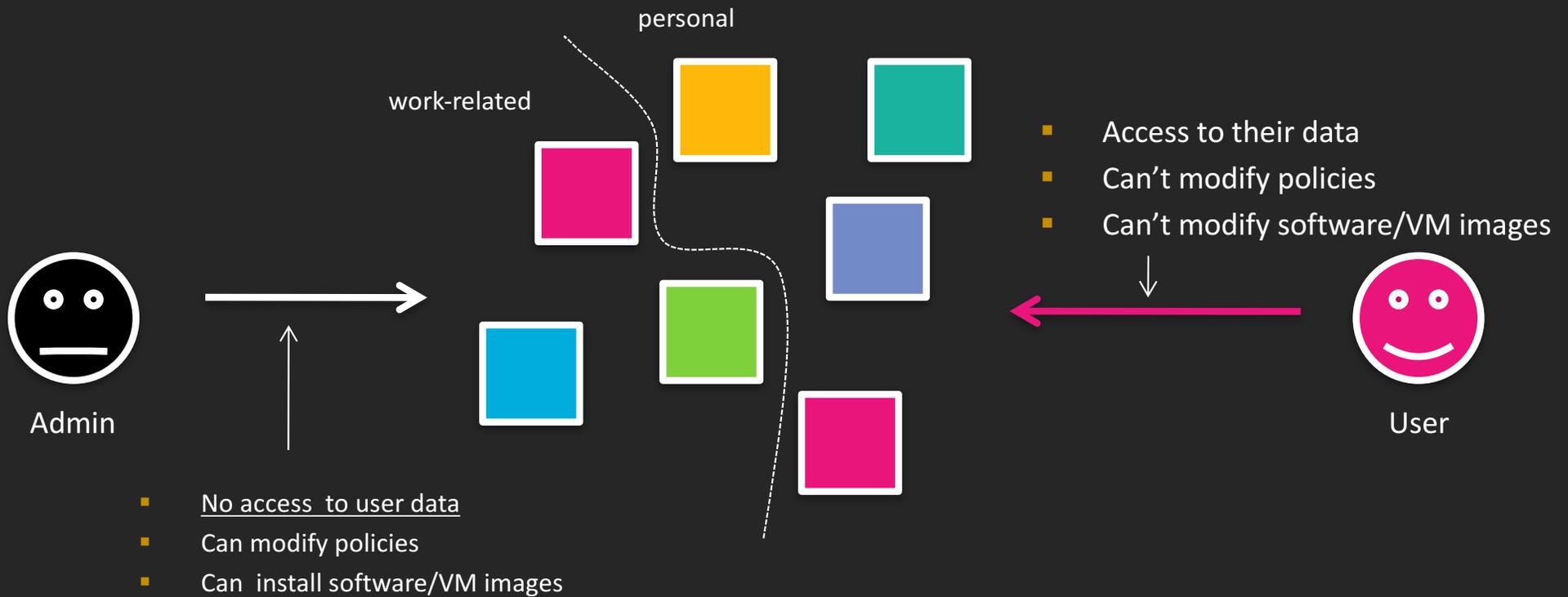
- Access to their data
- Can't modify policies
- Can't modify software/VM images



User



# What we want instead:



Check our [Qubes OS](#) new Admin API for implementation details



Occasionally mishaps happen still...

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# Example #8: Plan B

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Compromised machine



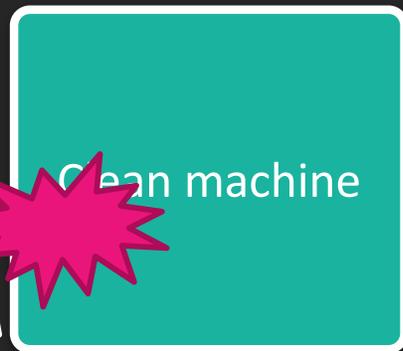
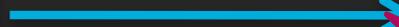
Backup  
(files, disks)



Clean machine

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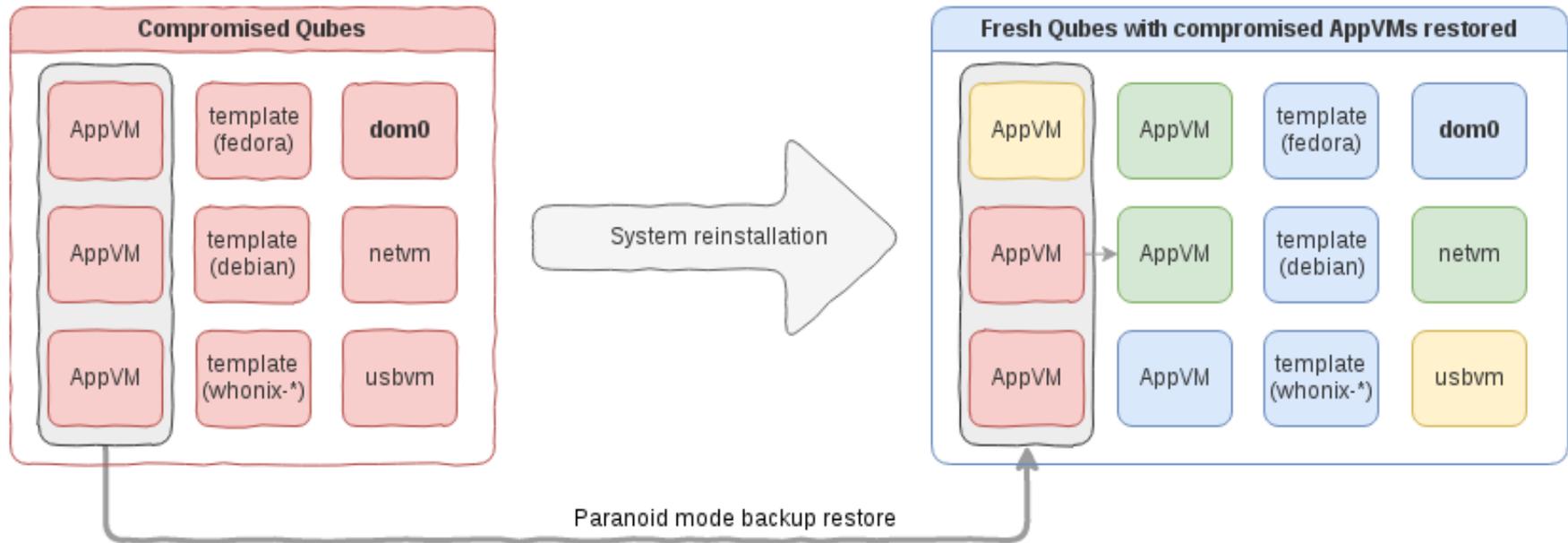




Restoring compromised backup is risky!



# Qubes (Paranoid) Backup Restore



# Security through Distrusting



## Division of Duty

- Mics (#1)
- Stateless laptop (#2)
- Multi signatures (#3/4)
- Tunneling (#5)



## Compartmentalization

- Qubes (#6/7)
- Tunneling (#5)
- Qubes Backup Restore (#8)



## Plan B having

- Qubes Backup Restore (#8)

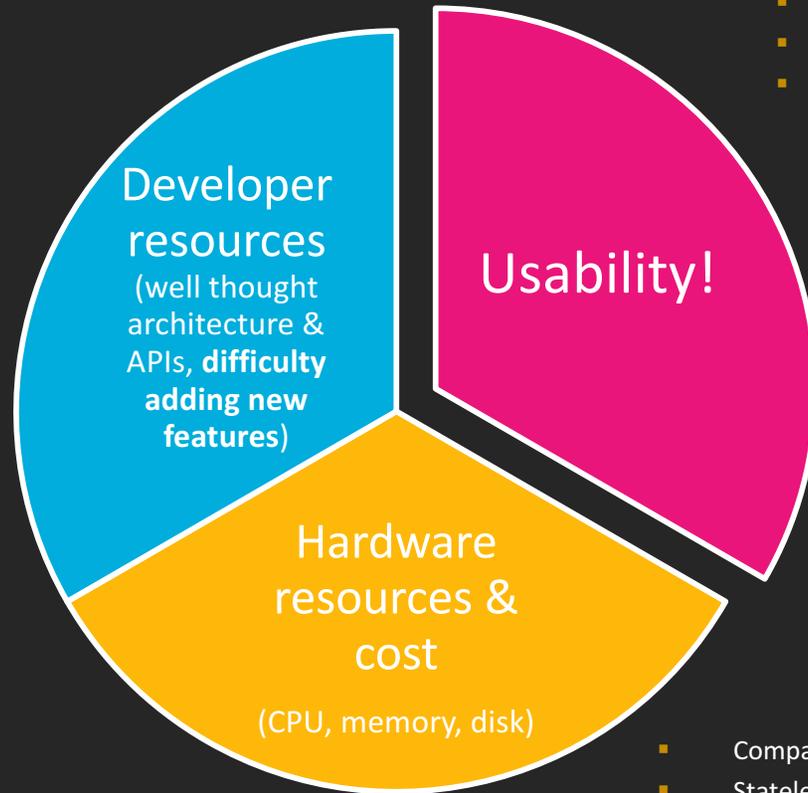


# Tradeoffs?

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- Compartmentalization
- Multisigs for binaries
- Stateless laptop (BOM costs)



- Mics
- Multisigs for wallets
- Compartmentalization (?)

- Compartmentalization
- Stateless laptop (BOM costs)



# Thanks!

\*\*\*

<https://qubes-os.org>

<https://invisiblethingslab.com>

<https://blog.invisiblethings.org>

<https://github.com/rootkovska>

@QubesOS // Twitter for Qubes OS

@rootkovska // Personal Twitter

427F11FD0FAA4B080123F01CDDFA1A3E36879494 // Qubes Master Key

ED727C306E766BC85E621AA65FA6C3E4D9AFBB99 // Personal Master Key

427F11FD 0FAA4B08 0123F01C DDF1A3E 36879494

